NSRIT Autonomous

ANSWER KEY & SCHEME OF EVALUATION

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B. Tech. (S3 Supplementary April 2022

ACADEMIC Regulation 2020

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Nadinpali Savananyatiz Roju Include of Technology (Autopologies, 10AG, 19ually Ivanaganan' System (OWs

NSRI

Semes	er End Supplem	entary Examination, April/May,	2 022	
Degree B. Tech. (U.G.)	Program		Academic Year	2021 - 2022
Course Code 20CE304			Semester	
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10 (b) Derive an expression expre	ss distribution for rectangular section.	6M	20CE304.2	<u></u> 3
A beam of lengt two point loads 11 respectively from load (ii). The po	e beam method.	6M	20CE304.3	
two point loads 11 respectively from load (ii). The po	ssion for the slope and deflection of a Simply with a point load at center.	6M	20CE304.3	Ľ
	OR n 6m is simply supported at its ends and carries of 48 kN and 40 kN at a distance of 1m and 3m n the left support. Find (I) Deflection under each int at which the maximum deflection occurs and tion using Macaulay's method.	: : 12M	20CE304.3	·
12 Derive an expre column are hing	The second s	12M	20CE304.4	Ľ
× 100 mm is ri	OR teel column of rectangular cross-section 120 mr gidity fixed at one end and hinged at the othe uckling load on the column and the correspondin	r. g 12M	20CE304.4	Ľ

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The stiffness of a closely coiled helical pring is 1.5 N/mm of compression under a maximum load of 100N. The maximum shearing stress produced in the wire of the spring is 130 N/mm². The solid length of the spring (when the coils are 112M touching) is given as 5cm. Find (I) Diameter of the wire (ii) Mean diameter of the coils and (iii) No. of coils required Take C=4.5X10⁴ N/mm².

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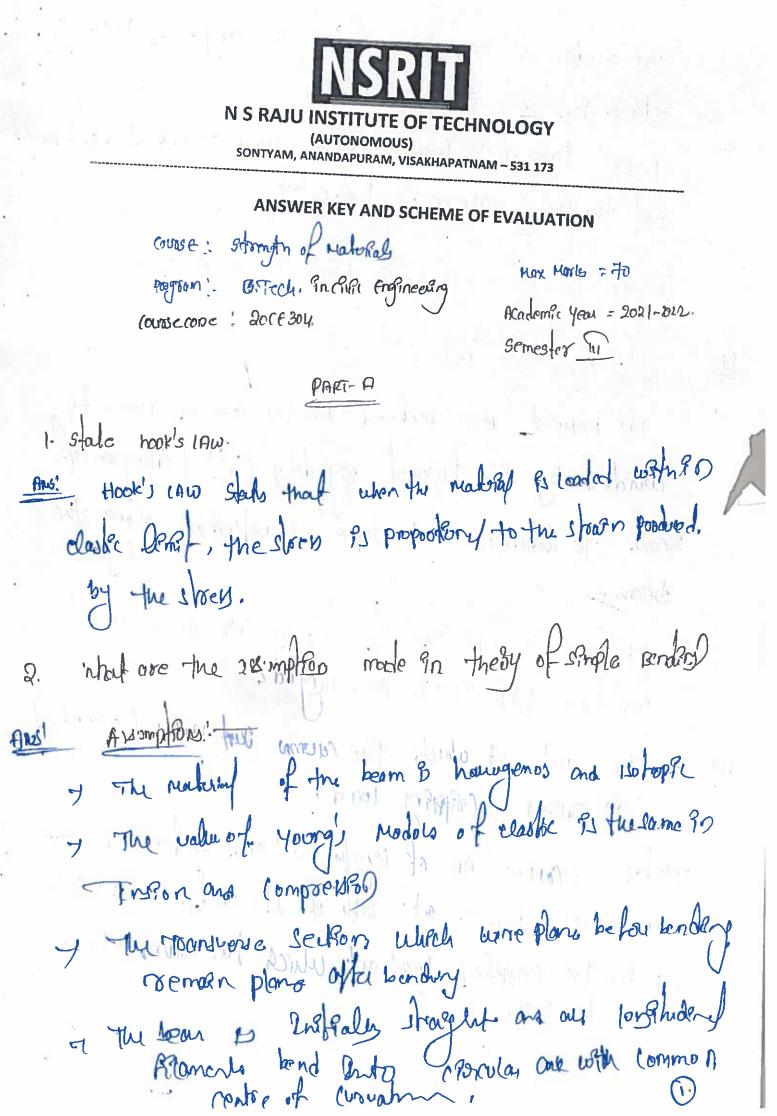
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20CE304.5

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N S RAJU INSTITUT TECHNOLOGY (AUTONOMOUS) SONTYAM, ANANDAPURAM, VISAKHAPATNAM – 531 173 ANSWER KEY AND SCHEME OF EVALUATION load when both ends of the column as haped (Sipplen The backellente. 1 P = P= MEI for effective length. he he effectives length as actual lengths be = l. betweens Mosed coll and open coll helper offerentiate Sporged !! Helfcal springs open coiled helican Closed- rolled helical spirit



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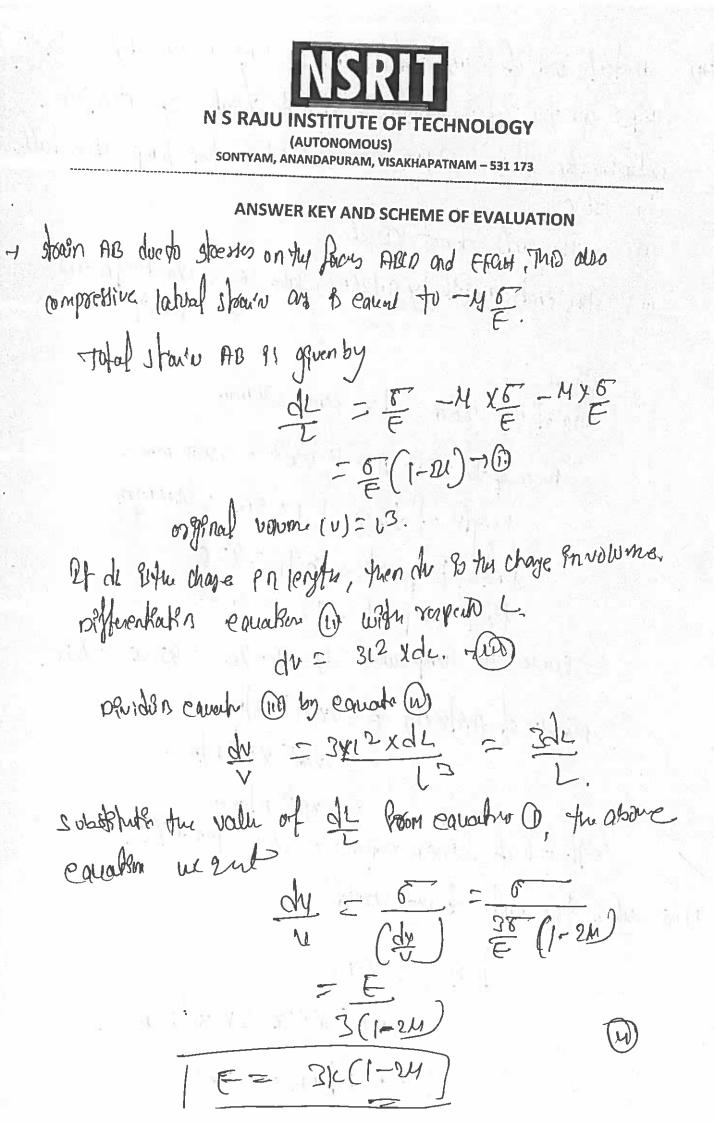
Open coiled helical springs closed colled helical springs y where of the open could " crosed corred hellcal sporry are the sporry hered sparg & wound not so In white helse angle fivery smart fequety onsthus sulfrand space The paper blue two addacent turory 21 small a goup eusin blu two adbaad on the nosed helfeal sporty correspond. -1 Piter of 11892 while as comparately azral load n The bendary expect Remarked. Larse a revuest larse helprogle m Qf. soblechato averas long and code . ~ 21 CONSENTS of posels TODSPORE/ STERIS and the stand of the second stands of the second seco NO VICE DESCRIPTION the second second second × 1 Marp : $\left\| \int_{X_{i}}^{X_{i}} \left\| \left\| f_{i} \right\|_{\infty} dx \int_{X_{i}}^{X_{i}} \left\| \left\| \left\| f_{i} \right\|_{\infty} dx \int_{X_{i}}^{X_{i}} dx \int_{X_$

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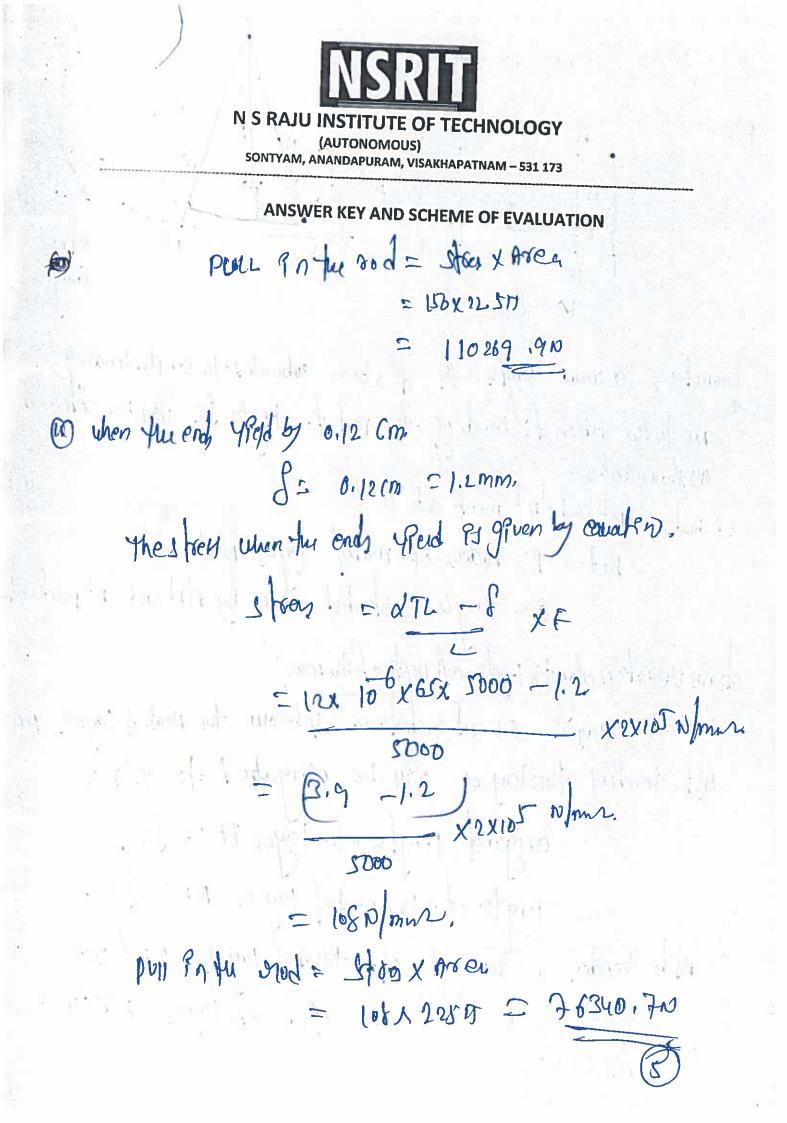
Start In Juliant

PART-B stearn curve Br mid steel rod. Explain Story -60) K-strain maderas 12 weeding -0 End Stock × fracture. ypedskeent Rie = PHZ = Slope Run 400 moduls Straga Dent of proportionality OA=

the relation between young's modules and modely of FINS ABCDFFGH Whitely & subjected to three mutually punder Promping, a cobe spects of earland Entensis, Vensile hut l= leigh of cobe dr = chape in jegger of the cube. yourgs module of the maker up of thake 5= Persia shress a Ray on the focs. M 2 Popson Jours. The volume of the cobe H= 13. ~ The spoorn (AB) of the side of the wabe three mortuals perpendicular The following three shares: -Jfred . Ostarn AB due to strend on the Bach AFHD. 1 @ strain AB dur to storms on the Bacs AEFB and DHAC The composensive Latial Joan Beaux to -ME.



try Aster and of 31m dramete and 5m long is connected to two goips one the rod & maintellind at a temprosture of 95 celetics petuntine the stress and prove encorted when the temporature falls to 38C. 1) the ends donot yead 1 the ends yld by 0.12(m.) Bale E= 2 X105 MIN / m2 and Gruen ! Any____ pla of the rod d= 3(m2 30mm Area of the rod, A= I xou2 = 22517 Mm2 Length of the god LESM = Soromm Robert tempersature 10T1 = 95C Final Remposalere T2 = 35-For4-9n Jemperature J=JI-J2 = 95-20 = 6FC. Model not fosterily F= 2x105 mm/m2. ~ 2×105 × 10 pm 2 2×106 p/m Coefficient of lenear apoising & = 12×159°C. a) a when the ends to not field !-Strey = diTIE = 12x10 6 x 65x 2x 10" N/m2. = 156 N mn2 (dentus





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Now bomby - win 1 = RXO and Elf' = (Rty) & (room of E'f'= Pty). But N'N = NN = fre Atena fu = RXO Inescon in the legth of the layer EF = EF - FF. = (Fty) O- PXO. FYXO. Stockin Pritu Layon EF I Invocane Rin the lergh organal legts $= \frac{y \times a}{EF} = \frac{y \times a}{E \times a} \cdot (: EF = fe = R \times a)$ = YR. n & constant . hence fuestion in layer & proportions to It's destones boory the neutral axes. the above caration Show the Variation of Straly alog the depthof the beam the parsiation of strang linens. B

stock vorstation! hut 6= street An tue layer EF E = yourg's moduls of the beam. from E= strengentus Loom EF Strain nin the Loga EF = C: stars 9nFF=MR (YR) TO = EXY = EXY Since Ears pare constant. therefore shes many lares fi durely propritioned to the distance of the law Commenter neutral lan us. if the constation of store is along the depth of the brown, the valiation of streng brows. 1 g=E Nation outs and moment of Restitutes the natival aris of any Toursverso section of abeau & defined a tur Dire of Intersection of the neutral lays with the towards sealson. It sworthing Dant -Fine 1979 Address public benefits and the provide a server in the set of the proof of the set and the second sec



ANSWER KEY AND SCHEME OF EVALUATION

If a section of a bears & sobored to purce sogging moment, then the store him will be compressive at an point above try neutral axis and tensive locking the neutral axis, those & no stress at the neutral axis. The stores at a destance y known the neutral axis & given by a the stores at a destance y known the neutral axis & given by a the stores at a destance y known the neutral axis & given by a the stores at a destance y known the neutral axis & given by a the stores at a destance y known the neutral axis & given by a the stores at a destance y known the neutral axis & given by a the stores at a destance y known the neutral axis & given by a the stores at a destance y known the neutral axis & given by a the stores beeken of abeans, but no he the neutral axis of the could aver a function. Consider a small layer of a destance y boom the neutral axis.

lonshall it small inger of the layer. het dA = Area of the layer. pour force on the wayer = stress on layor & Area of layer.

= 5 x dA = E x y x dA - JO (F=E XM). fore fore for the beam section & obtained patropath the alone caraba total fore on the beam sector D

E JYXdA=0 JyxdA =0 Moment of respirance pure perden, the layers above the NA are subjected to compressive shorrs whoreas the basers below this N. A corre subsected to tensile strend. of them force will have moment about the with. fills in layer = Exyx2A. Moment of the fole about Nite = Folce on lay u XY = EXYX dAxy = EXYZZA. Etter moment of the Bis orthe sections of the being = J ErynxdA = EjynxdA. ent M = Expiral Momente applied on the ben skelve. For candilitioning the moments of readours offered by Ju Sectus should be care to the cuton lendy mont M= EJYNXda)



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~ het M= Oxformal moment opplied on the bean section. Facaueldon -the moment of restation of forened by the Scotow should be thefonal berdy mount T. MO E MADD a suppression Jyr xdA sepresors the moment of shorthar of the arrea of the section about the neutral arra-Moment of Anosta be Dr but the M= FXIN M = F where or E Mart F M 95 eypressed 91 NMM; 290 mm/2 Of cuprested in Nmms; yinm EPS exproved & N/mmz, Rinmon

perfive shear strey distribution for concular section? 99)

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ANSWER KEY AND SCHEME OF EVALUATION

2 XEBX dy = &x JR2- y2 x dy [A ste FB = [12-y2) ~ moment of this areada about N.D 5 YXdA = 4x 2 J R2- y2 x dy = ey J R2-y2 dy. ~ moment of the whole shaded trea about N.A & obtained by Integration the above cauality to purps limits. Your R. AY = Jey VR2-y2 dy = - J (-2y) J p2-y 2 dy Now (-y) stru differential of (R2-y2). Integration of the obour construction becomes 3/2 Ay = (P2- y2)/2

$$= -\frac{2}{3} \left[(e^{2} - e^{2})^{3/2} - (e^{2} - y^{2})^{3/2} \right]$$

$$= -\frac{2}{3} \left[0 - (e^{2} - y^{2})^{3/2} \right] = \frac{2}{3} \left[e^{2} - y^{2} \right]^{3/2} \right]$$

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$$= \frac{1}{10} \left[e^{2} - y^{2} - y^{2} - y^{2} \right]^{3/2}$$

N S RAJU INSTITUTE OF TECHNOLOGY (AUTONOMOUS) SONTYAM, ANANDAPURAM, VISAKHAPATNAM - 531 173 ANSWER KEY AND SCHEME OF EVALUATION portue shear story distribution for rectorquian section? 967 BUS! Rectanguar section of beam of wetter b and depter a . Int FPS the stearface acting at the section a level EF at a distance y boorn the national ares. Considers - The thear stress at this level given by cauch. J 2= FX AY Rb where Az Area of the section above CHoded Bree ADEC) de-y - - A grax 5 = obtain of the C.C. of are Aloon neutral and 월 + 是 (d/e-y) = y+ = 프 = 블 + 북 월 (3+월) Actual whether of the seatson at the level EF I HOR of the whore section about N.A. Sapphin the value 9 retur abone caraho 62 Ex[= -2]xpxf[24 bxQ

$$= \frac{F}{3\Sigma} \left(\frac{dr}{dt} - b^{2} \right)$$

ATTOPEDER , ys 4/2 and how

$$P = \frac{F}{2\Sigma} \left[\frac{dr}{4} - \left(\frac{d}{2}\right)^{2} \right] = \frac{F}{3\Sigma} \times b \ge 0$$
ATT the ngubbed arts, y=0 hence.

$$P = \frac{F}{2\Sigma} \left[\frac{d2}{4} - 0 \right] = \frac{F}{2\Sigma} \times \frac{dn}{4}$$

$$= \frac{F^{2}}{8\Sigma} = \frac{Fd}{3\times b20}$$

$$= \frac{F}{5\times b1} = \frac{F}{5\times b2}$$
Now Avery Meadling The - Strength = $\frac{F}{5\times d}$.
Now Avery Meadling The - Strength = $\frac{F}{5\times d}$.

$$\frac{F^{2} = \Lambda T \times F_{5}}{T^{5}} = \frac{1}{7} \times \frac{f}{5d}$$

$$= \frac{F^{2}}{7^{5}} = \frac{1}{7} \times \frac{f}{5d}$$

$$\frac{F^{2} = \Lambda T \times F_{5d}}{T^{5}} = \frac{1}{7} \times \frac{f}{5d}$$

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N S RAJU INSTITUTE OF TECHNOLOGY (AUTONOMOUS) SONTYAM, ANANDAPURAM, VISAKHAPATNAM – 531 173

NSRIT

ANSWER KEY AND SCHEME OF EVALUATION

109 Oxplain Conjugate beam method?

mon the slopes and defection of bens can obtains by con origat been mend Consucrase Alton methods Before deversibling the conducate been method, het first dokne Consucade beau. a consucade is a smoginory beam of contra caual to the that of the original bean but to which the load drogsom is the M diagram. (the load at as point on the consticute been & also court to fim at the point drive d' by DE The rops and deplections at any sations of beams by conducted beans method is grue, The slope at my section of the given beau & caual to the thous force at the Concepondry Section of the conductate Beard O The deflection at an ichon for the gran beau. permito benden momente at the coorespondy of the consultate been ~ Befor Approxim the contraction bean method, (on Richarts bean D) constructed, the load on the conducate beam at any portet B Caund to the R.N. at that point divided by El Hence be loady on the conscients bean & lowowas

The sheatore at any point on the considerate bea gives the yope at the corresponding point of acture beau. the Breaton Port gate bean gen the Jope at the corresponding n The Rimat on point on the conjugate bear firs the deflection at the consemporating point of the adual tean. point of acheal beau. MER dingrow is diagrow which thous the vortation of MER over the longth of fur beau. 106) porrive on empression for the slope and perkerties of simply Sopported beau with point load at contre? An! A strupy supported bean AB off Lergetty i correging point load of w at the centre c. The B.M at A and BB Zoro and at the Centre B. M Wilbe Wyy. - The B. M vorsing according to shought line law. ROEWL WL L RB= W/2 WYy Bindiagra to a a a a d Load dioph THUR WE On 54 Cabepes et 5 the



ANSWER KEY AND SCHEME OF EVALUATION The Bim varies a cooley tostraght the law, the Bim dragman. > floor Inhy. a Nowthe constrate Rean AB canbe contructed the load on the Conducar Rean will be ablacted by devider the Bin af that Pont i the those of loady on the Conducate Bear will be some of Bim whe orderal of loady on conducto near will the events $\frac{M}{6\Sigma} = \left(\frac{WL}{4P}\right) = \frac{WXL}{4EP}$ House detaute at the Centre WILLE WL du Mounts. LA= Readin at A RA On Sucate Dan of fondy we LO = Realty at & For Ion Decate Dean -potal load onthe Conducate bear. = Area of the coor dagn. = YL XAOX (D = XIXWL = WL2 - Reaction at each support for the constructure half of the total lound,

RB = t x WIL FEE = WLZ 16ED 2. Man (ds/An)adA. hub DA 2510pe at A for the plue beau Hu Decoder to conducate bean method. OA 2 shear for at A fortur conducate bear = RA (S.F.at For conducute sean = RAS) = WLZ 1682 as yer simate forthe consulate beau. = RA XL - Load arrispondy to ACP. xoBrance of GG of ACD Brance $= \frac{W(2)}{16EL} \times \frac{L}{2} - \left(\frac{1}{2} \times \frac{L}{2} \times \frac{WL}{UEL}\right) \times \left(\frac{1}{2} \times \frac{WL}{2}\right)$ $2 \frac{WL^3}{3LER} - \frac{WL^3}{966R} = 3WL^2 - WL^3$ Jys wer uter And the set of a plant of ATT Consideration of n en en anti-se de série des

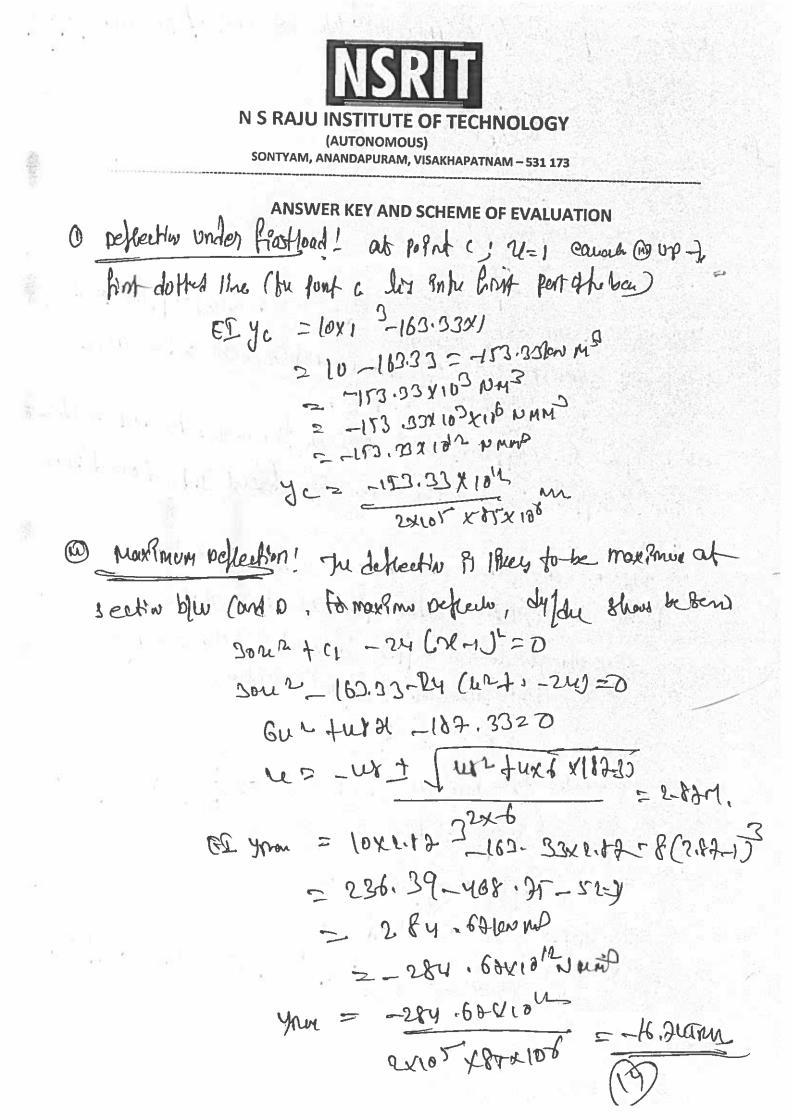


ANSWER KEY AND SCHEME OF EVALUATION 1) A been of lergety 6m is simply supposed at 21 and and avoiles two point loads of usion and wolow at a distance of Im and 3m repectively from the left sopport? E= 1×105 N/mm 2, I285× 108, MA (Deflection under cach load. Arsuned. most mon peffective. QUUN' ATTA I= 85×105 MMH ! E= 2×105 N MM Fint calcuat the reactor. RA and FB. Taled moment about A, we get RAX6 = U8X1 + U0X3 = 168 RB = 168/6 = 2800. RO = Dofeel roud - RB = (ustue)-28 = 60/002 Holand lut 6-1m -_____371 Consider the Section x inter last post of the bean (rem legger DB) at a distance se from the left supporte A. The B.M. about section is gen by El dry 2 Rore & -us(u-1). dur 2 Rore & -us(u-1). -40(21-3)

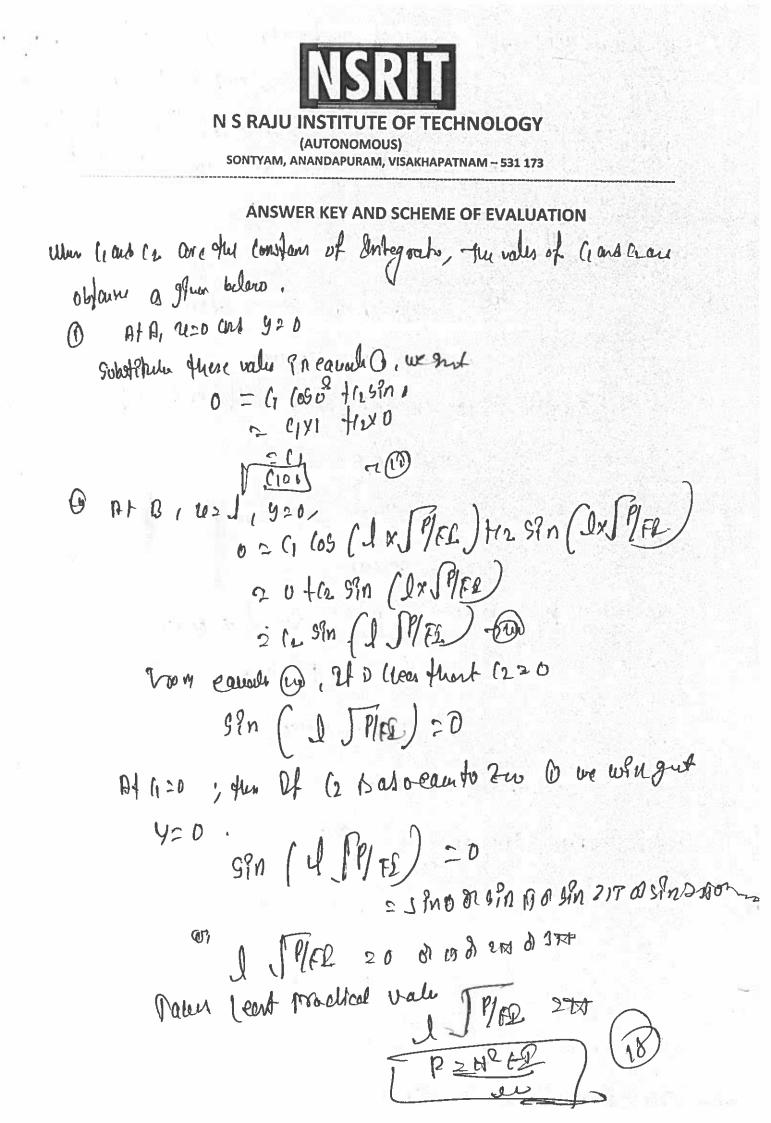
$$\frac{2609}{44} = -44 (24-1) = -40 (24-2)$$

$$\frac{1}{2} - 44 (24-1) = -40 (24-2)$$

$$\frac{1}{2} = \frac{1}{2} - \frac{1}{2} + \frac{1}{2} +$$



Derive an expression for crilling load when both ends of the coveneral 12/ hirged! Aro' The load at where the colum Dirt buckers (a bonds) & called the ppth band ~ CONSIGN a Colomo AND of 100445 I and Uniterry (NOD) saltime are hinged at Both of Efsend. Aga B. - but pre the (Rpp Brog load at about the column by Det builted ~ Doe to cripply load , the colors My deland anto & oned from consiste any seems at a distance reboon the end A. ACB but y = Deflectro (caloreds placent) at the skelvs pur normet due to rripply bout est the kelo = - py (-v & sign tale due to sign Canvention) M = FL dry Gauston two nonnes we have the dur FD 14 + PY= D 14 + 1 . y = D The yoluth of the above duffen carech yzy cos (21 JP/FR) tcl-Sin (22 JP/FL)





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ANSWER KEY AND SCHEME OF EVALUATION

0 0 000 - when a chirular shalf & subsceled to toxin, shear stress are schop in the making

of the shoult. - To determine the magnetise of shear strong at any point on the safe corners a kielt Bred at one end AA and Bree at the end BB as blown 9n by. - hut co is any line on the outer stops of the shelt.

" The short & us Seek & to forgue of at the end BB of thous they As a result of the toraws 7, the Halt at the end BB why relate areling an ever mosseching of tu sharbt Will be subtrand to shows true. outrupoint puting this to pland have time cours be deleas to co

of shown why.

posive tossion canafier

Ly)

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it the fine op while should to op'

Mt R= ladies of short L= confr of shelt T = Torque applies at the end BB Y= Shearstress Indiced at the sublan of the dett doe to -10094 nodors of stoppedan of the marker of the Stall, (=



1.10.10.1.1

(PD. YYL p.9 $\frac{co}{l} = \frac{c}{R}$ Y = PXCXO ~ now, given shaft subsect & to a given torane (m), the values Li Dians Lance constent. sheastes & produces to the rodes R. Par of Front If q githe theastring induced at a redusing "Poury the centre of the Keaht thy $\frac{N}{R} = \frac{\gamma}{R}$ $\frac{\gamma}{P} = \frac{CQ}{L}$ brin en of sheastern & maxim at outerstar an sheaster B 200 At the arts atulally.

15) The stiffney of a closely coiled helter sporg is 1.5 N/1007. of composition under a maximum load of 100N. The maximum, Shearry stress produced an the come of the sport of SUDDN/mm The bild length of the spring (when the coils are taveling) is give And O planch of the whole. I mean drame he of the coils. 1) No of Colls reasoned. Pale C=45 XIdeN/MM2 stillnes of spoint s= 15 N/m. AN . load onspory = w= 100 D. maximon sheash 7 = 130 mm2. sold length of spring = SIM = 50mm. Hodulu of ofgedory (= 4.5×104 N/mm. het d = dametre of wire. p = mean dra of coil R= mean radius of Carl 0/2=R n2 nomber of costs. S = (dy =(0) 1.5 = 4.5×10444 GURS. N GYXR3XN.



ANSWER KEY AND SCHEME OF EVALUATION d4= 1.5 x64 x R3XD 4.5×104 d " = 0,00 2133 R3xn -1(7) 9 = 10WXR 87 180-2 16 XLOOKR R = 130 Kird) -10 ×100 R = 0.4090 de -10 solution for value of Rin eauert O, we and dy = 0.00 2133 x (0.4 0906d3) 3 n = 0.002/33 × (0.4 0906)x29 xn = 0.000 14599 x divn 19.0 = 1 34 0.000 / 15 n= 1 0,000/45992 Jolid length = nxd di so=nxd $n = \frac{sp}{sp}$ Sobtinue top value of n cavallin

ds x 50 = _____ du- 1 x = 136.95 0.000/4199 To = 136.95 n= 10 = 50 = 14.025415 R= 0.40906d= = 0.40906 ×13.42) R=+6.36m mean dua of coil D=2R = 2x16.00 D=32.72m Thirty A at. IS & Chames weed & Call 32. A 2 P Were and mark Long Long I for a sheet I and 1. 1. 1. 14 for direct of product rate dealers

Nedropall Sevene avere (Republished) Technology (Automotions) 1046 (Ruelly Materiani System (DIMS)

Degree	B. Tech. (U. G.) Pro	nd Supplementar ogram Mi				demic Year	2021 -	2022
Course		st Duration 3	Hrs. Max. N	larks 70	Serr	nester	II	l
Course	Mechanics of Solids							
Part A (No.	hort Answer Questions 5 x 2 = Questions (1 through 5)	10 Marks)					(-)	D -1
1	Give the relationship between bu Poisson's ratio.	ulk modulus and mo	odulus of elasi	icity in term:	s of	Leaming Outco 20ME303	• •	Doł L1
2	Differentiate the point load, UDL	and VDL.				20ME303	.2	L1
3	Define the flexural rigidity and to	rsional rigidity.				20ME303		- L2
4 5	Mention the advantages of hollow Define buckling and stability.	w circular shafts ove	er solid circula	r shafts.		20ME303		Ľ
	ong Answer Questions 5 x 12 =	= 60 Marks)				20ME303	.1	Ľ
No.	Questions (6 through 15)	- 99 marx9)		Ма	rks	Learning Outco	mo /s\	Dol
6 (a)	Draw a neat stress- strain curve	diagram of stainles	s steel and ex	plain 6	M	20ME303		L
	A steel bar is 900 mm long ,its (two ends are 40 mi	m and 30 mm	in dia				
6 (b)	and the length of each rod is 20	00 mm The middle	portion of the	bar is	M	20ME303	2	
- (-/	15 mm in dia and 500 mm lo tensile load of 15 KN .Find total (ng ,if the bar is si extension Take E=	ubjected to an	ı axial	IAI S	201012303	.2	L
			200 GFa					
	A steel bar of 20 mm diameter is	s acted upon by the	forces. What	is the				
	elongation of the bar when you	ng's modulus, E =	210 GPa. Fi	nd net				
	elongation by principal of super A B B	CI D	8					
7 (a)				6	М	20ME303	.3	Ľ
	60KN 20kN -	30kN	50kN					_
	3m 1n	n 2m	1					
		11	1					80 80
	The stresses on two perpendicul	lar planes through a	a point in a bo	dy are				
7 (b)	120 MPa (Tensile) and 80 MPa and tangential stress on a pla	(Compression), De	etermine the r	ormal 6	М	20ME303	.3	L
	(ACW). Draw configuration and I	Mohr's diaorams	o with the v	enicai				_
8 (5)		-						
8 (a)	Write about the types of beams in Draw shear force and banding of	n detail with neat si	ketches		М	20ME303	.2	L
	Draw shear force and bending m and find shear force and bending	ioment diagrams to	r a cantilever l	beam				
	loads 3 kN and 6 kN at right end	and 0.5 m from right	carrying two p	oint				
0 /63	ioudo o kia and o kia at light end	6 kN 3 ki						
8 (b)	Α	C[0.5 m] B	•	6	M	20ME303	.2	Ľ
			,					
	L . 3 m							
		>						
	Differentiate the observations		DR					
9 (a)	Differentiate the shear force uniformly distributed and variably	and bending more	ments when	point, 6	М	20ME303	.2	L
	A cantilever beam 4 m long can	ries a VD. 2 kN/m	at the free en	d to 5				
	kN/m at the fixed end and draw \$	SFD and BMD						
				-				
9 (b)	6 kN A	J J B	3 kN	6	М	20ME303	2	L
	4 <u>4</u> <u>m</u>	>					÷.	
10 (a)	Write the sign convention of ch	our force and here	din =t	la ilea - O		00115000		
	Write the sign convention of sh				M	20ME303		Ľ
信行力	20 I rugadan Bara selijak se	nia dal ministra	in item	File and	1-1	a Kalin		
5 15.1	the site water water and the		erest for a second s	- Co	ntro	ler of Ev	mino	1

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10 (b)	left of the section. Draw the free body diagram, shear force and bending moment diagram of VDL on simply supported beam OR	6M	20ME303.3	L2
11 (a)	The simply supported beam is 6 m carries a point load 4 kN at a distance of 2.5 m from left where the UDL of 3 kN/m starts from point load for 1 m. Find the reactions and draw the SFD and BMD $\begin{array}{c} 4 \text{ kN} & 3 \text{ kN/m} \\ A & C & D \\ R_A & 2.5 \text{ m} \\ \end{array}$	6M	20ME303.3	L4 、
11 (b).	Derive the equations for simply supported beam with UDL an UVL	6M	20ME303.2	L2
• 13	Derive equation for moment of inertia for a I- section	5M	* 20ME303-2	L2
12'(a) ়ু 12 (b)	A soli circular shaft transmits a power of 60 KW at 200 rpm. Find the diameter of the shaft if the allowable shear stress is 50 MPa and allowable twist is 20 for every 10 mt length of shaft. Take C=80 GPa.	7M	20ME303.2	L2
	OR A cantilever beam of span L is subjected to a UDL of W KN/m at a distance 'a' from fixed end. Find the deflection of free end			
13 (a)	A BL C	5M	20ME303.2	L2
13 (b)	Explain the Macaulay's method in deflection of beams	7M	20ME303.3	L2
14 (a	Explain the buckling failure in columns with Rankine formula	6M	20ME303.2	L3
14 (b	Derive an expression for Euler's critical load with one lixed an the	6M	20ME303.2	L2
15 (a	Distinguish between circumferential and longitudinal stress in cylindrical shell when subjected to an internal pressure.	5M	20ME303.2	L2
15 (b	Differentiate thin and thick cylinders and write three applications of	7M	20ME303.2	L2
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	Contro	mer o	f Bxaminatione	

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Controller of Examinations NSRIT (A) Visakhapatnam

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Nadimpal I Satyanarayana Raju Institute of Technology (Autonomous): IQAC: Quality Management System (QMS)

NSRIT

Semester End Supplementary Examination, April/May, 2022

Degree	B. Tech. (U. G.)	Program	ME		:	Academic Year	2021 - 2022
Course Code	20ME303	Test Duration	3 Hrs.	Max. Marks	70	Semester	
Course	MECHANICS (OF SOLIDS		TAICH V.S.			

No. Questions (1 through 5)

3

Define the bulk modulus and modulus of elesticity.

Ap 15 (0. 202), Quantum Paparlantine and Semister Examination 1 Apadem CRepublicition 2020

- To put in more simple words, the bulk modulus is nothing but a numerical constant that is used to measure and describe the elastic properties of a solid or fluid when pressure is applied on all the surfaces. The bulk modulus of elasticity is one of the measures of the mechanical properties of solids Differentiate the point load, UDL and VDL.
- When placed in steel storage racks, a uniformly distributed load is one whose weight is evenly distributed over the entire surface of the rack's beams or deck. A point load is a one with its weight significantly concentrated in one (or more) places on the rack's beams or decks.

M/I=f/y=E/R – justify. Bending Equation. The axial deformation of the beam due to external load that is applied perpendicularly to a longitudinal axis is called the Bending Theory. The bending equation stands as f/y = E/R = M/I.

$$\frac{q_{\text{GAM}} \text{ fanding } [= q_{\text{GAM}} \text{ fan$$

Explain pure torsion

4

5

When the circular shaft is subjected to torque only without being acted upon by any bending moment or axial force, the circular shaft is said to be in the state of pure torsion. Define buckling and stabilitythe state of being stable

 $B = -\frac{M^2}{M} \times V$

Wheter

8 = Bulk modulus (psi) AP - Charge in pressure (sai) AV := Charge in volume

V = Intel volume

When the value of B is known (see reference table on next page), it is easy to calculate the effect of any pressure change en volume, or of any volume change on pressure.

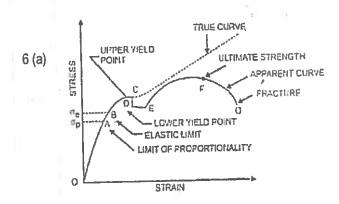
$$\sqrt{V} = \frac{V}{B} \times \sqrt{P}$$
 or $\sqrt{P} = \frac{V}{V} \times \sqrt{V}$

No.

6 (b)

2

Questions (6 through 15) Draw a neat stress- strain curve diagram of stainless steel and explain.



A steel bar of 900 mm long, its ends are 40mm and 30mm in dia. and the length of each rod is 200mm. The middle portion of the bar is 15mm in dia. If the cylinder is carrying a load of 20kN, find the total extension when modulus of elasticity is 200 Gpa.

Area of steel bar =2800mm2

Marks

6M

Data: Area of steel, As = 16 cm2Area of brass, Ab = 10 cm2Weight, W = 500 kg

 ΔI is same for steel and brass, so

2m

As=yl δl

The sum of forces of steel and brass is, fs+2fb=5000 ------(1)

Ss=(fs)/(as), sb=(fb)/(ab)

Fs=ss.as, fb=sb.ab

Ss=(ys)(is)(δl)

Ss is directly proportional to isys

Similarly, sb is directly proportional to ibyb

Where is and ib are the moment of inertia of steel and brass 2m

s is the stress

Sbss=ybisyslb=2×32=34

Ss=34sb

Ssas+2(ssab)=5000 ----- from eq. (1)

(ss)(as+2×43)(ss.ab)=5000

Ss (16+23×10)=5000

Ss(31)=5000

Ss=315000=161kg/cm2

Sb=43×115000

Sb=121kg/cm2

5

A steel bar of 20 mm diameter is acted upon by the forces. What is the elongation of the bar when young's modulus, E = 210 GPa. Find net elongation by principal of super position.

2M

Sol: Length, I = 5000 mm; cross sectional area, A= 22/7x20x20/4 =400mm2;

7 (a) Tensile force, P =50 KN = 50 x 103 N; Young's Modulus, E = 200 x 106 N/mm2.

6M

As per the young's modulus, $E = [(P/A) / (\delta I / I)]$

Therefore, change in length or elongation,

s,

 $\delta I = (50 \times 103 \times 103) / (400 \times 200 \times 103)$

 $= \frac{1}{2} = 0.6 \text{ mm}$

2M

7A. Area of steel bar =2800mm2

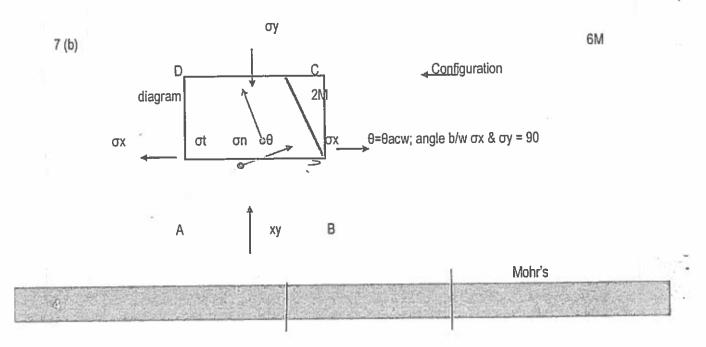
Data: Area of steel, As = 16 cm2 Area of brass, Ab = 10 cm2 Weight, W = 500 kg

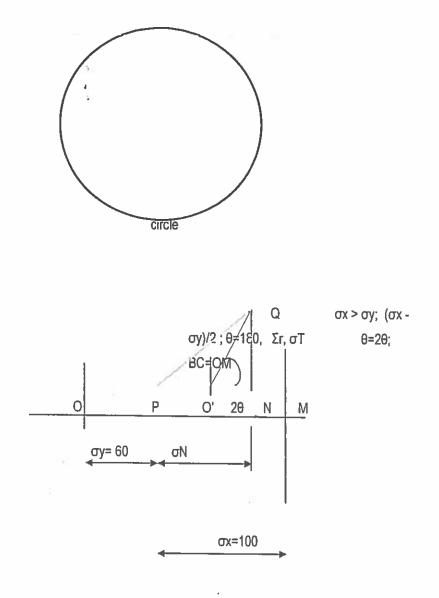
△I is same for steel and brass, so A B C D 60KN 20kN → 30kN 50kN 3m 1m 2m

The stresses on two perpendicular planes through a point in a body are 120 MPa(Tensile) and 80 MPa(Compression). Determine the normal and tangential stress on a plane at an angle 300 with the vertical (ACW).Draw configuration and Mohr's diagrams.

The stresses on two perpendicular planes through a point in a body are 90 MPa

(Tensile) and 70 MPa(Compression). Determine the normal and tangential stress on a plane at an angle 250 with the vertical (ACW).





 $\sigma N = ?; \sigma T = ?; \sigma R = \sigma R = \sqrt{\sigma N2 + \sigma T2} = 4M$

6**B**.

5

A steel rod of 5 m long and 20mm diameter is subjected to load 50 kN. Find the elongation when young's modulus is 2×106 N/mm2. Poison's ratio is 0.25

Sol: Length, I = 5000 mm; cross sectional area, A= 22/7x20x20/4 =400mm2; Tensile force, P =50 KN = 50 x 103 N; Young's Modulus, E = 200 x 106 N/mm2.

2M

As per the young's modulus, E = [(P/A) / (δ I / I)]



2M

$$\delta I = (50 \times 103 \times 103)/(400 \times 200 \times 103)$$

 $= \frac{1}{2} = 0.6 \text{ mm}$ Write about the types of beams. Simply supported beam Both end hinged beam Fixed beam 6M 8 (a) Continuous beam Cantilever beam 4 1 Propped cantilever beam Draw shear force and bending moment diagrams for a cantilever beam and find shear force and bending moments of span carrying two point loads 2 kN and 5kN at right end and 0.5m from right end. 3kN 6kN 6M .0.5m ₿ 8 (b) ΤC А 3m Differentiate the shear fo 6M

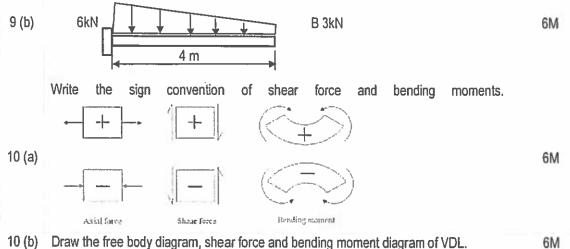
rce and bending moments when point, uniformly distributed and variably distributed

9 (a)

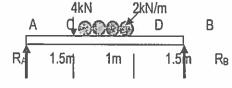
6

Load	Slope for shear force	Slope for bending Moment
Р	Constant	Linear
¥		\bigtriangleup
Uniformly distributed load	Linear	Parabolic
Uniformly varying load	Parabolic	Cubic
		\bigtriangleup

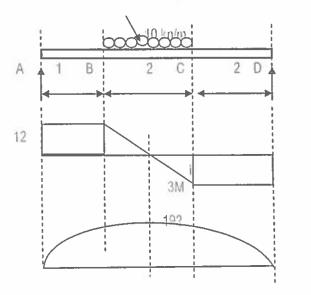
loads applied. Figure-1 Slopes for various types of loads A cantilever beam 4m long carries a VD, 2kN/m at the free end to 5kN/m at the fixed end and draw SFD and BMD.



10 (b) Draw the free body diagram, shear force and bending moment diagram of VDL. The simply supported beam is 5 m carries a point load 4 kN at a distance of 1.5m from left where the UDL of 2kN/m strarts from point load for 1 m. Find the reactions and draw the SFD and BMD.



11 (a)



Ma =0

Mc =12x 1=12 kn-m

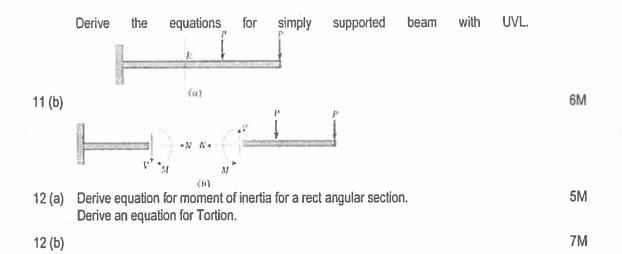
Md=0 =8x 2=16 kn-m

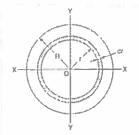
x/12 =(2-x)/8

x= 1.2 kn-m

.

Mm = 12(1+1.2)-10x1.2xq.2/2 = 19.2 kn-m 3M





Total torque could be easily determined by integrating the above equation between

limits 0 and R.

Polar moment of inertia is -----2

J=∫r2 dA

Therefore total torque transmitted by a circular solid shaft is,

T=tJ/R

T/J= t/R 4M

The expression for shear stress produced in a circular shaft subjected to torsion and therefore we have following result from that expression

 $T/J = C\theta/L$

Considering above two equations, the torsion equation for circular shaft is, $T/J = t/R = C\theta/L - ---3m$ 4M

Where,

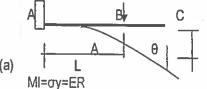
C = Modulus of rigidity

L = Length of shaft

 θ = Angle of twist

A cantilever beam of span L is subjected to a concentrated load W at a distance 'a' from fixed end. Find the deflection of free end.

y2



13 (a)

3

5M

Where, E = Modulus of Elasticity of beam material

R = Radius of Curvature of the beam

M = Moment of resistance

i = Moment of Inertia of beam c/s about N.A.

 $\sigma\sigma$ = Bending stress at the layer situated at a distance y from N.

3M

y = Distance of layer from N.A.

σ/y = E/R -----1

At the distance 'y', let us consider an elementary strip of very small thickness dy. ' σ ' is the bending stress in this strip.

o la tre bendrig atesa in tria sulp.

Let dA = area of this elementary strip.

The force developed in this strip = $\sigma_{s} dA$,

The elementary moment of resistance due to this elementary force is given by dM = f.dA.y 3M

Total moment of resistance due to all such elementary forces is given by

 $\int dM = \int \sigma x \, dA x y \text{ or } M = \int \sigma \times dA x y ----- 2$

From Eq. 1, $\sigma = y x (E/R)$ Putting this value of f in Eq. 2

we get where I = Moment of inertia of the whole area about the neutral axis N-A.

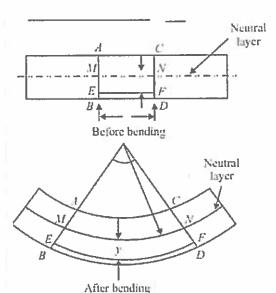
Where; M = Bending moment I = Moment of Inertia about the axis of bending i.e; Ixx y = Distance of the layer

The extreme distance of extreme fibre from N.A.)

E = Modulus of elasticity of the beam material.

R = Radius of curvature



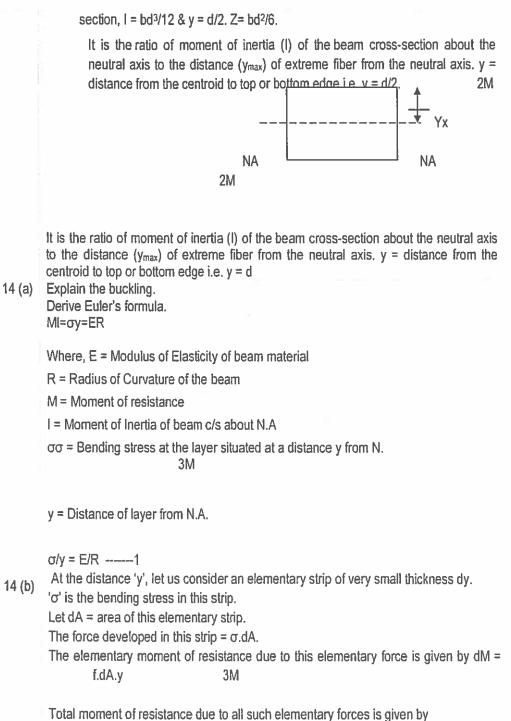


Explain the Macaulay's method in deflection of beams.

13 (b)

13

The modulus of section may be defined as the ratio of moment of inertia to the distance to the extreme fibre. It is denoted by Z. Z= I/y ; For rectangular



 $\int dM = \int \sigma x dA x y \text{ or } M = \int \sigma x dA x y ----- 2$

From Eq. 1, $\sigma = y \times (E/R)$ Putting this value of f in Eq. 2

we get where I = Moment of inertia of the whole area about the neutral axis N-A.

Where; M = Bending moment I = Moment of Inertia about the axis of bending i.e; Ixx y = Distance of the layer

The extreme distance of extreme fibre from N.A.)

E = Modulus of elasticity of the beam material.

R = Radius of curvature

14 (b)

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15 (a) What is circumferential and longitudinal stresses Define circumferential stresses, fil 5M

6M

The hoop stress is the force over area exerted circumferentially (perpendicular to the axis and the radius of the object) in both directions on every particle in the cylinder wall.

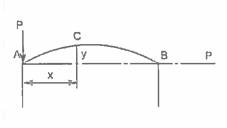
3M

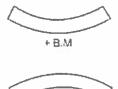
Define longitudinal stresses, f2

Longitudinal stress is defined as the stress produced when a pipe is subjected to internal pressure.

Longitudinal stress is also known as axial stress. . Differentiate thin and thick cylinders.

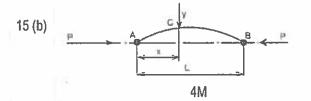
4M







According to sign convention



Consider an axially loaded strut, shown below, and is subjected to an axial load 'P' this load 'P' produces a deflection 'y' at a distance 'x' from one end.

Assume that the ends are either pin jointed or rounded so that there is no moment at either end.

Bsin {L \sqrt{p}/EI } =0 B=0 or sin {L \sqrt{p}/EI } =0 sin {L \sqrt{p}/EI } =0 or {L \sqrt{p}/EI } = π or nL= π 4M or { \sqrt{p}/EI } = π/L or P= $\pi^{2}EI/L^{2}$

12

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_	Semes	ter End Suppleme	ntary Examination, A	pril/May, 2	2022	
Degree	B. Tech. (U. G.)	Program	EEE	_		1 - 2022
Course Coo		Test Duration	3 Hrs. Max. Marks		emester	
Course	DC MACHINES &	TRANSFORMERS			5111C3(E)	u
Part A (Sho	rt Answer Questions 5	x 2 = 10 Marks)	Minimulation — 14 22:			27
No. Qu	estions (1 through 5)	i i i i i i i i i i i i i i i i i i i			Level of t	247
1 De	scribe multiply excited m	agnetic field system	h		Learning Outcome	
2 WI	hat is armature reaction	? What are the eff	ects of armature reacti	on on the	20EE304.1	L1
- Po	Hormanice of de machille	if in the second s			20EE304.2	L1
3 Wi	iy a starter is required to rter?	start a DC motor?	What is the essential ele	ement of a	2055204.2	~
316	nat is the function of a tra				20EE304.3	L2
	fine Sumpner's test.	instormer?			20EE304.4	II L1
	g Answer Questions 5	49 = 60 84-44-5			20EE304.5	L2
No. Qu	estions (6 through 15)	x 12 = 60 Warks)				
De De	rive the expression for	field energy and	co crorry in a daubi	Marks	Learning Outcome (s) Dok
6 (a) be	cited system assuming c	onstant current system	co energy in a doubl	у _{6М}	20EE304.1	L2
6 (b) De	rive the expression for to	rque in a sinoly exc	ited system	6M	2055204 4	
			OR		20EE304.1	L2
7 (a) Fo	a singly excited system	m derive the expre	ssion for magnetic field	. h		
CII	ay suied.			6M	20EE304.1	L2
7 (b) Wr	ite in brief about multiple	-excited magnetic fi	eld system.	6M	20EE304.1	L2
8 (a) Ev	dein the second of the tax					10.0
8 (a) Ex	plain the constructional d	etails of DC general	lor.	7M	20EE304.2	- L2
8 (b) 45	culate the emf generate	d by a 4 pole wave	wound armature having	9	×	
	slots with 18 conductors pole is 0.016Wb.	per slot. when dry	en at 1200 rpm the flu:	× 5M	20EE304.1	L3
*			OR			
9 (a) Exp	lain magnetization chan	acteristics of a DC s	bunt generator?	^а 6М	2055201.0	
9 (b) Exp	oress the EMF equation	of a DC generator.	nam generator ;	3 6M	20EE304.2	L2
				OIN	20EE304.2	L2
10 (a) Dei	ive the torque equation (of a DC motor.		6M	20EE304.3	L2
Ac	ynamo has a rated arm	lature current at 25	0A. What is the curren	+ 53	2022.004.0	12
in (n) hei	paul of the armature	If the armalure w	inding is lap or wave	€ 6M	20EE304.3	L3
COL	nected? The machine ha	as 12 poles.				
11 (a) Exp	lain speed control of DC	abu ta t	OR			
	lain in detail the losses i	shunt motor.	*	6M	20EE304.3	L2
(-)	idan in detail die 105565 1	n DC motor.		6M 📄	20EE304.3	L2
12 (a) Der	ive the EMF equation of	transformer				
A 4	00/230V, 50Hz, single	ohase transformat	had 200 turns on high	6M	20EE304.4	Ł2
12 (b) volt	age side. Find turns ratio), transformation, rat	io and number of turns	6M	0055004.4	
on	ow voltage winding. Also	find the flux develo	ped in the core	D CIVI	20EE304.4	L3
			08			
13 (a) Dra	w the various types of	three phase transf	ormer connections and			
DIIG	reach one of them.			0.01	20EE304.4	L3
13 (b) Ura	w and explain the phase	or diagram of single	e phase transformer on	cu		
1020	considering with windin	g resistance.		6M	20EE304.4	L2
D	e ile standt de se					
4 (a) Dra	w the circuit diagram of	Sumpner's test and	derive the equation for	6M	0055004 #	
The	iency of each transforme) [.			20EE304.5	L2
and	primary and secondary 600V respectively C	voltages of an auto	transformer are 1200V			
4 (b) and seco	600V respectively. Condary current is 120A	Draw the store	my of Cu when the	6M	20EE304.5	10
3500	indary current is 120A	Draw the circuit	and show the current	UNI	2014.0	L2

distribution in the winding.

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	A 10 kVA, 500/250 V, 50 Hz single-phase transformer gave the following test		6	
	data: OC Test (LV side): 250 V, 1.0 A, 80 W			
15 (a)	SC Test (HV side): 25 V, 12 A, 100 W Where LV refers to the low voltage and HV refers to high voltage side.	7M	20EE304.5	L3
	Determine the following:			
	 (i) Equivalent circuit referred to LV side (ii) Secondary load voltage at 0.8 p.f. lagging with full-load current. 			
15 (b)	Explain how the Scott connection can be used to obtain two phase supply from a three phase supply.	5M	20EE304.5	L2

C2. 12 8 5 13

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ALC: NO

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6a) Derive an expression for field and co energy in a doubly excited system assuming constant current system? Muttigly - excited Magnafte Aleba Bygloban		VI - No of C. Rutar	a prescribued angle. 3 Electro rechamical Transduces, have the specing 12 Syntements of Producing an electrical signal programments for conservery an electrical signal	Frequentized to electrical signal (current or veloage) > such Transduces require two excitation one ercitedion establishes a magnetic field while the	Hathomical) -> figure shows a raghetic field bystom with Two ellectoned erichabion - one on shabit and the other on robon	Ans) , Seeds of Three Independent visited in elblin of the strue
N S RAJU INSTITUTE OF TECHNOLOGY SONTYAM. ANANDAPURAM, VISAKHAPATNAM - 531 173	ANSWER KEY AND SCHEME OF EVALUATION DCMT SUPPLY EXAMINATION KEY & SCHEME	 Describe multiply excited magnetic field system? Ans: Electro-mechanical transducers have the special requirement of producing an electrical signal proportional to forces or velocities or producing force proportional to electrical signal (current or voltage). Such transducers require two excitation one excitation establishes a magnetic field of specified strength while the other excitation produces the desired signal (electrical or mechanical). Also continuous energy conversion devices motors and generators require multiple excitation. 	2) what is armature reaction ? what are the effects of armature reaction on the performance of dc machines? Ans: The current flowing through the armature conductors creates a magnetic field, which is called as ormature flux. This armature flux distorts and weakens the magnetic flux produced by the main poles. This effect of armature flux on the main flux is known as armature reaction.	 why a starter is required to start a dc motor? What is the essential element of a starter? Ans) Starters are used to protect DC motors from damage that can be caused by very high current and torque during startup. They do this by providing external resistance to the motor, which is connected in series to the motor's armature winding and restricts the current to an acceptable level. What is the function of a transformer? 	Ans) A transformer is a device that transfers electric energy from one alternating-current circuit to one or more other circuits, either increasing (stepping up) or reducing (stepping down) the voltage. 5) Define sumpners test? Ans) A full-load test on the large transformer is to be conducted to determine the maximum temperature rise of the transformer which is called as back-to-back test. The back-to-back test is also known as Commond	

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Assuming line asily The sent in the man where the field energy is given by soluting for 1 & 12 in borns of Mr. M. and substituting the egn (2) we know that is a mp (r,x) like this 12 - 2 My (21/22,0) 12 - 2 My (21/22,0) 322 - 201 いた い、た、の) こと PIL シーキアレントとちょうと Prin LILLANTE Tr = Juf (11, 2, 0) - 1) - PIN-PIL- (-11-42-Swillsoly OL = L112+1 21 - The flux produced by Ci links with its own collers seemhan

the self & multial induction as two entiting coils one functions of make 0 in linear case wit (11,12,10 = 2 41,11 + 42,111+ + 2222 If currents are used to describe the system state where inductances the functions of anyle 0 Coencert Why Chiles () = f 1 x1 d1+ f 2 d1-114 (1, 1, 1, 1) = J to 1, 1+ A2) d > + f (A2 + A2 + A2 +) 12= B-12, +B222; B-1=B12 TE = BINE (11, 12 0) - Billy 1 4 + belly 7 4 4 + fr 1 4 4 4

6b) Derive the expression for torque in a singly excited system.

Breause the merenerical changes diarid of the independent-Since dr is differented, is not preserve on left hand size of F dx = idx + / di + i dx - (2 m di + 2 m f d x) ond di is not present in the left hand side of coma E- X = [Jult - J] di + DWFdX - @ Its areflictions on wight hand side must be two lie We with the two (1,x) we are and a the area of the are Ffdx = - JWE JA + (i- JWA) dA Frax=- xdi + bug di - but dx J. NO I. X arx - time we have the is a white it is a first and a standard and the standard and t Pystem The mechanical workedme by the ALL when the asmedure indleated which drives the Muchanical system = field produces a medanical force To in the Interior michanical energy output = electrical energy input - inscale in field enorgy a disponce dx in positive direction is As per painciple of energy conservation Frax = 1d> - dwf - () No thou that wife = i & - we chy) - in with the egn () in egn (), we have grees every & rectanial date " dwn = Fr Jr tho Ve

Ans)

In fim eqn (1: 1's obvious that the fince aces in a direction petermination of meconital rivere values at Fr becaused finite Dx. obvisiosly Fristle Man thread cash marked force is given by the public demastres of Conversity of choosty Same in each case as $\Delta X - 70$ is $di = \int Lidi Linesh case <math>N_{f} R(i, X) = \int X di = \int Lidi$ othere there expressions will give slightly depleant Number then FF = - 34 (1,x) - 0 Fr =- AINC - 314 - D - (x'x) Jule = ! Et - ANTE - Firsting - D 1-051 1N1 (1,x) = 1 L(x) 12-(F) -16 111 7a) for a singly excited system derive the expression for magnetic field energy stored? To increase the induction of the exciting cot → mp((p, c) = f tras pt ⇒ いれ(x,x) = +x(n) mp (2) x 1 - 2 x - 2 (Tree) 22 - 2 (Tree) 24 - 2 x -P.F = - 2 mp = - 12 billing)

on incoverable loss in every takes placedue to by structure The assumption remaines the ideal colland the manage crevit × \ > ×1 > ×2 the relation 1-2 of F-2 , non linear The churry obserbed by the AUL for finite change BS the flue in the magnesic child unewgoes acycle -s the arcsoly absorbed by the magnesic system to whallish on a considered fight orosy interchange between flue of larthur tintages >) from initial two fluer is In flue linkages λ_1 . i. (x) dx = $\int_0^{0.2} F(\psi) d\phi$ I'm i'm lotreship is nogult toben cure which voics with the configuration volate x themaking so that the net enorgy is consumed. WAR = J IENAN = PECADAN of the alt gog blus the armabun many the Total reluidence (MT) & He cushe waller with position X of the magnetic publi decrains as and eddy current in the tren, d1-> d2-> d1 X memarky Brineline, where durp is the change in fich entry in time df. dure = eidt = dx idt = ids - Hinge the electric eventy If into the Ideal coil due to there gives all the 'x' mitteed to with blad Hurb KXE the flow of Lurrevo is in those dot is presenting that the time thing that the asmabuse is $\begin{bmatrix} dt & = i \ d(n\phi) \\ = ni \ d\phi \end{bmatrix}$ dp J J J dwa- ei dt due = dup = ci de (ei dt). (F= NI, mmf) strend in the magnesic field 1 696 ENERGY IN HYGNELIC STAEN シャンド v = 1242 the is the ロセット frue the eges

- As get the - frich an = fr FCOds the field energy is the most blue h-antis and int -> & chomps In > with fited × (uleuloro magnetic -rip x is allowed to champe with fixed > The second with the The 1-> relationship for would value of & A the mechanical system, (ellectro i muchinial) Indicaded in figure. emospy indictionale (V=iRtds enery will inductional blu the magnate check and the mobilionship can be expressed as Inf > is the independent valueble to as A HEND Term, co-energy Is now defined as The field cruzy is in growing a fundion of two restructed i.e. wh = wh(yx) (x) it is the Independent vosions le 1=1(1,1) King (1,2) = 1> - Mp(x,x) いた-いたついろ due = idx = Edp = dw maguic where in by expressions > as X(i, x), The independence The field energy is distributed throughout the space. Occupied by the field . Assoning NO losses contromatility vasiable of why becomes londy. arrea of the 1-7 cuare. IN/f= Ji X di The country in fig is shown to be the company LINIAN CASE Emmity density of the fick whe of Has the - the The fresh crossy can be expressed as which is L'all indubance L= À · テーントン = と F女 = と For have a pre-· · · P = F/\$ = reludance of the mospatic change INPCYN = FYF 2 I will write I as a with the MC = County inaris

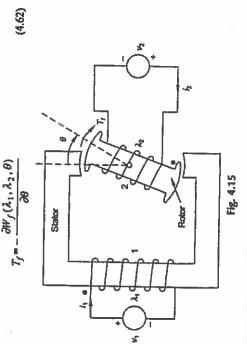
7b) write in brief about multiply excited magnetic system?

Ans) Multiply Excited Magnetic Field System – Singly-excited devices discussed earlier, are generally employed for motion through a limited distance or rotation through a prescribed angle.

Electro-mechanical transducers have the special requirement of producing an electrical signal proportional to forces or velocities or producing force proportional to electrical signal (current or voltage). Such transducers require two excitation one excitation establishes a magnetic field of specified strength while the other excitation produces the desired signal (electrical or mechanical).

Also continuous energy conversion devices motors and generators require multiple excitation.

Figure shows a magnetic field system with two electrical excitations one on stator and the other on rotor. The system can be described in either of the two sets of three independent variables ; $(\lambda_1, \lambda_2, \theta)$ or (i_1, i_2, θ) . In terms of the first set



where the field energy is given by

$$W_{j}(\lambda_{1}, \lambda_{2}, \theta) = \int_{0}^{\lambda_{1}} i_{j} d\lambda_{1} + \int_{0}^{\lambda_{1}} i_{j} d\lambda_{2}$$

Analogous to Eq. (4.28)
 $\partial W_{j}(\lambda_{1}, \lambda_{2}, \theta)$

$$i_1 = \frac{\partial W_f(\lambda_1, \lambda_2, \theta)}{\partial \lambda_1}$$
$$i_2 = \frac{\partial W_f(\lambda_1, \lambda_2, \theta)}{\partial \lambda_2}$$

Assuming linearity

(4.63)

$$\begin{split} \lambda_1 &= L_{11} i_1 + L_{12} i_2 \\ \lambda_2 &= L_{21} i_1 + L_{22} i_2; \ (L_{12} = L_{21}) \end{split} \tag{4.64b}$$

Solving for it and is in terms of λ_1, λ_2 and substituting in Eq. (4.63) gives upon integration

$$W_{j}(\lambda_{1}, \lambda_{2}, \theta) = \frac{1}{2} \beta_{11} \lambda_{1}^{2} + \beta_{12} \lambda_{1} \lambda_{2} + \frac{1}{2} \beta_{22} \lambda_{2}^{2}$$
 (4.65)

where

$$\begin{array}{l} \beta_{11} = L_{22}/(L_{11}L_{22}-L_{12}^2) \\ \beta_{22} = L_{11}/(L_{11}L_{22}-L_{12}^2) \\ \beta_{12} = \beta_{21} = -L_{12}/(L_{11}L_{22}-L_{12}^2) \end{array}$$

The Self- and mutual-inductance of the two exciting coils are functions are angle 0.

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If currents are used to describe the system state

$$T_f = \frac{\partial W_f(l_1, l_2, \theta)}{\partial \theta}$$
(4.66)

where the co energy is given by

$$W_{f}^{r}(l_{1}, l_{2}, \theta) = \int_{0}^{t_{1}} \lambda_{1} di_{1} + \int_{0}^{t_{2}} \lambda_{2} di_{2}$$
 (4.67)

In the linear case

$$W'_{f}(i_{1},i_{2},\theta) = \frac{1}{2}L_{11}i_{1}^{2} + L_{12}i_{1}i_{2}^{2} + \frac{1}{2}L_{22}i_{2}^{2}$$

where inductances are functions of angle 0

8a) Explain the constructional details of dc generator?

Ans) Construction of DC machines : A D. C. machine consists of two main parts 1. Stationary part: It is designed mainly for producing a magnetic flux. 2. Rotating part: It is called the armature, where mechanical energy is converted into electrical generate) or conversely electrical energy into mechanical (electric into) Parts of a De Generator: 1) Yoke 2) Magnetic Poles a) Pole Sche 3) Field Winding 4) Armature Corte 5) Armature winding 6) Commutator 7) Brushes and Bearings The stationary parts and rotating parts are separated from each other by an air gap.

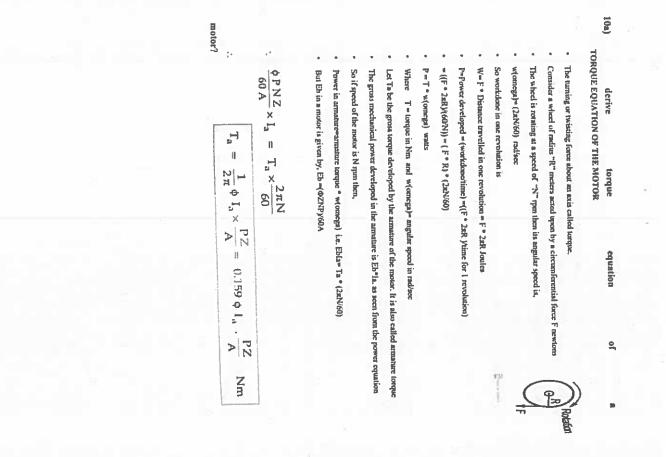
The stationary part of a D. C. machine consists of main poles, designed to create the magnetic flux, commutating poles interposed between the main poles and designed to ensure spark less operation of the brushes at the

magnetic-field flux Φ . When the armature winding rotates in this magnetic field so to cut the flux, the voltage Ea is	commutator surface by means of a spring, whose tension can be adjusted with the help of lever. A flexible copper
Ans : A DC generator, is an electrical machine that converts the mechanical energy of a prime mover (e.g. DC motor, AC induction motor or a turbine) into direct electrical energy. The generator shown in figure 1 is self-exciting. It more the voltage Fa generated by the machine to establish the field current If, which in turn gives rise to the	mica. Brushes and brush gear: Brushes are stationary and rest on the surface of the Commutator. Brushes are rectangular in shape. They are housed in brush holders, which are usually of box type. The brushes are made to press on the
9a) Explain magnetisation characteristics of a de shunt generator ?	cylindrical in shape and is made up of wedge shaped segments which are insulated from each other by thin layer of
=62208/120 ≈ 518.4 volts	case of motor. Choice of material: As it collects current from armature, it is also made up of copper segments. It is
emf generated = 0.016*810*1200*4\60*2	2. To convert internally developed alternating emf to in directional (dc) emf 3. To produce unidirectional torque in
no of conductors =18 *45 = 810	facilitate the collection of current from the armature conductors.
for wave winding A =2	Commutator. The rectification in case of dc generator is done by device called as commutator. Functions: 1. To
a = number of parallel paths	winding'.
p = number of poles	helps in producing the magnetic field i.e. exciting the pole as electromagnet it is called 'Field winding' or 'Exciting
Z = total number of armature conductors	Field winding: The field winding is wound on the pole core with a definite direction, runctions, to carry current was to the pole core on which the winding is placed behaves as an electromagnet, producing necessary flux. As it
n = annature speed in rpm	
$\Phi = $ flux in webers	winding in case of generators. 2. To carry the current supplied in case of dc motions, 3. 10 do the useful work it use
Ans $E = (N P \not \otimes Z)(60 A)$	armature conductors and emf gets induced in them. Function, 1. Generation of emf takes place in the armature
8b) calculate the emf generated by a 4pole wave wound armature having 45 slots 18 conductors per slot .when driven at 1200 rpm the flux per pole is 0.016 wb.	Armature winding: Armature winding is nothing but the inter connection of the armature conductors, placed in the slots provided on the armature core. When the armature is rotated, in case of generator magnetic flux gets cut by
	like cast iron or cast size.
COT MAJAGET	Armature: It is further divided into two parts namely, (1) Armature core (2) Armature winding. Armature core is cylindrical in shape mounted on the shaft. It consists of slots on its periphery and the air ducts to permit the air flow through armature which serves cooling purpose. Functions: 1. Armature core provides house for armature winding i.e., armature conductors. 2. To provide a path of low reluctance path to the flux it is made up of magnetic material
A A A A A A A A A A A A A A A A A A A	larger induced emf. To achieve this, pole core has been given a particular shape.
A Marine	winding which is necessary to produce the flux. 2. If alreads the flux produced unough an gap to animate one or the flux, pole 3. Pole shoe enlarges the area of annature core to come across the flux, which is necessary to produce
Milling Strange	Poles: Each pole is divided into two parts a) pole core b) pole shoe Functions: 1. Pole core basically carries a field
	nagnetic flux.
volt 1 york Arganiae	Yoke: 1. It saves the purpose of outermost cover of the dc machine so that the insulating materials get protected room harmful atmospheric elements like moisture, dust and various gases like SO2, acidic fumes etc. 2. It provides mechanical support to the poles. 3. It forms a part of the magnetic circuit. It provides a path of low reluctance for
conductor called pigtail is used to connect the brush to the external circuit. Functions: To collect current from commutator and make it available to the external circuit.	commutator and a frame / yoke. The armature is a cylindrical body rotating in the space between the poles and comprising a stotted armature core, a winding inserted in the armature core slots, a commutator and brush.

 C. and anthor of memory or molecular system of a condense version of a con	induced in the armature. This voltage is commonly referred to as the armature electromotive force or EMF. The induced EMF is proportional to the rate of cutting the flux and is given by emf formula.	Let P = number of poles Ø = flux/pole in webers
the D.C. Shunt Generator will fail to excite. traisfance becomes the critical field resistance. Critical field resistance traisfance becomes the critical field resistance. af determining the folge of the tangent to the linear position of the curve initial position of the O.C.C is neglected. all amount of EMF is generated EMF is directly proportional to the consequently generated EMF is directly proportional to the ansities where rel is canall iron path relutance becomes uppreciable and the amature, so it is called DC Shunt Generator. Due to residual urrent will be generated. This current while pussing into the field coils is will increase pole flux which will further increase the enfthe illbrium is reached at some point in graph between IF and EQ.		Z = total number of armature conductors. = number of slots x number of conductors/slot
		N = armature rotation in revolutions (speed for armature) per minute (rpm)
the D.C. Shunt Generator will fail to excite. resistance becomes the critical field resistance. Critical field resistance training the slope of the tangent to the linear position of the curve of determining the slope of the tangent to the linear position of the curve initial position of the O.C.C is neglected. all amount of EMF is generated even when IF = 0. Hence the curve the lower end is due to magnetic intertia. It has seen that the first part of the lower end is que to magnetic intertia. It has seen that the first part of the lower end is generated EMF is directly proportional to the consider where it is small iron path reluctance becomes appreciable and of the armature, so it is called DC Shunt Generator. Due to residual threat while generated. This current while passing into the field coils is will increase pole flux which will further increase the ent. In it is trached at some point in graph between IF and EG.		A = No, of parallel paths into which the 'z' no. of conductors are divided.
the D.C. Shunt Generator will fail to excite. resistance becomes the criticeal field resistance. Critical field resistance addetermining the slope of the tangent to the linear position of the curve i determining the slope of the tangent to the linear position of the curve i all amount of EMF is generated even when If = 0. Hence the curve all amount of EMF is generated EMF is directly proportional to the resides where it is small iron path reluctance becomes appreciable and flux and consequently generated EMF is directly proportional to the resides where it is small iron path reluctance becomes appreciable and flux and consequently generated EMF is directly proportional to the resides where it is small iron path reluctance becomes appreciable and flux and consequently generated EMF is directly proportional to the resides where it is small iron path reluctance becomes appreciable and flux and consequently between IF and EQ.		$\mathbf{E} = \operatorname{cmf}$ induced in any parallel path
the D.C. Shunt Generator will fail to excite. resistance becomes the critical field resistance. Critical field resistance addetermining the slope of the tangent to the linear position of the curve e initial position of the O.C.C is neglected. and amount of EMF is generated even when If = 0. Hence the curve and amount of EMF is generated EMF is directly proportional to the e lower end is due to magnetic inertia. It has seen that the first part of flux and consequently generated EMF is directly proportional to the the lower end is called DC Shunt Generator. Due to residual fur the farmature, so it is called DC Shunt Generator. Due to residual urrent will be generated. This current while passing into the field coils is will increase pole flux which will further increase the cnf. fibrium is reached at some point in graph between IF and EG.		Eg = emf generated in any one parallel path in the armature.
the D.C. Shunt Generator will fail to excite. resistance becomes the critical field resistance. Critical field resistance addetermining the slope of the tangent to the linear position of the curve initial position of the O.C.C is neglected. all amount of EMF is generated even when If = 0. Hence the curve be lower end is due to magnetic inertia. It has seen that the first part of flux and consequently generated EMF is directly proportional to the consequently generated EMF is directly proportional to the and amount of EMF is senated EMF is directly proportional to the constities where it is small inon path reluctance becomes appreciable and I to the annature, so it is called DC Shunt Generator. Due to residual urrent will be generated. This current while passing into the field coils is will increase plot flux which will further increase the cmf.the filbrium is reached at some point in graph between IF and EQ.		Average emf generated/conductor = dt0/dt volt
the D.C. Shunt Generator will fail to excite. resistance becomes the critical field resistance. Critical field resistance d determining the slope of the tangent to the linear position of the curve e initial position of the O.C.C is neglected. all amount of EMF is generated even when If = 0. Hence the curve is due to magnetic inertia. It has seen that the first part of the amature, so it is called EMF is directly proportional to the ensities where it is small iron path reluctance becomes appreciable and It the armature, so it is called DC Shum Generator. Due to residual urrent will be generated. This current while passing into the field coils is will increase pole flux which will further increase the enfithe if the information is reached at some point in graph between IF and EQ.		Flux current/conductor in one revolution $dt = d x p \ln one$ revolution,
the D.C. Shunt Generator will fail to excite. resistance becomes the critical field resistance. Critical field resistance d determining the slope of the tangent to the linear position of the curve e initial position of the O.C.C is neglected. all amount of EMF is generated even when If = 0. Hence the curve is due to magnetic inertin. It has seen that the first part of flux and consequently generated EMF is directly proportional to the consider where it is small iron path reluctance becomes appreciable and I to the armature, so it is called DC Shunt Generator. Due to residual urrent will be generated. This current while passing into the field coils is will increase pole flux which will further increase the enf. the ultimut is reached at some point in graph between IF and EQ.		the conductor will cut total flux produced by all poles = $d x p$
the D.C. Shunt Generator will fail to excite. Tresistance becomes the critical field resistance. Critical field resistance determining the slope of the tangent to the linear position of the curve e initial position of the O.C.C is neglected. all amount of EMF is generated even when If = 0. Hence the curve he lower end is due to magnetic inertia. It has seen that the first part of flux and consequently generated EMF is directly proportional to the ensities where it is small iron path reluctance becomes appreciable and I to the armature, so it is called DC Shunt Generator. Due to residual urrent will be generated. This current while passing into the field coils is will increase pole flux which will further increase the emf.the ilibrium is reached at some point in graph between IF and EG.		No. of revolutions/second = N/60
resistance becomes the critical field resistance. Critical field resistance determining the slope of the tangent to the linear position of the curve e initial position of the O.C.C is neglected. all amount of EMF is generated even when If = 0. Hence the curve be lower end is due to magnetic inertia. It has seen that the first part of flux and consequently generated EMF is directly proportional to the ensities where it is small iron path reluctance becomes appreciable and I to the armature, so it is called DC Shunt Generator. Due to residual urrent will be generated. This current while passing into the field coils is will increase pole flux which will further increase the enf. the illibrium is reached at some point in graph between IF and EG.	ch the D.C. Shunt Generator will fail to excite.	Therefore, Time for one revolution, dt = 60/N second According to Faraday's laws of Electromagnetic Induction, emf d Θ generated/conductor = $d\Theta \times p \times N / 60$ volts
e initial position of the O.C.C is neglected. and amount of EMF is generated even when If = 0. Hence the curve be lower end is due to magnetic inertia. It has seen that the first part of flux and consequently generated EMF is directly proportional to the ensities where it is small iron path reluctance becomes appreciable and I to the armature, so it is called DC Shunt Generator. Due to residual urrent will be generated. This current while passing into the field coils is will increase pole flux which will further increase the enfithe illbrium is reached at some point in graph between IF and EQ.	eld resistance becomes the critical field resistance. Critical field resistance and determining the show of the summer of starting of the summer of the summer of the summer of the summer of the	This is carf induced in one conductor. For a simplex wave-wound generator No. of parallel paths = 2 No. of
all amount of EMF is generated even when If = 0. Hence the curve be lower end is due to magnetic intertia. It has seen that the first part of flux and consequently generated EMF is directly proportional to the ensities where it is small iron path reluctance becomes appreciable and I to the armature, so it is called DC Shunt Generator. Due to residual urrent will be generated. This current while passing into the field coils is will increase pole flux which will further increase the emf.the ulthrum is reached at some point in graph between IF and EG.	e linear position of the	conductors in (series) in one path = $Z/2$
the lower end is due to magnetic inertia. It has seen that the first part of flux and consequently generated EMF is directly proportional to the ensities where it is small iron path reluctance becomes appreciable and I to the armature, so it is called DC Shunt Generator. Due to residual urrent will be generated. This current while passing into the field coils is will increase pole flux which will further increase the emf.the uitbrium is reached at some point in graph between IF and EG.	small amount of EMF is generated even when if = 0. there is a second of the second of	EMF generated/path = OPN/60 x Z/2 = O2PN/120 volt
flux and consequently generated EMF is directly proportional to the ensities where it is small iron path reluctance becomes appreciable and I to the armature, so it is called DC Shunt Generator. Due to residual urtent will be generated. This current while passing into the field coils is will increase pole flux which will further increase the enf.the ultiturn is reached at some point in graph between IF and EQ.	t the lower end is due to magnetic inertia. It has seen that the first part of	For a simple lap-wound generator Number of parallel paths = P Number of conductors in one path = Z/P
I to the armature, so it is called DC Shunt Generator. Due to residual urrent will be generated. This current while passing into the field coils is will increase pole flux which will further increase the cmf.the ultibrium is reached at some point in graph between IF and EG.	a flux and consequently generated EMF is directly proportional to the	EMF generated path = 0PNv60 (Z/P) = 0ZNv60
is will increase pole flux which will further increase the en- ultibrium is reached at some point in graph between IF and EG.	densities where it is small iton path reluctance becomes appreciable and lel to the armature, so it is called DC Shunt Generator. Due to residual current will be generated. This current while mession into the state and	A = 2 for simplex – wave winding $A = P$ for simplex lap-winding .
	his will increase pole flux which will further increase the enfulte utilibrium is reached at some point in graph between IF and EG.	

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sd Ia = 250 A , P = 12 (10b) A dynamo has a raped annulum HP II AZ=P number of perales paths with connected ! the machine has 12 poles. annabure winding is lop a wave per parts of the armabure if the IL = 150 = 20.83 Amps I - Ia current per poli in case of TI= 150 = 125 A ALT - number of parallel pathis with current at 250 A, what is the current AI wave winding in Amps In Currend-per path in case lap winding wave winding Az of Inp windry in Amps

11a) explain speed control of de shunt motor?

Ans) The speed equation of a DC motor shows that,

N a Eb/ ዋ

or N a (V - laRa) o

But the resistance of armature winding is small. Therefore the voltage drops laRa will be negligible as compared to the external supply voltage V. Therefore, the expression for the speed can be approximated as follows:

N a V/ (because V>> laRa)

From this expression we can obtain the dc shunt motor speed control methods.

The speed is inversely proportional to flux φ .

It is directly proportional to armature voltage drop (laRa).

It is directly proportional to applied Voltage V.

Speed Control of DC Shunt Motor:

So by varying one of these parameters, it is possible to change the speed of DC shunt motor. Depending on the parameter being controlled, methods of speed control of shunt DC Motor are classified as follows:

Flux Control Method:

speed control of dc shunt motor

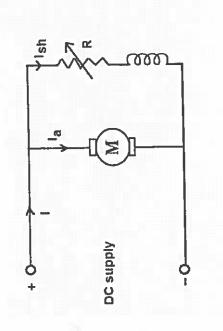
The set up for speed control of DC shunt motor using flux control technique is shown in the figure. In order to change the speed, we have to change flux. This can be achieved by changing the field current. The field current can be changed by changing the rhoostat R connected in series with the field At the time of starting the motor, we need to run the motor slowly, therefore, the flux should be maximum, because,

Na 1/9

To obtain maximum flux at the start, the field current should be maximum at the time of starting. To obtain this, the value of rheostat (R) should be minimum.

The speed of DC shunt motor can be varied by varying the field current. As we increase the resistance R of the rhoostat, the field current ish decreases. So the flux ϕ decreases. This results in increasing the speed of the motor. As the R is increased, the speed increases.

We can use this technique to control motor speed above its rated value. It is the most commercial method. There is a limit to the maximum obtainable speed by this method due to poor commutation at week fluxes. Most common maximum to minimum speed ratio is 6:1.



Speed Control of DC Shunt Motor by Armature:

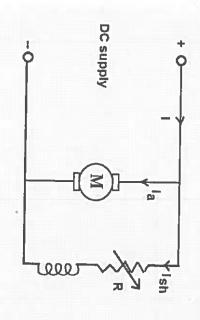
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speed control of dc shunt motor theory

The set up for speed control of DC shunt motor by armature voltage control method is shown in the figure. A rhoeostat is connected in series with armature winding. By varying the value of R we can vary the voltage across the armature.

Because speed N is directly proportional to armature voltage, it is possible to change the speed by changing the value of rheostat R.

method because, in this method, speed is reduced at the cost of power rheostat([a2R). We can use this technique to control motor speed below its rated value. But it is neither efficient nor economical



11b) explain in detail the losses in dc motor?

Ans:

The various losses in a dc machine whether it is a motor or a generator are classified into three

groups as:

1: Copper losses

2. Iron or core losses

3. Mechanical losses

Copper losses

• The copper losses are the losses taking place due to the current flowing in a winding

There are basically two windings in a dc machine namely armature winding and field winding.

The copper losses are proportional to the square of the current passing through these windings.

Armature copper loss = Ia'2 * Ra

Shunt field copper loss = Ish'2 * Rsh

Series field copper loss = lse'2 * Rse

Iron or core losses

is given by • These losses are also called magnetic losses , these losses include hysteresis loss and eddy current loss • The hysteresis loss is proportional to the frequency and the maximum flux density Bm in the air gap and

This loss is basically due to reversal of magnetisation of the armature core

Hysteresis loss H n B^{1.6} f V watts

۲ 11 Steinmetz hysteresis coefficient

∨ ≋ Volume of core in m³

where

h H Frequency of magnetic reversals

rddy current lass

The eddy current loss exists due to eddy currents.

This induced emf sets up eddy currents which cause the power loss , this loss is given by When attracture cure rotates, it cuts the magnetic flux and emf gets induced in the core.

These losses are almost constant

Eddy current loss = K Bm f² t² V watts

X ŝ Constant

where

11 || Thickness of each lamination

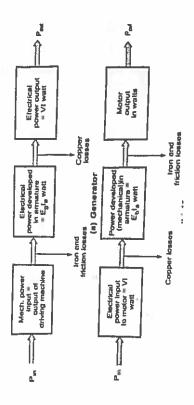
< Volume of core

f = Frequency of magnetic reversals

TOTAL LOSSES

For a dc machine

Total losses = constant losses + variable losses • The power flow and energy transformation diagrams at various stages,



12a) derive the emf eqn of a transformer?

Ans: E.M.F. Equation of a Transformer Consider that an alternating voltage V1 of frequency f is applied to the primary . The sinusoidal flux ϕ produced by the primary can be represented as: $\phi=\phi$ m sinust When the primary winding is excited by an alternating voltage V1, it is circulating alternating current, producing an alternating flux ϕ . ϕ

NI - Number of primary turns

N2 - Number of secondary turns

f - Frequency of the supply voltage

E1 - R.M.S. value of the primary induced e.m.f

E2 - R.M.S. value of the secondary induced e.m.f

The instantaneous e.m.f. el induced in the primary is From Faraday's law of electromagnetic induction - The flux increases from zero value to maximum value ϕ_m in 1/4f of the time period that is in 1/4f seconds. The change of flux that takes place in 1/4f seconds = $\phi_m - 0 = \phi_m$ weber.

 $\frac{d\varphi}{dt} = \frac{dt}{1/4f} = \frac{4f\varphi_{m}}{h_{m}} w_{m}/sec.$

Since that ϕ varies simuvaliably, the R.m.s value of the induced c.m.f is obtained by multiplying the average value with the form lizetur

Form factor of a sinwary Average value [,] }

R.M.S Value of c.m.f induced in one turns = $4\varphi_m f \approx 1.11$ Volus.

- 4,440 L Volta

R.M.S Vakie of a mf induced in primary winding = 4.444a f N1Volts.

K.M.S. Vakie of c.m.f artheed in secondary winding = 4,440m f N2Virks.

The cupter with of E_1 and E_2 are called c in f equation of a transformer

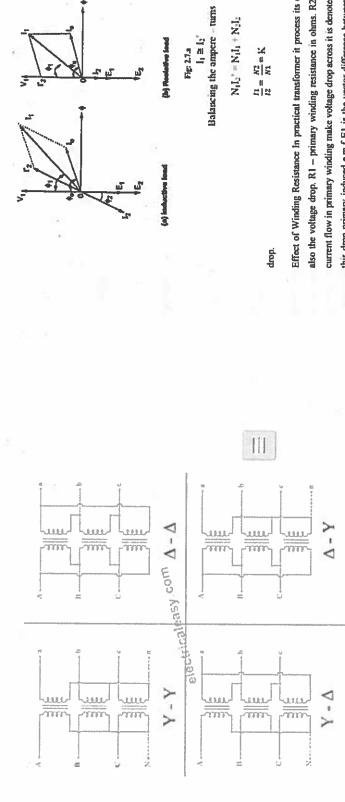
	alts.	olts.
l	N'NI"	~~~
	41	ŧ.
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l	Š,	11

N1 _ 130 X200 = 115 N1 _ 100 (TUMS M010) 1218 400 2300, SO HZ This takes X_{1}^{1} $N_{1} = 200 \text{ This}$ $K = \frac{1}{2} = \frac{N_{1}}{N_{1}} = \frac{N_{1}}{N_{1}}$ $K = \frac{1}{2} = \frac{N_{2}}{N_{1}} = \frac{1}{2} \frac{1}{2}$ $K = \frac{1}{2} = \frac{N_{2}}{N_{1}} = \frac{1}{2} \frac{1}{2}$ $K = \frac{1}{2} = \frac{N_{2}}{N_{1}} = \frac{1}{2} \frac{1}{2}$ Transon HU SILC FI- 400 V Chn = 4.00 4.00 - 4.44 × 50 × 100 4.00 - 4.44 × 50 × 100 4.00 - 4.44 × 50 × 100 Am = 0.00 00 wb

[3a] draw the various types of three phase transformer connections and brief each one of them?

Ans: Three Phase Transformer Connections Three phase transformer connections In three phase system, the three phases can be connected in either star or delta configuration. In case you are not familiar with those configurations, study the following image which explains star and delta configuration. In any of these configurations, there will be a phase difference of 120° between any two phases.

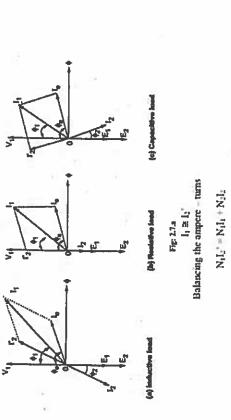
required phase/turns is relatively greater than that for star-star connection. . The ratio of line voltages on the primary balanced. Delta-Delta (Δ-Δ) • This connection is generally used for large, low-voltage transformers. Number of Line voltages on both sides are in phase with each other. • This connection can be used only if the connected load is line voltages on the primary side and the secondary side is equal to the transformation ratio of the transformers. • connection is 1/43 times of line voltage only). Thus, the amount of insulation required is also reduced. • The ratio of voltage transformers. Because of star connection, number of required turns/phase is reduced (as phase voltage in star These configurations are explained below. Star-Star (Y-Y) • Star-star connection is generally used for small, highconfigurations as (i) star -star, (ii) delta-delta, (iii) star-delta, (iv) delta-star, (v) open delta and (vi) Scott connection. Three Phase Transformer Connections Windings of a three phase transformer can be connected in various Δ) - The primary winding is star star (Y) connected with grounded neutral and the secondary winding is delta can continue to operate in open delta connection but with reduced available capacity. Star-Delta OR Wye-Delta (Y unbalanced loading. • Another advantage of this type of connection is that even if one transformer is disabled, system and the secondary side is equal to the transformation ratio of the transformers. • This connection can be used even for and the secondary winding is connected in star with neutral grounded. Thus it can be used to provide 3-phase 4-wire primary and secondary line voltages. Delta-Star OR Delta-Wye (Δ-Y) - The primary winding is connected in delta connected. . This connection is mainly used in step down transformer at the substation end of the transmission line. Delta (V-V) Connection Two transformers are used and primary and secondary connections are made as shown in t and secondary line voltages. Above transformer connection configurations are shown in the following figure. Open ratio of secodary to primary line voltage is v3 times the transformation ratio. • There is 30° shift between the primary service. • This type of connection is mainly used in step-up transformer at the beginning of transmission line. • The The ratio of secondary to primary line voltage is 1/v3 times the transformation ratio . There is 30° shift between the



13b) Draw and explain the phasor diagram of single phase transformer on load considering with winding resistance.

with 12 and phase of 12 is decided by the load. iv) Primary current 11 is vector sum of Io and 12' a) If load is Inductive, 12 lags E2 by \$2, shown in phasor diagram fig 2.7 (a). b) If load is resistive, 12 in phase with E2 shown in Ans) I Phasor Diagram i) Take (\$) flux as reference for all load iii) The load component 12', which is in anti-phase phasor diagram fig. 2.7 (b). c) If load is capacitive load, 12 leads E2 by \$2 shown in phasor diagram fig. 2.7 (c). For

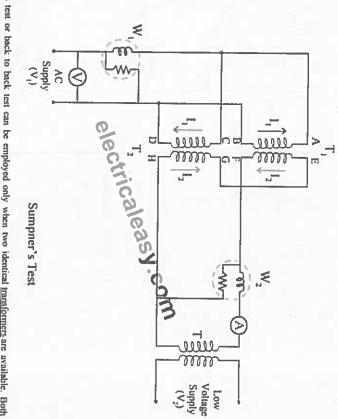
to V2 neglecting various is equal ជ here we assumed casy understanding at this stage



Effect of Winding Resistance in practical transformer it process its own winding resistance causes power loss and also the voltage drop. R1 - primary winding resistance in ohms. R2 - secondary winding resistance in ohms. The current flow in primary winding make voltage drop across it is denoted as IIR1 here supply voltage V1 has to supply this drop primary induced e.m.f E1 is the vector difference between V1 and 11R1. Similarly the induced e.m.f in secondary E2, The flow of current in secondary winding makes voltage drop across it and it is denoted as I2R2 here E2 has to supply this drop.

The vector difference between E2 and I2R2 (Assuming as purely resistive drop here.)

14a) draw the circuit diagram of sumpners test and derive the equation for efficiency of each transformer.



Ans)

Sumpner's test or back to back test can be employed only when two identical <u>transformers</u> are available. Both transformers are connected to supply such that one transformer is loaded on another. Primaries of the two identical transformers are connected in parallel across a supply. Secondaries are connected in series such that emfs of them are opposite to each other. Another low voltage supply is connected in series with secondaries to get the readings, as shown in the circuit diagram.

In above diagram, T₁ and T₂ are identical transformers. Secondaries of them are connected in voltage opposition, i.e. E_{EF} and E_{att} . Both the emfs cancel each other, as transformers are identical. In this case, as per superposition theorem, no current flows through secondary. And thus the no load test is simulated. The current drawn from V₁ is 2la, where lo is equal to no load current of each transformer. Thus input power measured by wattmeter W₁ is equal to iron loss per transformer Pi = W₁/2.

Now, a small voltage V_2 is injected into secondary with the help of a low voltage transformer. The voltage V_2 is adjusted so that, the rated current I_2 flows through the secondary. In this case, both primaries and secondaries curry rated current. Thus short circuit test is simulated and wattmeter W_2 shows total full load copper losses of both transformers.

copper loss per transformer P_{co} =

W1/2

5

From above test results, the full load efficiency of each transformer can be given as -



14b) The primary and secodary voltages of an autotransformer are 1200v and 600v respectively .calculate the economy of cu when the secondary current is 120A .draw the circuit and show the current disturbiion in the winding .

The economy of cu when the scondary (1)b) the polynowy and secondary withager, of an cuedo one secondary withager, Current is 120 A. Draw the circuit and show the current district builtion in the souting = kw = 0.5× wo = 0.5 Wo TI = KIL = 0.5 X120A=GA K= Ve = 8000 = 7 = 0.5 and 600 V respectively. calculate i pur condage of saving = 0.5 X100 1200 3 120 A 600V -.20.-1204 winding. Goa 607 23

Ans)

15a) A 10 kva 500/250 v 50 hz single phase transformer gave the following test data

Oc test (LV test):250 V,1.0A,80W

SC TEST (H V SIDE):25 V,12A,100W where LV refers to the low voltage and HV side refers to high voltage side .determine the follwing:

Ans) Xu = V2 = 250 = 263.8800 JL Imay 0.9474 = 263.8800 JL To = 1:0 A TW = WO = 80 SE TEST (LV Side) , 250V 1.0A 80W SE TEST (HV Side) : 25 V, 12A, 100AU VIGE) ILES, 126, 42 202 Themportur rating = 10 KV A open chicale test perfectmenter on w site Imag = VII-TUT G1 = 5001, E2 = 2501 Pu= 1/2 = 150 = 781.25 J_ -- Stell - VI-0:1024 1= J(1,0)- (0.3) = J 0.8976 = 0.9474A 1.1.2.0 Lalliers courses - 0.32A L

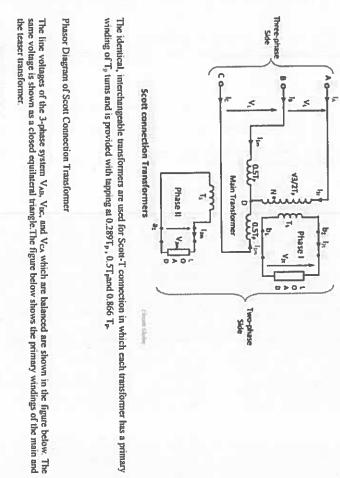
15 b) explain how the scott connection can be used to ontain two phase supply from a three phaswe supply? Key - Kr Kep = (0.5) 1.9645 - 6.4960 Pris = K² Rep = 6.5² 0.694450.17 Vi = Gi-Ji Rescust - Fixersing full load currie on Luride Ty 1000 z רר וו perinmeters referred to LU SIJE V2= 232.66V Ans)

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10.69 L shart circuit act performed on the stip 0-6944 1282.0 love. 4 -25 - 2.0833-1-1-105 VI--EL-ILRES CUSTO -INRESSING Transformation valio = 4 - 5 - 950 Xep = V(Zep)-leps = Jle. USAJfull load current on the side -10×1000 252 00 1 - 20 A 1296.15 (12)2when pp costs = 0.8 las 00 200 2.0= duis Zep = 1(20) ((Ist)) A 4 L(130) у, З All Philade Ker

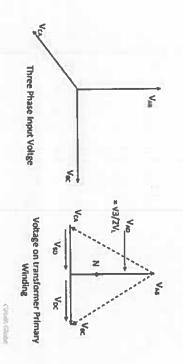
The Scott-T Connection is the method of connecting two single phase transformer to perform the 3-phase to 2-phase conversion and vice-versa. The two transformers are connected electrically but not magnetically. One of the transformers is called the main transformer, and the other is called the auxiliary or feaser transformer.

The figure below shows the Scott-T transformer connection. The main transformer is centre tapped at D and is connected to the line B and C of the 3-phase side. It has primary BC and secondary ata. The teaser transformer is connected to the line terminal A and the centre tapping D. It has primary AD and the secondary b₁b₂.



+ V22

 $|V_{AB}| = |V_{BC}| = |V_{CA}| = |V_L|$



The D divides the primary BC of the main transformers into two halves and hence the number of turns in portion $BD = the number of turns in portion <math>DC = T_p/2$. The voltage V_{120} and V_{10c} are equal, and they are in phase with V_{10c} .

$$V_{BD} = V_{DC} = \frac{1}{2} V_{BC} = \frac{1}{2} V_L < 0^*$$

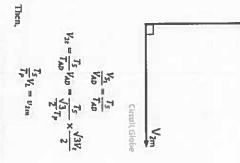
The voltage between A and D is

$$V_{\theta D} = V_{DC} = \frac{1}{2} V_{\theta C} = \frac{1}{2} V_L < 0^\circ$$
$$V_{AD} = V_{AB} + V_{\theta D}$$

$$V_{AB} = V_L \left(-\frac{1}{2} + j \frac{\sqrt{3}}{2} \right) + \frac{1}{2} V_L$$

$$V_{AD} = V_L \left(j \frac{\sqrt{3}}{2} \right) = 0.866 V_L < 90$$

The teaser transformer has the primary voltage rating that is $\sqrt{3/2}$ or 0.866 of the voltage ratings of the main transformer. Voltage V_{AD} is applied to the primary of the teaser transformer and therefore the secondary of the voltage V_{AD} of the teaser transformer will lead the secondary terminal voltage V_{AD} of the main transformer by 90° as shown in the figure below.



For keeping the voltage per turn same in the primary of the main transformer and the primary of the teaser transformer, the number of turns in the primary of the teaser transformer should be equal to 43/2 T_p.

Thus, the secondaries of both transformers should have equal voltage ratings. The V_{2n} and V_{2n} are equal in magnitude and 90° apart in time; they result in the balanced 2-phase system.

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Position of Neutral Point N

A

The primary of the two transformers may have a four wire connection to a 3-phase supply if the tapping N is provided on the primary of the teaser transformer such that

The voltage across AN = V_{AN} = phase voltage = V_VV3 .

Since the voltage across the portion AD.

 $V_{AD} = \frac{\sqrt{3}}{2} V_L$

the voltage across the portion ND

$$V_{ND} = V_{AD} - V_{AN} = \frac{\sqrt{3}}{2}V_L - \frac{V_L}{\sqrt{3}} = \frac{V_L}{2\sqrt{3}}$$

The same voltage turn in portion AN, ND and AD are shown by the equations,

$$T_{AN} = \frac{T_P}{\sqrt{3}} = 0.5777_P$$
$$T_{ND} = \frac{T_P}{2\sqrt{3}} = 0.28817_P$$
$$T_{ND} = \frac{\sqrt{3}T_P}{2} = 0.8667_P$$
$$\frac{T_{AN}}{2} = \frac{T_P}{\sqrt{3}} + \left(\frac{T_P}{2\sqrt{3}}\right) = 2$$

The equation above shows that the neutral point N divides the primary of the teaser transformer in ratio.

AN: ND=2:1

CONTEST Instruction

CEMP RADUL)

(Or, R.S.R. Krishman midu) HUD-CER VALLA

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Degree Course	B. Tech. (U. G.) Program ECE Code 20EC304 Test Duration 3 Hrs. Max. I	Marks 70	Academic Year 202 Semester	21 - 2022 III
Course	Random Variables and Stochastic Process			
Part A (Short Answer Questions 5 x 2 = 10 Marks)			
No.	Questions (1 through 5)		Learning Outcome (s)	DoK
1	State Bays theorem.		20EC304.1	L1
2	What does the Chebychev's inequality states?		20EC304.2	L1
3	List four properties of Joint distribution function.		20EC304.3	£1
4	Differentiate between Autocorrelation and Cross-correlation fu	Inction	20EC304.4	L1
5	Define narrow Pass Process.		20EC304.5	L1
Part B ((Long Answer Questions 5 x 12 = 60 Marks)			
No.	Questions (6 through 15) Let XX be a random variable with PDF given by	Marks	Learning Outcome (s)	DoK
6 (a)	$f_X(x) = \left\{egin{array}{cc} cx^2 & x \leq 1 \ 0 & ext{otherwise} \end{array} ight.$	6M	20EC304.1	L2
6 (b)	Find constant c, Variance (X) and $P(X \ge \frac{1}{2})$ Elaborate on Discrete and continuous sample spaces.	6M	20EC304.1	L2
0(0)	OR	OIM	2020304.1	LL
	Let X be a continuous random variable with PDF $\int dr^3 = 0$ for $n \neq 1$			
7 (a)	$f_X(x) = egin{cases} 4x^3 & 0 < x \leq 1 \ 0 & ext{otherwise} \end{cases}$	6M	20EC304.1	L2
7 (b)	Find P(X \leq 2/3 X >1/3) Elaborate on Conditional density and their properties.	6M	20EC304.1	L2
• /				
8 (a)	State and prove Chebychev's inequality If X be a random variable and P(X-x) is the probability mas	6M is	20EC304.2	L3
	function given as below table			
	X 0 1 2 3 4 5 6 7			
8 (b)	0 k 2k 3k 4k K ² 2 7	6M	20EC304.2	L3
	P(K ² K ²			
	X- +			
	x) k			
	Find the value of k and P(X≤6) OR			27
9 (a)	Let Y=2X+3, If the random variable is uniformly distribute	ed 6M	20EC304.2	L3
	over [-1, 2], determine fy(y) Explain about the characteristic function and state	ite		
9 (b)	properties	¹¹³ 6M	20EC304.2	L2
10	State and explain the properties of Marginal distribution	on 12M	20EC304.3	L3
	functions, Conditional distribution and density function. OR			
11 (a)	Consider two random variables X and Y with joint Probabil Mass Function given in Table below.	^{ity} 6M	20EC304.3	L3
	made i anotori giron in rabio bolow.	- (2. Keljani	
		Contro	oller of Examinati	ons
			NSRIT (A)	

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	Y = 2	Y = 4	Y = 5
X = 1	12	$\frac{1}{21}$	$\frac{1}{24}$
$X \simeq 2$	t. ti	1 12	L N
$\lambda^{'} = 3$	l) L	1 =	12

Find $P(X \le 2, Y \le 4)$, P(Y = 2|X = 1) and check X and Y are independent.

State and prove central limit theorem for equal distributions 11 (b) case

Explain about Second-order and wide-sense stationary 12 (a) process. Derive the Relationship between cross-power density 12 (b) spectrum and cross-correlation function.

OR

- Explain how random processes are classified with neat 13 (a) sketches. Derive the relationship between power spectral density and
- 13 (b) autocorrelation function.
- A random processes X(t)= Asin(ωt + θ), where A, ω are 14 (a) constants and θ is a uniformly distributed random variable on the interval $(-\pi, \pi)$ find average power?
- Derive the relation between input PSD and output PSD of an 14 (b) LTI system:

OR

Explain the following i) Noise Figure ii) Noise Sources. 15 (a) Derive the expression for average cross power between two 15 (b) random process X(t) and Y(t).

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20EC304.3

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PART-A:	ANSWE KEY AND SCHEME OF EVALUATION
1. State	Baye's theorem - 2m
2. What	does the Chebycher's inequality statement - 2m
	four properties of foort distribution-function - 2m each propaty - 1/2 mark.
4. D;	fferentiale Autocorretation_1m
	function
5. De.	
5. De. PART-B'	-function - 1 m
	-function - 1 m
PART-B!	function - 1 m Cross concelation - 1 m fine Narrow Based process - 2m.
PART-B'. 6. (a)	function Cross concelation - 1 m fine Narrow Based process - 2m. find constant c - 3m g (6m) Variance (x) - 3m.

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PART - A

1. State Baye's Theorem.

-Ans:

Baye's Theorem: If B, , B2, B3, --- Bn are n mutually exclusive and

Exhaustive events such that $P(B_i) > 0$ (i=1,2,3,...n) in a sample Space S and A is any other event in S intersecting every B' (i.e. A Can only occur in combination with any one of the events $B_i, B_2, B_3, ..., B_n$) Such that P(A) > 0 then the conditional protoability of B' given A is

2. what does the chebycher's in mequality states? Ans: Chebycher's Inequality:-

If x is a scandom variable with mean and variance r^2 then (i) $p \{ |x-\mu| \ge k\sigma \} \le \frac{1}{k^2}$ and (ii) $p \{ |x-\mu| < k\sigma \} \ge 1 - \frac{1}{k^2}$

Ans!

Auto Correlation :

Consider a random process x(t). Let x_1 , and x_2 be two random variables defined at times t_1 and t_2 with density function $f_x(x_1, x_2; t_1, t_2)$. The Correlation of x_1 and x_2 , $E[x_1, x_2] = E[x(t_1), x(t_2)]$ is auto correlation.

$$R_{x_{x}}(t_{1},t_{2}) = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} \alpha_{1} \alpha_{2} f_{x}(\alpha_{1},\alpha_{2};t_{1},t_{2}) d\alpha_{1} d\alpha_{2}$$



Cross Correlation:

Consider two random processes XII) and Y(t) defined with random variables X and Y at times t, and to respectively. Then the correlation of x and Y

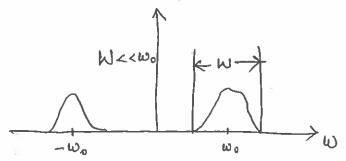
$$E[xy] = E[x(t_1) y(t_2)] \text{ is called the Cross}$$

Correlation function.
$$R_{xy}(t_1,t_2) = \int_{-\infty}^{\infty} xy f_{xy}(x,y;t_1,t_2) dxdy$$

Ans !

Narrow Band Process :

When the bandwidth 'W' of band limited process is Vory less than 'Wo' i.e if W< Wo, where wo 'y the frequency at which power spectrum is maximum, then it is said to be narrow band process.



PART-B

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Now find Variance of, or = G[x"] - (E[x]] -) () Now find $E[x^n] = \int_{-\infty}^{\infty} x^n f_x(x) dx$ = $\int x^{n} \cdot C x^{n} dx$ = c f xy da = 3 [357] $= \frac{3}{10} \left[\frac{1}{1} - \frac{1}{10} \right]^{5} = \frac{3}{10} \left[\frac{1}{2} \right]$ frm () $\therefore E[9^N] = \frac{3}{5}$ $\therefore Variance x' = \frac{3}{5} - 0^N$ $i - e = \frac{5}{5}$ Now find $P(X \neq b) = P(X \neq 0.5) = \int c x^{2} dx$ $= c \int dx dx$ = Z []]] ... $=\frac{3}{2}\left[1-(0.5)^{3}\right]$ $= \frac{1}{2} \left[1 - 0.125 \right]$ $\left[P(X \ge \frac{1}{2}) \stackrel{\text{P}}{=} 0.4375 \right]$

6) 5) Clabosiate on Discrete and Continuous Sample Spaces.
Ans: Discrete Sample Space:
A Sample Space is said to be descrete and
finite of the set
Eq: $S = d^{1}, 2, 3, 4j^{2}$ a finile Set
* If the set in the sample space have infinite
* If the set in the sample space is descrete and finite no-of elements then the sample space is descrete and finite
no-of elements then the sample sparse of
$e_{j!}$ $S = d_{1,2,3,-\cdots} j_{ij}$ infinite set

If the sample space contains infinite number of values with continuous values with a given range then it is called continuous lample space

Il S= { 0,1,2,3, --- 100} . The set of numbers Oto 100 the elements are enfinite and continuous.

N S RAJU INST ECHNOLOGY (AUTONOMOUS) SONTYAM , ANANDAPURAM, VISAKHAPATNAM - 531 173 7 (0) Let X be a continuous erandom variable with PDF fx(x) = { 423 , OCACI Find P(x==/x>1). Given $f_x(x) = \begin{cases} 4\pi^3, 0 \le \pi \le 1 \\ 0, 0 \text{ there is e} \end{cases}$ Ans ! $P(X = \frac{1}{2}/x > \frac{1}{3}) = \frac{P(x = \frac{1}{3}, x > \frac{1}{3})}{P(x - \frac{1}{3})}$ $P(X \leq \frac{2}{3}, X > \frac{1}{3}) = \int 4\pi^{3} dx$ $= 4 \int_{1/3}^{2/3} x^3 dx$ = 4 $\int_{1/3}^{1/3} x^3 dx$ = $4 \left[\frac{24}{4} \right]_{1/3}^{2/3}$ = (音)-(音) 0.189 - 0.0119 = 0.1771 $P(x>_{3}) = \int 4x^{3} dx = \# [\frac{x^{4}}{y}]_{y_{2}}^{y_{1}}$ 13 = [1-13] = [1 - 0.019]P(X>=)= 0.9881

$$P(X \leq \frac{2}{3} / X > \frac{1}{3}) = \frac{0.1771}{0.9881}$$

$$P(X \leq \frac{2}{3} / X > \frac{1}{3}) = 0.1779$$

F(6) Elaborate on conditional density and their properties. Ans! Conditional density function:

Conditional density function of a standom variable × as the derivative of the Conditional distribution function If we denote this density by $-f_{\rm X}(\pi/B)$, then $f_{\rm X}(\pi/B) = \frac{1}{4\pi} \left[F_{\rm X}(\pi/B) \right]$

If $F_{X}(\mathcal{A}_{\mathcal{B}})$ contains step discontinuities, as when X is a discrete or mixed random variable, we assume that impulse functions are present in $f_{X}(\mathcal{A}_{\mathcal{B}})$ to account for the derivatives at the discontinueties.

1)
$$f_{x}(x|B) \ge 0$$

2) $\int_{-\infty}^{\infty} f_{x}(y|B) dy = 1$
3) $F_{x}(x|B) = \int_{-\infty}^{\infty} f_{x}(x|B) dx$
4) $p(x_{1} \le x \le x_{2}|B) = \int_{x_{1}}^{x_{2}} f_{x}(x|B) dy$

N S RAJU INSTITU (AUTONOMOUS) SONTYAM , ANANDAPURAM, VISAKHAPATNAM – 531 173 8 a) State and Prove chebychev's Inequality Ans: Statement: If Kis a random variable with meanly and variance of then (i) P{12-11>koj<1 and (ii) P & 12-11 < K- 2 = 1- 1 Proof: - Given that Xis a random variable with mean 11 and variance or we know that $\sigma^{2} = V(x) = E [(x-u)^{2}]$:: |x-m] = Ko = $\int_{-\infty}^{\infty} (x-\mu)^2 f(x) dx$ X-MI ±Ko 2=4±KG $\sigma^{Y} = \int (\chi - \mu)^{\gamma} f(x) dx + \int (\chi - \mu)^{\gamma} f(x) dx + \int (\chi - \mu)^{\gamma} f(x) dx + \int (\chi - \mu)^{\gamma} f_{\chi}(x) dx$ $\chi = \mu - k\sigma$ $\chi = \mu - k\sigma$ $\mu - k\sigma$ $\mu + k\sigma$ $\chi = \mu - k\sigma$ $\mu + k\sigma$ $\chi = \mu - k\sigma$ $\mu - k\sigma$ Jhen $\sigma^{\gamma} = \int_{-\infty}^{M-K\sigma} (x-\mu)^{\gamma} f_{\chi}(x) dx + \int_{-\infty}^{\infty} (x-\mu)^{\gamma} f_{\chi}(x) dx + 0$ $M+K\sigma$ for the first integral in egn ()

$$\begin{split} \mathcal{X} &\leq \mu + k \sigma \\ \mathbf{x} \cdot \mathcal{U} &\leq -k \sigma^{-} \\ \mathbf{multiply} \cdot \mathbf{on} \ b \cdot \mathbf{s} \ b \mathbf{y}^{-1} \\ &-(\mathbf{x} \cdot \mathcal{U}) \geqslant k \sigma^{-} \\ \mathbf{Squaring} \ on \ both \ Sides \\ \left[-(\overline{\alpha} \cdot \mathcal{U})^{\gamma}\right] \Rightarrow (k \sigma)^{\gamma} \\ &(\overline{\alpha} - \mathcal{U})^{\gamma} \geqslant k^{\gamma} \sigma^{\gamma} \rightarrow (\overline{\mathbf{e}}) \\ \mathbf{for} \ dw \ second \ integral \ in \ eq^{-}(\mathbf{D}) \\ &\mathbf{q} \geqslant \mathcal{U} + k \sigma \\ &\mathbf{a} - \mathcal{U} \geqslant k \sigma \\ \mathbf{squaring} \ in \ b \cdot \mathbf{s} \\ &(\overline{\alpha} - \mathcal{U})^{\gamma} \geqslant \kappa^{\gamma} \sigma^{\gamma} \rightarrow (\overline{\mathbf{e}}) \\ \\ \text{Using equations } \quad (\overline{\mathbf{e}}) & \overline{\mathbf{e}} \otimes (\overline{\mathbf{e}}) \ eq^{n} \otimes \mathbf{e} \\ \sigma^{\gamma} \geqslant \int_{-\infty}^{\mathcal{U} - \mathcal{K}} \int_{-\infty}^{\infty} \frac{\mathcal{E}}{\mathcal{E}} (\mathbf{x}) d\mathbf{x} + \int_{-\infty}^{\infty} \mathcal{E}_{\mathbf{x}} (\mathbf{x}) d\mathbf{x} \\ &\sigma^{\gamma} \geqslant k^{\gamma} \sigma^{-} \frac{\mathcal{E}}{\mathcal{E}} (\mathbf{x}) d\mathbf{x} + \int_{-\infty}^{\infty} \mathcal{E}_{\mathbf{x}} (\mathbf{x}) d\mathbf{x} \\ \mathbf{u} + k \sigma \\ \sigma^{\gamma} \geqslant k^{\gamma} \sigma^{-} \left[\int_{-\infty}^{\mathcal{U} - \mathcal{K} \sigma} \frac{\mathcal{E}}{\mathcal{E}} (\mathbf{x}) d\mathbf{x} + \int_{-\infty}^{\infty} \frac{\mathcal{E}}{\mathcal{E}} (\mathbf{x}) d\mathbf{x} \\ \mathbf{u} + k \sigma \\ \mathbf{v} \geqslant k^{\gamma} \sigma^{-} \int_{\mathbf{x}} \frac{\mathcal{E}}{\mathcal{E}} (\mathbf{x}) d\mathbf{x} + \int_{-\infty}^{\infty} \frac{\mathcal{E}}{\mathcal{E}} (\mathbf{x}) d\mathbf{x} \\ \mathbf{u} + k \sigma \\ \mathbf{u} + k \sigma \\ \mathbf{u} = k^{\gamma} \left[\int_{-\infty}^{\mathcal{U} - \mathcal{K} \sigma} \frac{\mathcal{E}}{\mathcal{E}} (\mathbf{x}) d\mathbf{x} + \int_{-\infty}^{\infty} \frac{\mathcal{E}}{\mathcal{E}} (\mathbf{x}) d\mathbf{x} \\ \frac{\mathcal{E}}{\mathcal{E}} \geqslant \int_{-\infty}^{\mathcal{E}} \frac{\mathcal{E}}{\mathcal{E}} \mathbf{u} - \mathcal{E} = \mathbf{e} + \mathbf{e} \right] + \mathbf{p} \left[\mathbf{x} \ge \mathcal{U} + \mathbf{e} \sigma \right] \\ \frac{1}{k^{\gamma}} \geqslant \mathbf{p} \left[\mathbf{x} \le \mathcal{U} \le \mathbf{e} - \mathbf{e} - \mathbf{e} \right] + \mathbf{p} \left[\mathbf{x} \ge \mathcal{U} + \mathbf{e} \sigma \right] \end{aligned}$$

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 $\frac{1}{\nu\nu} \ge P\left[-(\alpha-\mu) \ge k\sigma\right] + P\left[\alpha-\mu \ge k\sigma\right]$ $\frac{1}{L^{2}} \geq P\left(|X-M| \geq k\sigma\right)$ $i = P \left[|x - \mu| \ge k - \right] \le \frac{1}{k^{\nu}} \rightarrow \oplus$ Now multiply 1-1 on both sider egn (+) - P [12-M] = Ko-] = - 1 adding 1 on bothsides 1- P [1x-1] > Ko] > 1- 1 i.e P[19-11< Ko-7 > 1- tr Hence proved.

8) b) If x be a random vareable and P(x=2) is the Probability mass function given as below table

	Х	0	1.1	2	3	୍ୟ	5	6	7	
	P(X=12)	0	ĸ	212	3ĸ	ų۴	k²	252	FRAIC	
Ļ	Find the	vali	e of	k a	und	P(XS	6)			

Ans!-

Given table

×	0	ſ	2	3	4	5	6	7	
P(x=n)	0	k	2 k	31	4 K	62	2Kr	ŦĸŦĸ	

we know that
$$\stackrel{\infty}{\underset{n=0}{\overset{\times}{\simeq}}} P(n) = 1$$

$$P(x=0) + p(x=1) + p(x=2) + p(x=3) + p(x=4) + p(x=5) + p(x=6) + p(x=7) = 1$$

0+k+2k+2k+3k+k+2k+ Ft+k=)

$$10k^{v} + 9k = 1$$

$$10k^{v} + 9k = 1$$

$$10k^{v} + 9k = 0$$

$$10k^{v} + 10k - 1k = 0$$

$$(10k + 1) - 1(k + 1) = 0$$

$$(10k + 1)(k + 1)(k + 1) = 0$$

$$(10k + 1)(k + 1)(k + 1) = 0$$

$$(10k + 1)(k + 1)(k + 1)(k + 1) = 0$$

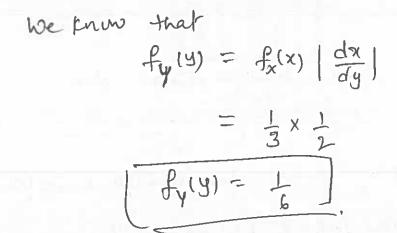
$$(10k + 1)(k + 1)(k + 1)(k + 1)(k + 1) = 0$$

$$(10k + 1)(k + 1)(k + 1)(k + 1)(k + 1)(k + 1) = 0$$

$$(10k + 1)(k + 1)(k$$

 $P(x \le b) = P(x = b) + P(x = 1) + P(x = 2) + P(x = 3) + P(x = 4)$ + P(x = 5) + P(x = 6) $= 0 + k + 2k + 2k + 3k + k^{2} + 2k^{2}$ $= 8(+ + 3k^{2})$ $= 8(+ + 3k^{2})$ = 8(+ + 3

N S RAJU INST TECHNOLOGY (AUTONOMOUS) SONTYAM , ANANDAPURAM, VISAKHAPATNAM - 531 173 9(a) Let Y = 2x+3 °f the random variable is uniformly distoubuted over [-1,2] determine fy1y). Ans! Riven y = 29+3 given interval (-1,2] a b 2 f random variable is uniformly distributed $f_{\chi}(\chi) = \begin{cases} \frac{1}{b-a}, & \alpha < \alpha < b \\ 0, & \text{otherwise} \end{cases}$ $= \int \frac{1}{a^2 - (-1)} - 1 < a < 2$ $f_{\chi}(\chi) = \begin{cases} \frac{1}{3}, -1 \leq \chi \leq 2 \\ 0, & \text{otherwise} \end{cases}$ The tranformation 4 = 2×+3 2x = y-3x = Y-3 It can be written of 2 = 4-3 diff writ my $\frac{dy}{dy} = \frac{1}{2} \frac{d}{dy} \left[\frac{y-3}{2} \right]$ $= \frac{1}{2} \left[1 - 0 \right]$ $\frac{dy}{dy} = \frac{1}{2} = \frac{1}{2} = \frac{1}{2} = \frac{1}{2}$



9(5) Explain about the characterstic function and state its properties.

Any! Characteristic function ?-Consider a sendom variable x costs a p.d.f $f_{X}(x)$, then the expected value of the function $e^{j\omega x}$ is called characteristic function . It is expressed as $f_{X}(\omega) = E [e^{j\omega x}]$ It is a function of a seal variable, $-\infty = \omega - \infty$ where j is an imaginary operator. $i = \phi_{X}(\omega) = E [e^{j\omega x}] = \leq e^{j\omega x_{i}} \rho(y_{i}) (for discrete)$ $\phi_{X}(\omega) = \int_{-\infty}^{\infty} e^{j\omega x} f_{X}(x) dx (for continuou)$

1. The characterstic function is Unity at
$$\omega = 0$$
 and given by
 $\varphi_{\chi}(\omega)|_{\omega=0} = \varphi_{\chi}(0) = 1$

The distantion function $F_{XY}(x,y)$ of a single orandom Variable (i.e either X or y) which is obtained by adjusting one of its value to imfinity is known as marginal distribution function.

The marginal distribution function of x and y are given by

$$F_{x}(x) = F_{xy}(x_{i,\infty}) = P(x \le x) = P(x \le x, y < \infty)$$

$$F_{y}(y) = F_{xy}(\infty, y) = P(y \le y) = P(x \le \infty, y \le y)$$

$$i \cdot e \quad f_{x}(x) = \underset{y}{\le} P(x = x, y = y)$$

$$F_{y}(y) = \underset{x}{\le} P(x = x, y = y)$$
for a continuous bivarente searcher wave given by

$$F_{x}(x) = \underset{x}{=} \underset{x}{\xrightarrow{y}} P(x = x, y = y)$$
for a continuous bivarente searcher by

$$F_{x}(x) = \int_{\infty}^{\infty} \int_{x_{y}}^{\infty} f_{x_{y}}(x, y) dx dy$$

$$F_{x}(x) = \int_{\infty}^{y} \int_{x_{y}}^{\infty} f_{x_{y}}(x, y) dx dy$$

$$F_{y}(y) = \int_{x}^{y} \int_{-\infty}^{\infty} f_{x_{y}}(x, y) dx dy$$

$$Margenal density function:$$

$$Morginal density function:$$

$$Morginal density function:
$$f_{x}(x) = \frac{1}{2\pi} \left[f_{x}(x) \right]$$

$$f_{y}(y) = \frac{1}{2\pi} \left[f_{x}(x) \right]$$

$$f_{y}(y) = \frac{1}{2\pi} \left[f_{x}(x) \right]$$

$$(ohne f_{x}(x) and f_{y}(y) are the marginal density functions
$$of x and y despectively$$

$$d can als o be expressed in terms of finit classity functions
$$f_{x}(x) = \int_{x}^{x} f_{xy}(x, y) dy - f_{y}(x, y) dy - f_{y}(y) = \int_{x}^{x} f_{xy}(x, y) dy - f_{xy}(x, y) dy - f_{y}(y) = \int_{x}^{x} f_{xy}(x, y) dy - f_{xy}(x, y) dy - f_{y}(y) = \int_{x}^{x} f_{xy}(x, y) dy - f_{xy}(x, y) dy - f_{y}(y) = \int_{x}^{x} f_{xy}(x, y) dy - f_{xy}(x, y) dy - f_{y}(y) = \int_{x}^{x} f_{xy}(x, y) dy - f_{y}(y) = \int_{x}^{x} f_{xy}(x, y) dy - f_{y}(x, y) dy - f_{y}(x, y) dy - f_{y}(x, y) dy - f_{y}(x, y) = \int_{x}^{x} f_{xy}(x, y) dy - f_{xy}(x, y) dy - f_{y}(x, y) dy - f$$$$$$$$

 $k \in$

Properties of Conditional distribution functions: 1. $F_x(-\infty/B) = 0$ 2. $F_x(\infty/B) = 1$ 3. $0 \le F_x(x/B) \le 1$ 4. $F_x(\sqrt{\pi/B}) \le F_x(\sqrt{\pi/B})$ if $\pi_i \le \pi_i$ 5. $P(\pi_i < x \le \pi_i/B) = F_x(\sqrt{\pi/B}) - F_x(\sqrt{\pi/B})$ 6. $F_x(x^{+}/B) = F_x(\sqrt{\pi/B})$

11a) Consider two signdom Varlables x and y with foint probability Mass function given in Table below. $\frac{1}{x=1} \quad \frac{1}{12} \quad \frac{1}{24} \quad \frac$

Find $P(X \leq 2, Y \leq 4)$, $P(x \leq 2/x = 1)$ and check x and y are independent

Ans '	Given	1	Y = 2	Y=4	Y=5
~	-1 ,	X=1	1/12	1 24	1 24
		×=2	1-6		tes
		x=3	4	18	12

We know that

$$F_{X,Y}(x_1,y) = P(x \le n, p \le q)$$

 $f_{X,Y}(x_1,y) = P(x \le n, p \le q)$
 $f_{X,Y}(x_1,y) = P(x \le n, y \le q)$
 $f_{X,Y}(x_1,y) = P(x \le n, y \le q)$
 $f_{X,Y}(x_1,y) = P(x \le n, y \le q)$
 $f_{Y,Y}(x_1,y) = P(x \ge n, y \le q)$
 $f_{Y,Y}(x_1,y)$

- 11) b) State and prove central limit theorem for equal distributions.
 - Ans: <u>Central limit theorem</u>: According to the central limit theorem (CLT), the distonbution of random process which is cumulative effect of a large number of independent noise sources can be assumed to be Gaussian.

For example in communication systems, the noise is always modelled as a random variable with gauseian distribution. This is a valied assumption since the noise in a communication System is the cumulative effect of many random noise Sources.

Equal distributions;
Consider N continuous standom variables
$$x_n$$
, $n=1,2,3,-\infty$
have the same distribution and density furtions.
Let $Y = x_1 + x_2 + x_3 + \cdots + x_N$

Also Let w be normalized and on variable $i-e w = \frac{Y-\overline{Y}}{\overline{y}}$

where $Y = \underbrace{X}_{n=1}^{Y} x_n$, $\overline{Y} = \underbrace{X}_{n}$ and $\overrightarrow{Y} = \underbrace{X}_{n=1}^{Y} x_n$

So
$$M = \sum_{n=1}^{N} x_n - \sum_{n=1}^{N} \overline{x_n}$$

 $\left(\sum_{n=1}^{N} \overline{v_n}\right)^{1/2}$

Since all random variables have the same distribution

Then wis a gaussian random Variables.

12) a) Explain about second order and winde sense Stationary process:

And: Second order or Wide Sense Stationary process's-

If the second order probability density function
of trandom process is Endependent of time (ie does not charge
with time) then it is called as second order stationary
process
i.e.
$$f_x(x_1, x_2; t_1, t_2) = f_x(x_1, x_2; t_1, t_1, t_2, t_2)$$

Where, 7 is a real number

Since second order probability density function determines the first order density function, a second order stationary process consists of first order stationary process.

The autocorrelation function of the scandom process X (t) for t, and to is given as,

$$R_{X_{K}}(t_{1}, t_{2}) = E \left[x(t_{1}) x(t_{2}) \right]$$

If the auto correlation function of the orandom process is the function of time difference $(T'=t_q-t_1)$ but not absolute time, then it is called as second order (N) wide lenge stationary

process (-wss)
i-c
$$R_{xx}(t_i, t_i+T') = E\left[x(t_i) \cdot x(t_i+T')\right]$$

 $\therefore R_{xx}(t_i, t_i+T') = R_{xx}(T')$

A random process is said to be wide sense stationary or weak stationary if it satisfies the following conditions,

1. The mean of the vandom process is constant i.e. $E[x_{it}] = \overline{x} = Constant$.

 \mathcal{Q} . Its cuto correlation functions depends only on $\mathcal{P}(i \in t_2 - t_1)$ but not on t. $i \in \mathbb{E}[x(t) \times (t_1 \mathcal{P})] = \mathcal{R}_{X_X}(\mathcal{P})$ 12) b) Derive the Relationship between Cross-power density spectrum and Cross - asselation function.

$$\begin{split} & \underbrace{\operatorname{free}}_{Y_{T}}^{T}(t) \text{ and } Y_{T}(t) \text{ are ensemble members of the} \\ & praces X(t) \text{ and } Y(t) \text{ suspectively} \cdot \operatorname{free}_{t} \text{ forware bars}(f^{t}) \\ & of X_{T}(t) \text{ is given by}, \\ & X_{T}(w) = \int^{T} X(t) e^{-jwt} dt \quad \rightarrow \mathbb{O} \\ & \operatorname{free}_{t} \text{ forware brans}(f^{t}) \text{ of } Y_{T}(t) e^{-jwt} dt, \quad \rightarrow \mathbb{O} \\ & \operatorname{free}_{t} \text{ forware brans}(f^{t}) \text{ of } Y_{T}(t) e^{-jwt} dt, \quad \rightarrow \mathbb{O} \\ & \operatorname{free}_{t} \text{ forware brans}(f^{t}) \text{ of } Y_{T}(t) e^{-jwt} dt, \quad \rightarrow \mathbb{O} \\ & \operatorname{free}_{t} \text{ forware transform of } X_{T}^{*}(t) \text{ is given by} \\ & = \int_{-T}^{T} Y(t_{T}) e^{-jwt} dt, \quad \rightarrow \mathbb{O} \\ & \operatorname{free}_{t} \text{ forware transform of } X_{T}^{*}(t) e^{jwt} dt \quad \rightarrow \mathbb{O} \\ & \operatorname{free}_{t} \text{ forware transform of } X_{T}^{*}(t) e^{jwt} dt \quad \rightarrow \mathbb{O} \\ & \operatorname{free}_{t} \text{ forware transform of } X_{T}^{*}(t) e^{jwt} dt \quad \int \mathbb{O} \\ & \operatorname{free}_{t} \text{ forware transform of } X_{T}^{*}(t) e^{jwt} dt \quad \int \mathbb{O} \\ & \operatorname{free}_{t} \text{ forware transform of } X_{T}^{*}(t) e^{jwt} dt \quad \int \mathbb{O} \\ & \operatorname{free}_{t} \text{ forware transform of } Y_{T}(w) = \int_{-T}^{T} X(t) e^{jwt} dt \quad \int -T \quad \rightarrow \mathbb{O} \\ & \operatorname{free}_{t} \text{ forware transform of } Y_{T}(w) = \int_{-T}^{T} X(t) e^{jwt} dt \quad \int -T \quad \rightarrow \mathbb{O} \\ & \operatorname{free have} \\ & \operatorname{$$

Now applying Enverse transform on both side, we get,

$$\frac{1}{2T} \int S_{xy}(w) e^{jwT} dw = \frac{1}{2\pi} \int L_{T} \int J_{T} \int J_{xy}(t,t_{1}) e^{jw(t_{T}t)} dt dt_{1}$$

$$= L_{T} \int J_{T} \int J_{T} \int J_{T} \int R_{xy}(t,t_{1}) \int J_{T} \int e^{jw(T-t_{1}+t)} dw dt dt_{1}$$

$$= L_{T} \int J_{T} \int R_{xy}(t,t_{1}) \int J_{T} \int e^{jw(T-t_{1}+t)} dw dt dt_{1}$$

$$= L_{T} \int J_{T} \int R_{xy}(t,t_{1}) \int J_{T} \int e^{jw(T-t_{1}+t)} dw dt dt_{1}$$

$$= d(t_{1}-T-t) dt_{1} dt \begin{bmatrix} \cdots & J_{T} & J_{T} & J_{T} \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$= L_{T} \int J_{T} R_{xy}(t,t_{T}+T) dt_{1} dt$$

$$= J_{T} \int R_{xy}(t,t_{T}+T) dt_{1} dt$$

. The ocelation becomes

$$\frac{1}{2\pi} \int S_{xy}(w) e^{9wr} dw = dt \quad \frac{1}{2T} \int R_{xy}(t_i t_i t_i \tau) dt$$

$$T \to \infty \qquad T \to \infty \qquad -T$$

She above 'relations is valed for $-T < t_i \tau < T$.

(OR) (I3) a) Explain how random processes are classified with neat Sketches.

Ant!- Deffinition of Random Perocess 5-
A standom Variable cohich is a function of sample
Space and time is called standom process
$$\{x(t,s)\}$$
.
It can be represented as a time function over the
entire time i.e $[x(t,s) = x(t)]$.

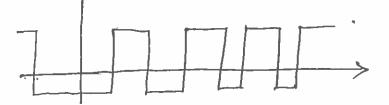
Classification of Random process:-

Random processes are mainly classified into four types based on time it' and amplitude of trandom Variable 'X'. 1. Continuous Random process (CRP) 2. Discrete Random process (DRP) 3. Continuous Random Sequence (CRS) 4. Discrete Random Sequence (DRS) 1. Continuous Random process :-

A random process is said to be continuous if trandom variable X and lime are continuous over the entire time (12H)

2. Discrete Random process :-

In déscrete random process, the random vareable × has only déscreté value white time t'és continuous



c) Continuous siandom process &-

In Continuous sandom sequence, erandom variable X is Continuous but time t has discrete values. although the the salues of the

d) Discrete Random Sequence_

13) b) Derive the relationship between power spectral density and auto Correlation.

Ans! Relationship between power spectral density and auto covelation .

Let X(t) be an ensemble of anadom process X(t). Let us define $X_{T}(t)$ as $X_{T}(t) = \int_{0}^{\infty} x(t) - for - T \le t \le T$ $\int_{0}^{\infty} 0 \text{ herwise}$ The fourier transform of $X_{T}(t)$ is given by $X_{T}(t) = \int_{-T}^{T} x_{T}(t) e^{-Swt} dt$

Using parseval's thebur, we can work
$$\int_{-T}^{T} x^{T}(t) dt = \frac{1}{2\pi} \int_{-\infty}^{\infty} 1x_{T}(t) dt$$
The average power is given by
$$P(T) = \frac{1}{2\pi} \int_{-T}^{T} x^{T}(t) dt = \frac{1}{2\pi} \int_{-\infty}^{\infty} \frac{1x_{T}(t)}{2\pi} dt$$
To find the average power of the senders process, we take the expected value with T tending to infiniting in the above equation:
$$P_{XX} = \frac{1}{2\pi} \int_{-\infty}^{\infty} \frac{1}{2\pi} \int_{-T}^{T} E[x^{n}(t)] dt$$

$$= \frac{1}{2\pi} \int_{-\infty}^{\infty} \frac{1}{2\pi} \int_{-T}^{T} E[x^{n}(t)] dt$$

$$We \quad the above equation, we can obtain
$$S_{XX}(w) = \frac{1}{2\pi} \int_{-T}^{T} \int_{-T}^{T} E[x^{n}(t)] e^{\frac{1}{2}wt} dt$$

$$S_{XX}(w) = \frac{1}{2\pi} \int_{-T}^{T} \int_{-T}^{T} E[x^{n}(t)] e^{\frac{1}{2}wt} dt$$

$$V(t) = \frac{1}{2\pi} \int_{-T}^{T} \int_{-T}^{T} E[x^{n}(t)] e^{\frac{1}{2}wt} dt$$

$$S_{XX}(w) = \frac{1}{2\pi} \int_{-T}^{T} \int_{-T}^{T} e^{\frac{1}{2}} e^{\frac{1}{2}wt} dt$$

$$S_{XX}(w) = \frac{1}{2\pi} \int_{-T}^{T} \int_{-T}^{T} e^{\frac{1}{2}} e^{\frac{1}{2}wt} dt$$$$

Now the involue formier transform of
$$S_{XX}(\omega)$$
 is

$$F^{-1}[S_{XX}(\omega)] = \frac{1}{2\pi T} \int_{-\infty}^{\infty} S_{XX}(\omega) e^{-j\omega T} d\omega$$

$$= \frac{1}{2\pi T} \int_{-\infty}^{\infty} \frac{1}{2\pi T} \int_{-T}^{T} \int_{-T}^{T} R_{XX}(t_{1}, t_{2}) e^{-j\omega T} d\omega dt_{2}dt_{1}$$

$$= \frac{1}{2\pi T} \int_{-T}^{T} \int_{-T}^{T} R_{XX}(t_{1}, t_{2}) \frac{1}{2\pi} \int_{-\infty}^{\infty} e^{\frac{1}{2}\omega(t_{2}-t_{1}+T)} d\omega dt_{2}dt_{1}$$

$$= \frac{1}{2\pi T} \int_{-T}^{T} \int_{-T}^{T} R_{XX}(t_{1}, t_{2}) \frac{1}{2\pi} \left[a\pi d(t_{2}-t_{1}-T) d\omega dt_{2}dt_{1} \right]$$

$$= \frac{1}{2\pi T} \int_{-T}^{T} R_{XX}(t_{1}, t_{2}) \frac{1}{2\pi} \left[a\pi d(t_{2}-t_{1}-T) d\omega dt_{2}dt_{1} \right]$$

$$= \frac{1}{2\pi T} \int_{-T}^{T} R_{XX}(t_{1}, t_{2}) \frac{1}{2\pi} \left[a\pi d(t_{2}-t_{1}-T) d\omega dt_{2}dt_{1} \right]$$

$$= \frac{1}{2\pi T} \int_{-T}^{T} R_{XX}(t_{1}, t_{2}) \frac{1}{2\pi} \left[a\pi d(t_{2}-t_{1}-T) d\omega dt_{2}dt_{1} \right]$$

$$= \frac{1}{2\pi T} \int_{-T}^{T} R_{XX}(t_{1}, t_{2}) \frac{1}{2\pi T} \left[a\pi d(t_{2}-t_{1}-T) d\omega dt_{2}dt_{1} \right]$$

$$= \frac{1}{2\pi T} \int_{-T}^{T} R_{XX}(t_{1}, t_{2}) \frac{1}{2\pi T} \left[a\pi d(t_{2}-t_{1}-T) d\omega dt_{2}dt_{1} \right]$$

$$= \frac{1}{2\pi T} \int_{-T}^{T} R_{XX}(t_{1}, t_{2}) \frac{1}{2\pi T} \left[a\pi d(t_{2}-t_{1}-T) d\omega dt_{2}dt_{1} \right]$$

$$= \frac{1}{2\pi T} \int_{-T}^{T} R_{XX}(t_{1}, t_{2}) \frac{1}{2\pi T} \left[a\pi d(t_{2}-t_{1}-T) d\omega dt_{2}dt_{1} \right]$$

$$= \frac{1}{2\pi T} \int_{-T}^{T} R_{XX}(t_{1}, t_{2}) \frac{1}{2\pi T} \int_{-T}^{T} R_{XX}(t_{1}, t_{2}) \frac{1}{2\pi T} \left[a\pi d(t_{2}-t_{1}-T) d\omega dt_{2}dt_{1} \right]$$

$$= \frac{1}{2\pi T} \int_{-T}^{T} R_{XX}(t_{1}, t_{1}) \frac{1}{2\pi T} dt_{1}$$

$$= \frac{1}{2\pi T} \int_{-T}^{T} R_{XX}(t_{1}, t_{2}) \frac{1}{2\pi T} \frac$$

1 a 1

.

Therefore

$$S_{XX}(w) = \int_{-\infty}^{\infty} R_{XX}(t) e^{-f_W t} dt \rightarrow 0$$
and

$$R_{XX}(t) = \int_{0}^{\infty} \int_{-\infty}^{\infty} S_{XX}(w) e^{f_W t} dw \rightarrow 0$$
i-e power spectral density & fourier transform of
euto correlation function and eqn $0 \ge 0$ is also known as
where - Khintchine orelations.

14) a) A random perocesses X(t) = A sin(wt+0), where A, ware Constants and O is Uniformly distributed random rbriable on the Enterval (-17, 71). find average power?

Ans:-

$$P_{XK}(\tau^{2}) = E[X^{2}t^{2}t^{2}]$$

$$= \int_{-\infty}^{\infty} x^{2}(t) f_{p}[\theta] d\theta$$

$$= \int_{-\pi}^{\pi} \left[k \sin(\omega t + \theta)\right]^{2} \cdot \frac{1}{4\pi} d\theta$$

$$= \frac{1}{4\pi} \int_{-\pi}^{\pi} A^{2} \sin^{2}(\omega t + \theta) d\theta$$

$$= \frac{A^{2}}{4\pi} \int_{-\pi}^{\pi} \left[\frac{1 - \cos(2\omega t + 2\theta)}{2}\right]^{2} d\theta$$

$$= \frac{A^{2}}{4\pi} \int_{-\pi}^{\pi} \left[0 - \frac{\sin(2\omega t + 2\theta)}{2}\right]_{-\pi}^{\pi}$$

$$= \frac{A^{2}}{4\pi} \left[w[\theta]_{-\pi}^{\pi} - \frac{1}{2}\left[\sin(2\omega t + 2\theta)\right]_{-\pi}^{\pi}\right]$$

$$= \frac{A^{2}}{4\pi} \left[v[\theta]_{-\pi}^{\pi} - \frac{1}{2}\left[\sin(2\omega t + 2\theta)\right]_{-\pi}^{\pi}\right]$$

$$= \frac{A^{2}}{4\pi} \left[x[\theta]_{-\pi}^{\pi} - \frac{1}{2}\left[\sin(2\omega t + 2\theta)\right]_{-\pi}^{\pi}\right]$$

$$= \frac{A^{2}}{4\pi} \left[x[\theta]_{-\pi}^{\pi} - \frac{1}{2}\left[\sin(2\omega t + 2\theta) + \sin(2\omega t - 2\pi)\right]\right]$$

$$= \frac{A^{2}}{4\pi} \left[2\pi - \frac{1}{2}\left[\sin(2\omega t - 2\pi) + \sin(2\omega t - 2\pi)\right]\right]$$

$$= \frac{A^{2}}{4\pi} \left[2\pi - \frac{1}{2}\left[\sin(2\omega t - \sin(2\omega t) + \sin(2\pi - 2\omega t)\right]\right]$$

$$= \frac{A^{2}}{4\pi} \left[2\pi - \frac{1}{2}\left[\sin(2\omega t - \sin(2\omega t) + \sin(2\pi - 2\omega t)\right]\right]$$

$$= \frac{A^{2}}{4\pi} \left[2\pi - \frac{1}{2}\left[\sin(2\omega t - \sin(2\omega t) + \sin(2\pi - 2\omega t)\right]\right]$$

$$= \frac{A^{2}}{4\pi} \left[2\pi - \frac{1}{2}\left[\sin(2\omega t - \sin(2\omega t) + \sin(2\pi - 2\omega t)\right]\right]$$

$$= \frac{A^{2}}{4\pi} \left[2\pi - \frac{1}{2}\left[\sin(2\omega t - \sin(2\omega t) + \sin(2\pi - 2\omega t)\right]\right]$$

$$= \frac{A^{2}}{4\pi} \left[2\pi - \frac{1}{2}\left[\sin(2\omega t - \sin(2\omega t) + \sin(2\pi - 2\omega t)\right]\right]$$

$$= \frac{A^{2}}{4\pi} \left[2\pi - \frac{1}{2}\left[\sin(2\omega t - \sin(2\omega t) + \sin(2\pi - 2\omega t)\right]\right]$$

1416) Derive the relationship between input PSD and output PSD of an LTE System. Power Density spectrum of Response ;-Ans! Consider that a random process X(t) is applied on an LTI system having a transfor furtion H(w). The output response is 41+). If the power spectrum of the input is Ixx(w), then thun power spectrum of the output susponse is given by $S_{yy}(\omega) = |H(\omega)|^{\gamma} S_{xx}(\omega)$ Poroof: Let RXX(T) be the autocorrelation of the output response y It). then the power spectrum of the response is the fourier tranform of Ryy (7). \therefore Syy(w) = F[Ryy(r)]= Joo Ryy(T) = Swid T. We know that $R_{yy}(t') = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} R_{xx}(t'+t'_1-t'_2) Rh(t'_1) h(t'_2) dt'_1 dt'_2$ $Jhen (S_{1}(W)) = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} R_{1}(T+Y_{1}-Y_{2})h(T)h(T_{2})dT, dT_{2} = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} R_{1}(T+Y_{1}-Y_{2})h(T_{2})h(T_{2})dT, dT_{2} = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} R_{1}(T+Y_{1}-Y_{2})h(T+Y_$ $= \int_{-\infty}^{\infty} h(T_{i}) \int_{-\infty}^{\infty} R_{X_{x}}(T_{i}+T_{i}-T_{y}) e^{j\omega T} d\tau d\tau_{y} d\tau_{y} d\tau_{y}$

$$= \int_{-\infty}^{\infty} h(T_1) e^{j(w)T_1} \int_{-\infty}^{\infty} h(T_2) e^{j(w)T_2} \int_{-\infty}^{\infty} R_{xx} (T+T_1-T_2) e^{j(w)T_1} e^{j(w)T_1} dT dT_1 dT_2$$

$$ket \quad T+T_{1}-T_{2}=t, \ dT = dt$$

$$S_{yy}(\omega) = \int R(T_{1}) e^{j\omega T_{1}} dT_{2} \int R(T_{2}) e^{j\omega T_{2}} \int Rxx(t) e^{j\omega T_{2}} dt$$

We down that

$$S_{yy}(w) = H^{f}(w) H(w) S_{xx}(w) = H(-w) H(w) S_{xx}(w)$$

 $S_{yy}(w) = (H(w))^{r} S_{xx}(w)$

(OR) 15) a) Explain the following (i) Noise Figure (ii) Noise Saurces.

Ans!

Noise figure ?-
Noise figure give the amount of noise internally
generated by the seystem. It is the state of the power
densily of the total noise available at the output of the
network to the power density available at the output only
due to the input noise Source. Noise figure gives a
measure of the system performance of the noise.
It is mathematically expressed as

$$F = \frac{S_{n_0}(\omega)}{S_{n_0}(\omega)} = \frac{S_{n_0}(\omega) + S_{n_0}^{"}(\omega)}{S_{n_0}(\omega)} = 1 + \frac{S_{n_0}(\omega)}{S_{n_0}(\omega)}$$

Where Snolwi = the total noise power spectral density at the output Sho(00) is the noise power Spectral density at the output due to Enput noise and Sno (w) = noise power spectral density at the output due to the noise generated internally by the system. and Sno(w) = sno(w). then F = 1Not: If F>1, the System & said to be a moisy system. These mange of Fis 1 < FLOD. As Fincreases, the System becomes noisy. (11) Notse Sources :-There are two types of noise sources 1. External Moise 2. Internal noise. External Noises-Noise whose sources are external to the preceiver is called external noise. Most external noise is added Ento the desired signal is communication channels, In external noise having (i) Atmospheric Noise (ii) Extraterestatal Norse

(117) Industrial Noise

Internal Noise :-

The noise Created with in a other device or a system is called internal noise. -) Internal noise generated by any of the active (or) passive devices found in systems. This noise is also called function noise Internal noise harring -> (i) Shot noise (ii) Teansit - time noise (iii) Flicker Noise (iv) Thermal Noise ..

15) b) Derive the expression for average cross power between two scendom process XIt) and YIt)

$$\frac{4}{2} \frac{1}{2} \frac{1}$$

Consider two seal scandom process Xits and Yits. If
X(t) and Yits) are fointly wass scandom process them the
Cross power density spectrum & defined as F.T of the cross
Corelation function of x(t) and Y(t)

$$S_{XY}(w) = \frac{1}{T \to \infty} \frac{E\left[1 \times \frac{1}{T}(w) \times \frac{1}{T}(w)\right]}{2T}$$

 $S_{YK'(w)} = \frac{1}{T \to \infty} \frac{E\left[1 \times \frac{1}{T}(w) \times \frac{1}{T}(w)\right]}{2T}$

Average Cross power:
Jhe average cross power of a wide serve
Stationary examplos process with is the any power is
the Cross consolution function at
$$P=D$$

 $P_{xy} = R_{xy}(P) | T=D$
 $= E [x(t) Y(t+P)]/T=D$
 $P_{xy} = \frac{1}{2T} \int_{-\infty}^{\infty} S_{xy}(w)dw$
(M) Time average if and moment.
 $P_{xy} = A [E E[x(t)] Y(t+J]]$
 $= \frac{L}{T-T} \int_{-T}^{T} E [x(t) Y(t)] dr$
 $P_{xy} = \frac{L}{T-T} \int_{-T}^{T} R_{xx}(b) dr$

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n Deg	Semester End Supplementary Examination, A	pril/May	, 2022	
-	Irea Codo DOCCODA		Academic Year	2020 - 2021
Cou		s 70	Semester	III II
Pari	A (Short Answer Questions 5 x 2 = 10 Marks)		-	
NC	Questions (1 through 5)		Learning Outcor	na (s) Del
1	- inventious formatica and anticipatione line formatione		20CS304,1	
3			20CS304.2	L1
4	Define Pointers.		20CS304.3	
5	The and the differences between vectors & lists		20CS304.4 20CS304.5	
No	B (Long Answer Questions 5 x 12 = 60 Marks) Questions (6 through 15)		2000001.0	() L I
		Marks	Learning Outcon	ne (s) DoK
6 (a	r rogramming and object-offented programming	6M	20CS304.1	L1 ⁼
6 (t	Explain the advantages of object-oriented programming.	6M	20CS304.1	L2
7 (a	OR Explain the different types of Data types with example.	-	81.63 B	
i 7 (b	What is code reusability? Explain different C++ features that	7M	20CS304.1	L2
. (-	enable reusability.	5M	20CS304.1	L1
1 8 (a		414	0000004.0	25 25 30 100000
8 (b	what is dynamic binding? How it is different from static binding?	4M	· 20CS304.2	
	List some advantages of dynamic binding over static binding.	8M	20CS304.2	L1
9 (a		4M	0000001.0	
9 (b	Define inline function. Write a C++ program for finding the grap of a		20CS304.2	
•	triangle using inline functions.	8M	20CS304.2	L1
10 (a		6M	20CS304.3	С. т.а. ¹ 1
10 (t	Write a C++ program to find the area of circle rectangle and			→ ·
SI 12	triangle using function overloading.	6M -	20CS304.3	L3
11 (a	What is inheritance? Present the advantages and disadvantages of			
11 (b	innentance.	6M	20CS304.3	L1
11 (6) Write a C++ Program to overload + operator to add two matrices.	6M	20CS304.3	L3
12 (a	Explain the role of this pointer in C++ with a programming			
12 (b	example.	5M	20CS304.4	L2
	OR	7M	20CS304.4	L3
13 (a	What is a virtual destructor? Explain with an example	6M	20CS304.4	L2
13 (b)	Write a C++ program to illustrate catching all exceptions.	6M	20CS304.4	L3
14 (a)		6M	2000204 5	
14(b)	Write a function template for finding the minimum value in an array.	6M	20CS304.5 20CS304.5	L2 L3
	OR			70
15 (a)	Define template. What is the need for templates in programming? Write C++ code that declares a Template class.	5M	20CS304.5	L2
15 (b)	Write a C++ program that fills a vector with random numbers.	7M		
			20CS304.5	L3
	Co	-i) ntrollo	Kalyani	tions

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N S RAJU INSTITUTE OF TECHNOLOGY

(AUTONOMOUS)

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SCHEME OF VALUATION & ANSWER KEY

Degree	B. Tech. (U. G.)	Program	CSE/CSM/CSD	Test	end Exam	Academic Year	2021 - 2022
Course Code	20CS304	Test Duration	180 Min.	Max. Marks	70	Semester	III
Course Assessment Pa		- •	rough c++				
R (L1):	U (L2):	Apply (L3):	Analyze (L4):		E (L5):	C (Le	5)

Part A (Short Answer Questions 5 x 2 = 10 Marks)

No.	Questions (1 through 5)		Learning Outcome (s)		DoK
1	Differentiate formatted and unformatted console I/O operations.	2M	20CS304.1	L1	
	Unformatted input /output is the most basic form of input/output.				

Unformatted I/O transfers the internal binary representation of the data directly between memory and the file.

Formatted output converts the internal binary representation of the data to Ascii characters which are written to the output file.

Formatted input reads characters from the input file and converts them to internal form.

Unformatted I/O operations:cout,cin,put(),get(),getline(),write()

Formatted I/O operations :ios class functions and flags,manipulators,user-defined output functions

2

Define class and object

20CS304.2 L1

- A Class is a user defined data-type which has data members and member functions.
- Data members are the data variables and member functions are the functions used to manipulate

these variables and together these data members and member functions defines the properties and behavior of the objects in a Class.

An **Object** is an instance of a Class. When a class is defined, no memory is allocated but when it is instantiated (i.e. an object is created) memory is allocated.

List any four types of inheritance Single inheritance Multiple inheritance Hierarchical Inheritance Hybrid Inheritance Multilevel inheritance

Define Pointers?

4

5

3

2m 20CS304.4

2m

20CS304.3

L1

L1

A pointer is a variable that stores the memory address of an object. Pointers are used extensively in both C and C++ for three main purposes: to allocate new objects on the heap, to pass functions to other functions, to iterate over elements in arrays or other data structures.

Vector	List
t has contiguous memory.	While it has non-contiguous memory.
t is synchronized.	While it is not synchronized.
Vector may have a default size.	List does not have default size.
In vector, each element only requires the space for itself only.	In list, each element requires extra space for the node which holds the element, including pointers to the nex and previous elements in the list.
Insertion at the end requires constant time but insertion elsewhere is costly.	Insertion is cheap no matter where in the list it occurs
Vector is thread safe.	List is not thread safe.

Part B (Long Answer Questions 5 x 12 = 60 Marks)

No.	Questions (6 through 15)		Learning Outcome (s)	DoK
6 (a)	Write difference between procedural oriented probject –oriented programming ? 6m procedure or oriented-3m	rogramming and iented-3m object	20CS304.1	L1
	Procedural Oriented Programming	Object-Oriented	Programming	
	In procedural programming, the program is divided into small parts called <i>functions</i> .	In object-oriented divided into smal		
	Procedural programming follows a top-down approach.	Object-oriented p <i>up approach</i> .	programming foll	ows a <i>bottom-</i>
	There is no access specifier in procedural programming.	Object-oriented p specifiers like pri		
	Adding new data and functions is not easy.	Adding new data	and function is	easy.
	Procedural programming does not have any proper way of hiding data so it is <i>less secure</i> .	Object-oriented p hiding so it is mo		vides data

6 (b) Explain the advantages of object oriented programming 6M

20ESX02.1 L2

- We can build the programs from standard working modules that communicate with one another, rather than having to start writing the code from scratch which leads to saving of development time and higher productivity,
- OOP language allows to break the program into the bit-sized problems that can be solved easily (one object at a time).
- The new technology promises greater programmer productivity, better quality of software and lesser maintenance cost.
- OOP systems can be easily upgraded from small to large systems.
- It is possible that multiple instances of objects co-exist without any interference,
- It is very easy to partition the work in a project based on objects.

7 (a) Explain the different types of data types with example. 7M Datatypes-4m examples-3m

20CS304.1 L2

3

Data types in C++ is mainly divided into three types

1. **Primitive Data Types**: These data types are built-in or predefined data types and can be used directly by the user to declare variables. example: int, char, float, bool etc. Primitive data types available in C++ are:

- Integer
- Character
- Boolean
- Floating Point
- Double Floating Point
- Valueless or Void
- Wide Character

2Derived Data Types: The data-types that are derived from the primitive or built-in datatypes are referred to as Derived Data Types. These can be of four types namely:

- Function
- Array
- Pointer
- Reference

3 Abstract or User-Defined Data Types: These data types are defined by user itself. Like, defining a class in C++ or a structure. C++ provides the following user-defined datatypes:

- Class
- Structure
- Union
- Enumeration
- Typedef defined DataType

using namespace std;

```
int main()
```

```
cout << "Size of char : " << sizeof(char)
   << " byte" << endl:
  cout << "Size of int : " << sizeof(int)
   << " bytes" << endl;
  cout << "Size of short int : " << sizeof(short int)
   << " bytes" << endl;
  cout << "Size of long int : " << sizeof(long int)
    << " bytes" << endl;
                       of signed
                                        long int
  cout << "Size
<<sizeof(signedlongint)
    << " bytes" << endl;
  cout << "Size of unsigned long int :
                                                        <<
sizeof(unsigned long int)
    << " bytes" << endl;
  cout << "Size of float : " << sizeof(float)
    << " bytes" <<endl;
  cout << "Size of double : " << sizeof(double)</pre>
    << " bytes" << endl;
  cout << "Size of wchar_t : " << sizeof(wchar_t)
    << " bytes" <<endl;
```

return 0;

7 (b) What is code reusability ?explain different c++ features that enable 20ESX02.3 L3 reusability? -5m

- C++ strongly supports the concept of reusability. The C++ classes can be reused in several ways. Once a class has been written and tested, it can be adapted by another programmer to suit their requirements. This is basically done by creating new classes, reusing the properties of the existing ones.
- software re-usability is primary attribute of software quality. C++ strongly supports the concept of reusability. C++ features such as classes, virtual function, and templates allow designs to be expressed so that re-use is made easier there are many advantage of reusability. They can be applied to reduce cost, effort and time of software development. It also increases the productivity, portability and reliability of the software product

8 (a) With an example explain the syntax for declaring objects 20CS304.2 L2 syntax:2m example 2m 4MAn Object is an instance of a Class. When a class is defined, no memory is allocated but when it is

instantiated (i.e. an object is created) memory is allocated.

Svntax:

ClassName ObjectName;

In C++, an object is created from a class. We have already created the class named MyClass, so now we can use this to create objects.

To create an object of MyClass, specify the class name, followed by the object name.

To access the class attributes (myNum and myString), use the dot syntax (.) on the object:

class MyClass { // The class

public: // Access specifier

int myNum; // Attribute (int variable)

string myString; // Attribute (string variable)

```
};
```

int main() {

MyClass myObj; // Create an object of MyClass

// Access attributes and set values

myObj.myNum = 15;

```
myObj.myString = "Some text";
```

// Print attribute values
cout << myObj.myNum << "\n";
cout << myObj.myString;
return 0;</pre>

8 (b) What is Dynamic binding ?how it is different from static binding?list some advantages of dynamic binding over static binding? 8m DYNAMIC BINDING-2M DIFFERENCE-2M ,Advantages-4m

Dynamic binding refers to linking a procedure call to code that will execute only once. The code associated with the procedure is not known until the program is executed, which is also known as late binding.

1. Static binding happens when all information needed to call a function is available at the compiletime. Dynamic binding happens when the compiler cannot determine all information needed for a function call at compile-time.

Advantages:

1. Static Binding execution is more efficient & faster than Dynamic. This Binding compiler knows that

these types of methods can not be overridden.

- 2. In the Static Binding, the type is used while Dynamic Binding uses objects for Bindings.
- 3. One of the major advantages of Dynamic Binding is flexibility; due to the flexibility, a single function

can handle different types of an object at runtime.

4. In Static Binding, All information needed before the compilation time, while in Dynamic Binding, no information remains available before run time.

5. Static Binding can take place using normal functions, while Dynamic Binding can be achieved using virtual functions.

OR

Explanation-1m,syntax-2m ,example-1m Explain Friend Function ? 20CS304.2 L2 4m

A friend function in C++ is defined as a function that can access private, protected and public members of a class.

The friend function is declared using the friend keyword inside the body of the class.

Friend Function Syntax:

1 class className { 2 3 friend returnType functionName(arguments); 4 5 }

By using the keyword, the 'friend' compiler understands that the given function is a friend function.

We declare friend function inside the body of a class, whose private and protective data needs to be accessed, starting with the keyword friend to access the data. We use them when we need to operate between two different classes at the same time.

To declare a function as a friend of a class, precede the function prototype in the class definition with keyword friend as follows -

class Box { double width:

> public: double length; friend void printWidth(Box box); void setWidth(double wid);

};

```
Define Inline function .write a c++ program for finding the area of a
                                                                           20ESX02.2
                                                                                         L1
       triangle using inline functions.
                                        Inline function-2m program -6m
9 (b)
       8m
```

Inline function in C++ is an enhancement feature that improves the execution time and speed of the program. The main advantage of inline functions is that you can use them with C++ classes as well.

// calculate an area of triangle using the inline function #include <iostream> using namespace std;

9 (a)

```
// inline function, no need prototype
inline float triangle_area(float base, float height)
float area;
area = (0.5 * base * height);
return area;
}
int main(void)
float b, h, a;
b = 4;
h = 6;
// compiler will substitute the inline function code here.
a = triangle_area(b, h);
cout<<"Area = (0.5*base*height)"<<endl;
cout<<"where, base = 4, height = 6"<<endl;
cout<<"\nArea = "<<a<<endl;
return 0;
}
Output example:
Area = (0.5*base*height)
where, base = 4, height = 6
Area = 12
```

10 (a) What is a constructor ?Explain with an example constructor-2m example-2m 4m

20CS304.3 L1

A constructor is a special type of member function of a class which initializes objects of a class. In C++, Constructor is automatically called when object(instance of class) is created. It is special member function of the class because it does not have any return type.

using namespace std;

class construct

{

public:

```
int a, b;
  // Default Constructor
    construct()
    {
      a = 10;
      b = 20;
    }
 };
  int main()
 {
    // Default constructor called automatically
    // when the object is created
    construct c;
    cout << "a: " << c.a << endl
       << "b: " << c.b;
    return 1;
 }
Output:
a: 10
```

b: 20

10 (b) Write a c++ program to find the area of circle ,rectangle and triangle 20CS304.3 L3 using function overloading. 6m

* C++ program to find Area using Function Overloading */

```
#include<iostream>
using namespace std;
int area(int);
int area(int,int);
float area(float);
float area(float,float);
int main()
{
     int s,l,b;
      float r,bs,ht;
      cin>>s;
      cin>>l>>b;
```

cout<<"Enter side of a square:";

cout<<"Enter length and breadth of rectangle:";

cout<<"Enter radius of circle:";

cin>>r;

cout<<"Enter base and height of triangle:";

cin>>bs>>ht;

cout<<"Area of square is"<<area(s);

cout<<"\nArea of rectangle is "<<area(l,b);

```
cout<<"\nArea of circle is "<<area(r);
```

cout<<"\nArea of triangle is "<<area(bs,ht);

```
int area(int s)
```

```
retum(s*s);
```

}

{

}

{

```
int area(int l,int b)
```

```
return(l*b);
```

```
}
float area(float r)
{
```

```
return(3.14*r*r);
```

```
}
float area(float bs,float ht)
{
    return((bs*ht)/2);
}
OUTPUT ::
/* C++ program to find Area using Function Overloading */
```

Enter side of a square:2 Enter length and breadth of rectangle:3 6 Enter radius of circle:3 Enter base and height of triangle:4 4

Area of square is4 Area of rectangle is 18 Area of circle is 28.26 Area of triangle is 8

OR

11 (a) What is inheritance ? present the advantages and disadvantages of 20CS304.3 L1 inheritance? Defination -2m, advantages-2m, Disadvantages-2m In C++, inheritance is a process in which one object acquires all the properties and behaviors of its parent

object automatically. In such way, you can reuse, extend or modify the attributes and behaviors which are defined in other class.

In C++, the class which inherits the members of another class is called derived class and the class whose members are inherited is called base class. The derived class is the specialized class for the base class.

Advantages:

- Inheritance promotes reusability
- Inheritance allows us to inherit all the properties of base class and can access all the functionality of inherited class. It implements reusability of code.

Disadvantages:-

- Inherited functions work slower than normal function as there is indirection.
- Improper use of inheritance may lead to wrong solutions.
- Often, data members in the base class are left unused which may lead to memory wastage.
- Inheritance increases the coupling between base class and derived class. A change in base class

will affect all the child classes.

12

```
Write a c++ program to overload + operator to add two matrices.
                                                                                                          L2
                                                                                          20ESX02.3
         Program -5m output-1m 6m
11 (b)
         #include<iostream>
         using namespace std;
         class Matrix
         {
              int a[3][3];
            public:
              void accept();
               void display();
              void operator +(Matrix x);
         };
         void Matrix::accept()
          {
               cout<<"\n Enter Matrix Element (3 X 3) : \n";
               for(int i=0; i<3; i++)
               {
                    for(int j=0; j<3; j++)
                    {
                         cout<<" ";
                         cin>>a[i][j];
                    }
               }
          }
          void Matrix::display()
          {
               for(int i=0; i<3; i++)
               {
                     cout<<" ";
                     for(int j=0; j<3; j++)
                     {
                          cout<<a[i][j]<<"\t";
```

```
}
           cout<<"\n";
     }
}
void Matrix::operator +(Matrix x)
{
     int mat[3][3];
     for(int i=0; i<3; i++)
     {
           for(int j=0; j<3; j++)
           {
                mat[i][j]=a[i][j]+x.a[i][j];
           }
     }
     cout<<"\n Addition of Matrix : \n\n";
     for(int i=0; i<3; i++)
     {
           cout<<" ";
           for(int j=0; j<3; j++)
           {
                cout<<mat[i][j]<<"\t";
           }
           cout<<"\n";
     }
}
int main()
{
     Matrix m,n;
     m.accept();
                     // Accepting Rows
                   // Accepting Columns
     n.accept();
     cout<<"\n First Matrix : \n\n";
     m.display(); // Displaying First Matrix
     cout<<"\n Second Matrix : \n\n";
     n.display(); // Displaying Second Matrix
               // Addition of Two Matrices. Overloaded '+' Operator
     m+n;
     return 0;
```

}

789

Addition of matrix:

5 7 9 5 7 9 14 16 18

12 a) Explain the role of this poiter in c++ with a programming example 5m This pointer -2m example-3m

20CS304.4 L2

The 'this' pointer is passed as a hidden argument to all nonstatic member function calls and is available as a local variable within the body of all nonstatic functions. 'this' pointer is not available in static member functions as static member functions can be called without any object (with class name).

using namespace std;

/* local variable is same as a member's name */

class Test

{

private:

int x;

public:

```
void setX (int x)
```

{

// The 'this' pointer is used to retrieve the object's x

// hidden by the local variable 'x'

this-x = x;

}

```
void print() { cout \leq x = \leq x \leq endl; }
        };
        int main()
        {
          Test obj;
         int x = 20;
          obj.setX(x);
          obj.print();
          return 0;
        }
        Output:
        x = 20
        What is meant by late binding ?how is it implemented in c++ ?
                                                                                                         7m
12(b)
        late binding-2m implementation-5m
                                                                                          20CS304.4
                                                                                                             L3
        Late Binding : (Run time polymorphism) In this, the compiler adds code that identifies the kind of object
        at runtime then matches the call with the right function definition (Refer this for details). This can be
        achieved by declaring a virtual function.
        using namespace std;
        class Base
        {
        public:
          virtual void show() { cout<<" In Base \n"; }
        };
        class Derived: public Base
        {
        public:
          void show() { cout<<"In Derived \n"; }</pre>
        };
        int main(void)
        Ł
           Base *bp = new Derived;
           bp->show(); // RUN-TIME POLYMORPHISM
          return 0;
        }
        Output:
        In Derived
                                                           OR
        What is virtual destructor ?explain with an example? Virtual destructor-2m example-4m 6m
 13
        20CS304.4
                       L2
```

15

A virtual destructor is used to free up the memory space allocated by the derived class object or instance while deleting instances of the derived class using a base class pointer object. A base or parent class destructor use the virtual keyword that ensures both base class and the derived class destructor will be called at run time, but it called the derived class first and then base class to release the space occupied by both destructors.

#include<iostream> using namespace std; class Base

```
public:
  Base() // Constructor member function.
  cout << "\n Constructor Base class"; // It prints first.
virtual ~Base() // Define the virtual destructor function to call the Destructor Derived function.
  cout << "\n Destructor Base class"; /
}:
// Inheritance concept
class Derived: public Base
```

{

ł

}

public:

Derived() // Constructor function.

{

cout << "\n Constructor Derived class" ; /* After print the Constructor Base, now it will prints. */

}

~Derived() // Destructor function

ł

cout << "\n Destructor Derived class"; /* The virtual Base Class? Destructor calls it before calling the Base iss Destructor. */

```
<u>}:</u>
int main()
```

}

Base *bptr = new Derived; // A pointer object reference the Base class. delete bptr; // Delete the pointer object.

OUTPUT:

}

CONSTRUCTOR BASE CLASS CONSTRUCTOR DERIVED CLASS DESTRUCTOR DERIVED CLASS DESTRUCTOR BASE CLASS

14 Discuss about STL programming model

20CS304.5 L2

The Standard Template Library (STL) is a set of C++ template classes to provide common programming data structures and functions such as lists, stacks, arrays, etc. It is a library of container classes, algorithms, and iterators. It is a generalized library and so, its components are parameterized. A working knowledge of template classes is a prerequisite for working with STL.

- STL has four components
- Algorithms
- Containers
- Functions
- Iterators

Algorithms

The header algorithm defines a collection of functions especially designed to be used on ranges of elements. They act on containers and provide means for various operations for the contents of the containers.

- Algorithm
 - <u>Sorting</u>
 - Searching
 - Important STL Algorithms
 - Useful Array algorithms
 - Partition Operations
- Numeric
 - valarray class

Containers

Containers or container classes store objects and data. There are in total seven standard "first-class" container classes and three container adaptor classes and only seven header files that provide access to these containers or container adaptors.

- Sequence Containers: implement data structures which can be accessed in a sequential manner.
 - <u>vector</u>
 - <u>list</u>
 - <u>deque</u>
 - <u>arrays</u>
 - <u>forward_list(</u> Introduced in C++11)
- Container Adaptors : provide a different interface for sequential containers.
 - <u>queue</u>
 - priority queue
 - <u>stack</u>
- Associative Containers : implement sorted data structures that can be quickly searched (O(log n) complexity).
 - <u>set</u>
 - <u>multiset</u>
 - <u>map</u>

6m EXPLANATION STL PROGROGRAMING-6M

<u>multimap</u>

Functions

The STL includes classes that overload the function call operator. Instances of such classes are called function objects or functors. Functors allow the working of the associated function to be customized with the help of parameters to be passed.

<u>Functors</u>

Iterators

As the name suggests, iterators are used for working upon a sequence of values. They are the major feature that allow generality in STL.

<u>Iterators</u>

OR

DEFINE TEMPLATE WHAT IS THE NEED FOR TEMPLATES IN PROGRAMMING?WRITE C++ CODE THAT DECLARES A TEMPLATE CLASS TEMPLATE -2M NEED-2M CODE-1M 20CS304.5 5M

A **template** is a simple and yet very powerful tool in C++. The simple idea is to pass data type as a parameter so that we don't need to write the same code for different data types. For example, a software company may need sort() for different data types. Rather than writing and maintaining the multiple codes, we can write one sort() and pass data type as a parameter.

C++ adds two new keywords to support templates: 'template' and 'typename'. The second keyword can always be replaced by keyword 'class'.

- **15 a)** It allows you to define the generic classes and generic functions and thus provides support for generic programming. Generic programming is a technique where generic types are used as parameters in algorithms so that they can work for a variety of data types. Templates can be represented in two ways: Function templates
 - // C++ program to demonstrate the use of class templates
 - #include <iostream>

using namespace std;

// Class template

template <class T>

class Number {

private:

// Variable of type T

T num;

public:

Number(T n) : num(n) {} // constructor

T getNum() {

return num;

}

};

int main() {

// create object with int type

Number<int> numberInt(7);

// create object with double type

Number<double> numberDouble(7.7);

cout << "int Number = " << numberInt.getNum() << endl;</pre>

cout << "double Number = " << numberDouble.getNum() << endl;</pre>

return 0;

}___

Run Code

Output

int Number = 7

double Number = 7.7

15 b) WRITE A C++ PROGRAM THAT FILLS A VECTOR WITH RANDOM NUMBERS 7M PROGRAM -6M OUTPUT-1M 20CS304.5 L3 // C++ program to generate the vector // with random values #include <bits/stdc++.h> using namespace std;

// Driver Code int main() {

// Size of vector
int size = 5;

// Initialize the vector with
// initial values as 0
vector<int> V(size, 0);

// use srand() for different outputs
srand(time(0));

// Generate value using generate
// function
generate(V.begin(), V.end(), rand);

cout << "The elements of vector are:\n";

```
// Print the values in the vector
for (int i = 0; i < size; i++) {
    cout << V[i] << " ";
}</pre>
```

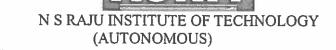
Nasimpall Salyanarayana Raju institute of Technology (Autonomous). IQAC: Quality Management System (QMS)

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Semester End Supplementary Examination, April/May, 2022 2021 - 2022 Academic Year Mechanical Engg. B. Tech. (U. G.) Program Degree 111 70 Semester 3 Hrs. Max, Marks Test Duration 20ME305 **Course Code** MANUFACTURING PROCESS Course Part A (Short Answer Questions 5 x 2 = 10 Marks) Learning Outcome (s) DoK Questions (1 through 5) No. L1 20ME305.1 Why is a loose piece pattern used? Give its problems. 1 20ME305.2 What are the essential conditions that are to be kept in mind while designing L2 2 risers? L2 20ME305.3 Write any four differences between the welding and soldering. 3 20ME305.4 11 Why do most welding failures occur in HAZ? Explain. 4 L1 20ME305.5 Differentiate between fullering and edging in forging operation 5 Part B (Long Answer Questions 5 x 12 = 60 Marks) DoK Learning Outcome (s) Marks No. Questions (6 through 15) 20ME305.1 L2 Explain the types of patterns with a neat sketch 6M 6 (a) 20ME305.1 L1 6M With help of sketch explain gating system. 6 (b) OR L2 20ME305.1 6M Explain injection molding and Blow molding. 7 (a) 20ME305.1 L1 6M 7 (b) Explain about hand molding process. Explain the construction and working principle of Cupola L2 6M 20ME305.2 8 (a) Furnace with a neat sketch. 20ME305.2 L2 6M Explain the principle of gating and gating ratio 8 (b) OR 20ME305.2 L.2 6M What are the types of centrifugal casting? 9 (a) 20ME305.2 for which Which centrifugal casting method is used L2 6M 9 (b) application? L2 7M 20ME305.3 How submerged arc welding process takes place. 10 (a) 20ME305.3 L2 5M Describe its advantages and applications. 10 (b) **OR** 20ME305.3 Explain the TIG systems of arc-welding give the applications L2 6M 11 (a) of each. 20ME305.3 Explain the MIG systems of arc-welding give the applications L2 6M 11 (b) of each. L2 5M 20ME305.4 Describe the electro slag welding process 12 (a) 20ME305.4 L2 7M Describe the electron beam welding process 12 (b) OR L2 20ME305.4 6M Explain the method and application of friction stir welding. 13 (a) L2 20ME305.4 6M Explain any two destructive testing of welds. 13 (b) L1 7M 20ME305.5 Describe the wire drawing process. 14 (a) L2 20ME305.5 5M Describe the tube drawing process. 14 (b) OR 20ME305.5 L1 4M Enumerate the typical applications of cold working. 15 (a) L2 20ME305.5 **8**M Explain the blanking and piercing process with a neat sketch. 15 (b)

Automotive Party Parity in the optimal Autom

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SONTYAM, ANANDAPURAM, VISAKHAPATNAM – 531 173

ANSWER KEY AND SCHEME OF EVALUATION

MANUFACTURING PROCESS (20ME 305)

ME-III sem

Part-A (5*2=10)

1. Why is loose piece pattern used? Give its problems.

Loose piece of a pattern which is attached to the pattern during molding, remains in the mold during lift-off and is only then removed separately. It is used for undercuts if the mold joint cannot be positioned in this cross section.

2. What are the essentials conditions that are to be kept in mind while designing Risers? 1) The modulus of the riser should be larger than the modulus of the casting, which encourages directional solidification, insuring that feed metal will be available to counteract shrinkage in the casting throughout solidification

2) The riser should have enough volume to provide the required feed metal to the casting.

3. Write any four differences between Welding and Soldering?

S.NO	Welding	Soldering
1	metal fabricators melt the base	metal fabricators heat the metal to be
	metal.	bonded but never melt them
2	Welding joints are the strongest, followed by soldered joints then brazed joints.	Soldering is most similar to brazing because it uses capillary action to flow the metal into the joint until it cools and hardens.
3	Welding requires about 6,500 degrees Fahrenheit, while	soldering requires about 840 degrees Fahrenheit.
4	Workpieces and the metal base are heated and melted in welding	Soldering requires no heating of the workpieces.

4. Why do most welding failures occur in HAZ? Explain?

The heat produced in the weld bead area causes chromium carbides to precipitate around the grain boundaries in the HAZ, causing the local chromium content to drop below 10.5%, at which point the steel loses the ability to form a passive film and is no longer stainless.

5. Differentiate between fullering and edging in forging operation?

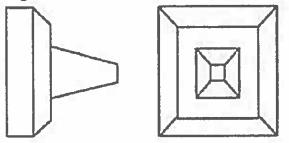
the flow of material on impact of the die is the opposite of fullering. When the concave die deforms the workpiece the material flows into the die area from both sides. The process is called edging because it is usually carried out on the edges of the workpiece.

Part-B

6(a). Explain the types of patterns with neat sketch? (6M)

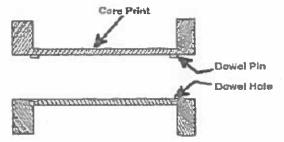
We use 10 different types of patterns in the casting process. Single piece pattern, two piece pattern, gated pattern, multi piece pattern, match plate pattern, skeleton pattern, sweep pattern, lose piece pattern, cope and drag pattern, shell pattern.

There have more details about 10 different types of patterns. Single Piece Pattern



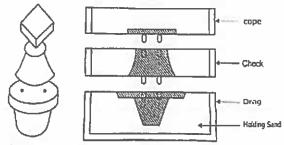
Single piece pattern, also called solid pattern is the lowest cost casting pattern. It is very suitable for simple process, and small scale production and the large casting manufacturers prefer it because this kind of casting pattern make casting process just needing simple shapes, flat surfaces like simple rectangular blocks. One flat surface is used to separate planes.

Two-Piece Pattern

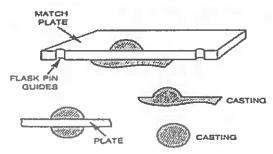


Two-piece pattern also called split piece pattern is a common casting pattern for intricate casting. This kind of pattern has parting planes which may have flat or irregular surface, and the exact position of the plane was decided by the shape of the casting. There are two pieces of the split piece pattern. One of the parts is molded in drag and another is molded in cope. And the cope part always has dowel pins. With the dowel pins, the two halves of split piece pattern can be aligned.

Multi Piece Pattern

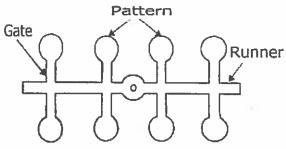


Multi piece pattern is a good solution for complex designs which is hard to make. This kind of pattern includes 3 or more pattern which helps you achieve mold making. Take the three-piece pattern as an example. The pattern is made of the top, bottom, and middle parts. The top part is cope, the bottom part drag, and the middle parts are called as checkbox. Match Plate Pattern



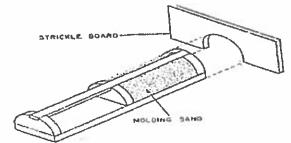
Match plate pattern has a metallic plate to divide the cope and drag areas into the opposite face of the plate. This kind of pattern nearly has no hard work and can provide high output. It is widely used in the manufacturing industry, and usually has an expensive cost, precise casting and high yield. And this kind of casting pattern is widely used in metal casting like aluminum.

Gate Pattern



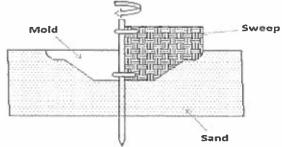
Gate pattern can consist of one or more patterns into a molding pattern. It is designed for the mold which makes multiple components at one casting process. The gates are used to combine the different patterns, and runners to create a flow way for the molten materials. When the gates and runners have already attached, the patterns are loosing. This kind of pattern is expensive, and it is usually used for small castings.

Skeleton Pattern



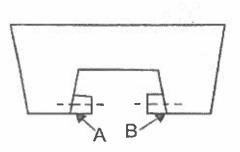
Skeleton pattern is large in size, and it is a good choice for the casting which has the simple size and shape. This kind of casting pattern is expensive and not versatile. It is not the best choice from the aspect of economic, while is very efficient in extra sand removing. If you want to use this casting pattern you should highlight the wood frames when you casting. The skeleton pattern is widely used in the industries of pit or floor welding.

Sweep Pattern



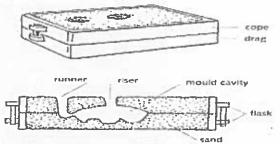
Sweep pattern uses a wooden board with proper size to rotate along one edge to shape the cavity. This kind of casting pattern creates a cavity in the vertical direction and the base of it is attached with sand, and it also creates casting in a very short time, and it has consisted of three parts: spindle, base and sweep which also called wooden board.

Loose Piece Pattern



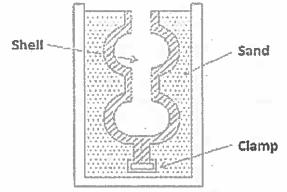
Loose piece pattern can help manufacturers remove one piece of solid pattern which is above or below the parting plane of the mold. This kind of pattern needs extra skilled labor work, so it is expensive casting pattern in castings.

Cope and Drag Pattern



Just like its name, cope and drag pattern has consisted of two separate plates, and it has two parts which can be separately molded on the molding box, and these parts create the cavity. This kind of pattern has a bit similar with the two-piece pattern and is usually used in large casting.

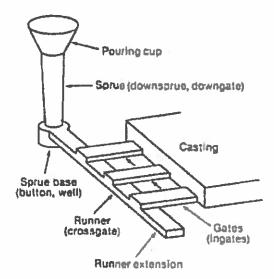
Shell Pattern



Shell pattern is a good choice to create hollow shaped structure. It parts along the center and dowels the resultant halves.

6(b). With help of sketch explain gating system? (6M)

- Elements of Gating System
 - Pouring Cup
 - Spruce
 - Spruce Well
 - Cross-gate or Runner
 - Ingate or Gates



- Pouring Cup It is the funnel-shaped opening, made at the top of the mold. The main purpose of the pouring basin is to direct the flow of molten metal from ladle to the sprue.
- Spruce It is a vertical passage connects the pouring basin to the runner or ingate. It is generally made tapered downward to avoid aspiration of air. The cross section of the sprue may be square, rectangular, or circular.
- Spruce Well It is located at the base of the sprue. It arrests the free fall of molten metal through the sprue and turns it by a right angle towards the runner.
- Runner It is a long horizontal channel which carries molten metal and distribute it to the ingates
 .It will ensure proper supply of molten metal to the cavity so that proper filling of the cavity takes
 place.
- Gate These are small channels connecting the mould cavity and the runner. The gates used may vary in number depends on size of the casting.
 Function of Gating System
- A good gating system should help easy and complete filling of the mould cavity.
- It should fill the mould cavity with molten metal with least amount of turbulance.
- It should prevent mould erosion.
- It should establish proper temperature gradient in the casting.
- It should promote directional solidification.
- It should regulate the rate of flow of metal into the mould cavity.

(OR)

7(a). Explain Injection molding and blow molding? (6M)

Injection molding is a method to obtain molded products by injecting plastic materials molten by heat into a mold, and then cooling and solidifying them.

The method is suitable for the mass production of products with complicated shapes, and takes a large part in the area of plastic processing.

The process of injection molding is divided into 6 major steps as shown below.

- 1. Clamping
- 2. Injection
- 3. Dwelling
- 4. Cooling
- 5. Mold opening
- 6. Removal of products

The process is proceeded as shown above and products can be made successively by

repeating the cycle.

Injection molding machine is divided into 2 units i.e. a clamping unit and an injection unit.

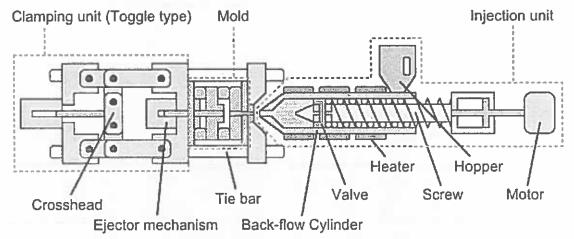
The functions of the clamping unit are opening and closing a die, and the ejection of products. There are 2 types of clamping methods, namely the toggle type shown in the figure below and the straight-hydraulic type in which a mold is directly opened and closed with a hydraulic cylinder.

The functions of the injection unit are to melt plastic by heat and then to inject molten plastic into a <u>mold</u>.

The screw is rotated to melt plastic introduced from the hopper and to accumulate molten plastic in front of the screw (to be called metering). After the required amount of molten plastic is accumulated, injection process is stared.

While molten plastic is flowing in a mold, the machine controls the moving speed of the screw, or injection speed. On the other hand, it controls dwell pressure after molten plastic fills out cavities.

The position of change from speed control to pressure control is set at the point where either screw position or injection pressure reaches a certain fixed value.



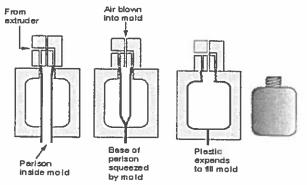
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Blow molding (or moulding) is a manufacturing process for forming and joining together hollow <u>plastic</u> parts. It is also used for forming <u>glass bottles</u> or other hollow shapes.

In general, there are three main types of blow molding: extrusion blow molding, injection blow molding, and injection stretch blow molding.

The blow molding process begins with softening plastic, by heating, and forming it into a <u>parison</u> or, in the case of injection and injection stretch blow molding (ISB), a preform. The parison is a tube-like piece of plastic with a hole in one end through which compressed air can pass.

The parison is then clamped into a <u>mold</u> and air is blown into it. The air pressure then pushes the plastic out to match the mold. Once the plastic has cooled and hardened the mold opens and the part is ejected. Water channels within the mold assist cooling.



7(b). Explain about hand molding process? (6M)

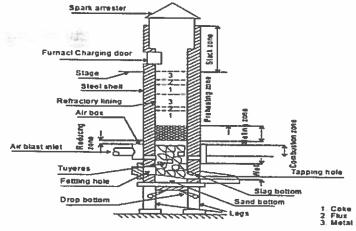
Whatever the force required for ramming and compressing of molding sand, if it is obtained from a human hand, then it is called as hand molding.

The properties of Hand Moulding method are as follows:

- It is the cheapest method of mold making operation.
- Complex shapes of the patterns will be easily used in mold making without any damages to the pattern.
- Strength and Hardness of the mold are non-uniform because of the non-uniform force applied by hand.
- The production rate is low.

The other method used for preparing the mold is Machine Molding method. It generally consists of 4 operations. By using these operations only, the mold can be prepared.

8(a). Explain the construction and working principle of cupola furnace with a neat sketch? (6M)



Working Principle of Cupola Furnace:

The Cupola furnace works on the principle where we generate heat from burning coke and when the temperature of the furnace is above the melting point of the metal then the metal is melt.

The charge introduced in the cupola consists of pig iron, scrap, casting rejection, coke, and flux. Coke is the fuel and limestone are added as a flux to remove undesirable materials like ash and dirt. The scrap consists of Steel and cast iron rejections.

The working of Cupola furnace is, Over the sand Bottom, Coke in charged extending up to a predetermined height. This serves as the coke bed within which the combustion takes place.

Cupola operation is started by igniting the coke bed at its bottom. After the Coke bed is properly Ignited, alternate charges of limestone, pig iron, and coke are charged until the level of the charging Door.

Then the air blast is turned on and combustion occurs rapidly within the coke bed. Within 5 to 10 minutes after the blast is turned on the first molten cast iron appears at the tap hole.

Usually, the first iron which comes out will be too cold to pour into sand molds. During the cupola operation, molten metal may be tracked every 10 minutes depending on the melting rate and the capacity. All the oxygen in the air blast is consumed by the combustion, Within the combustion zone.

-

The chemical reaction takes place which is,

C + O2 (from the air) $\rightarrow CO2 + Heat$

This is an exothermic reaction. The temperature in this zone varies from 1550 to 1850 degree Celsius.

Then hot gases consisting principally of Nitrogen and carbon dioxide moved upward from the combustion zone, where the temperature is 1650 degree Celsius.

The portion of the coke bed if the combustion zone is reducing zone. It is a protective zone to prevent the oxidation of the metal charge above and while dropping through it. As the hot carbon dioxide gas moves upward through the hot coke, some of it is reduced by the following reaction.

 $3Fe + 2CO \rightarrow Fe3C + CO2$

This is an endothermic reaction.

The first layer of iron above the reducing zone is the melting zone where the solid iron is converted into the molten state. A significant portion of the carbon is picked up by the metal also takes place in this zone.

The hot gas is passed upward from the reducing and melting zones into the preheating zone which includes all layers of charge above the melting zone up to the charging Door.

Since the layer of the charge is preheated by the outgoing gases which exist at the top of the cylindrical shell, this temperature is this zone is around 1090 degrees Celsius.

8(b). Explain the principle of gating and gating ratio? (6M)

A gating system should avoid sudden or right angle changes in direction. A gating system should fill the mould cavity before freezing. The metal should flow smoothly into the mould without any turbulence. A turbulence metal flow tends to form dross in the mould.

Gating ratio

Gating ratio is the ratio between the cross-sectional area of the sprue to the total cross-sectional area of the runners to the total cross-sectional area of the ingates.

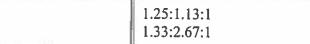
The formula for the gating ratio is As: Ar: Ag.

With the Pressurized Gating System, the gating ratio is usually 1: 2: 1 or 1: 0.75: 0.5. This system is called a "Gate control system" because ingates control the flow of the metal.

With the Unpressurized Gating System, the gating ratio is usually 1: 2: 2 or 1: 3: 3 or 1: 1: 3. This system is called a "*Choke control system*" because the choke controls the flow of the metal.

Table of gating ratio for various of materials:

Materials	Gating ratio		
Aluminum	1:2:1 1:1.2:2 1:2:4 1:3:3 1:4:4 1:6:6		
Aluminum bronze	1:2.88:4.8		
Brass	1:1:1 1:2:3 1.6:1.3:1		
Copper	2:8:1 3:9:1		
Ductile iron	1.15:1.1:1		



Gating ratio of materials

(OR)

9(a). What are the types of centrifugal casting? (6M)

Working Principle of Centrifugal Casting

In this process molten metal is poured into the spinning mold preheated to a certain temperature. The mold is placed vertically or horizontally based on the required shape of product. Once poured it is then continued to rotate about its central axis.

Due to the rotational motion of the mold; a centrifugal force is acted upon the molten metal just poured into the spinning mold. This force displaces the molten metals towards the periphery forcing them to deposit on the walls.

The molten metal is spread uniformly on to the walls of the die; thanks to the centrifugal force 100 times greater than of gravity.

As the process continues with more and more metal poured into the mold; the relatively denser element tends to deposit on towards the wall while lighter elements and slug deposit at the center.

The mold is then left to rotate till the whole mold solidify and then other light elements like slag are separated from the center.

The whole process itself leads to a reduction in defects due to slags, irregular grain structure, and trapped air. The final product have closed grain structure with improved elongation, tensile strength, and yield strength. Fig: Construction of True centrifugal casting Machine

Advantages Of Using Centrifugal Casting

The whole process of centrifugal casting does not rely on gates and risers. This results in the continuous availability of molten metal at the center during the solidification process.

Unlike conventional casting where the mold solidifies from both inside and outside; molten metal at the center ensures unidirectional solidification (Outside to Inside).

This helps get rid of defects such as blow holes, shrinkage cavity and gas pockets. The process allows for subsequent savings on initial capital (*Non-machinery cost*) and manufacturing costs.

This allows for the economic advantage which gave enough flexibility to produce various shapes and sizes of products.

The combination of deciding factors such as unidirectional solidification, solidification under pressure, and impurities displacement to the central axis; results in superior quality products with high soundness compared to other industrial processes.

Such Advantages result in increased life and endurance of the product without fracturing.

Other benefits of centrifugal casting includes:

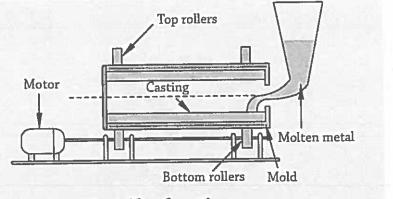
- 1. Impurities collected at the center are relatively easier to remove than other industrial or manufacturing process.
- 2. Mass production of symmetrical product was made possible at a much lower cost.
- 3. It provide dense metal with mechanical soundness.
- 4. Relatively less temp of molten metal is required for the process saving money and energy.
- 5. Gates and risers are no longer needed.
- 6. After the casting process is heavily reduced.

Fig: Separation of impurities toward the center during the whole process.

Types of Centrifugal Casting

The process of using centrifugal force for casting can be divided into three major parts; True centrifugal, semi-centrifugal and centrifuging.

1) True Centrifugal Casting



Centrifugal casting

True centrifugal casting or commonly known as normal centrifugal casting is used to produce a symmetrical hollow structure with round holes. The key feature of this process is to produce a symmetrical hollow structure without using any cores. It is achieved by pure centrifugal force by rotating mold about its vertical or horizontal axis.

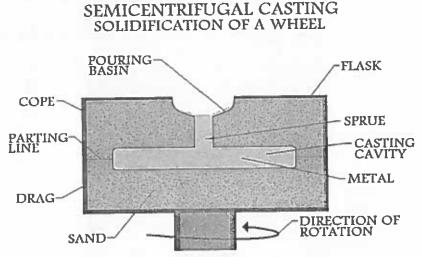
The shape of the mold can be either circular, square, rectangular or hexagonal; as long as they are symmetrical about its vertical or horizontal axis of rotation.

Centrifugal force acting on the molten metal introduced in the system; force it towards the wall of the mold/die. The casting of long parts such as pipes and liners are done along the horizontal axis while for others along the vertical axis.

To avoid molten metal to take the parabolic path along the mold while hardening due to gravity; the mold is subjected to high speed rotation to produce centrifugal force 100 times stronger than of gravity.

The mold is then left to rotate on its axis till it solidifies unidirectionally. This process is generally used for producing large and medium-sized parts such as cylinder liner, hollow pipes, and bushes.

2) Semi-Centrifugal Casting

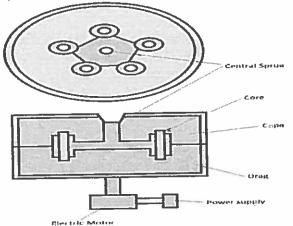


A core is placed inside the mold in a semi centrifugal casting to produce a hollow structure. The process itself includes the rotation of mold along its vertical axis with the core inserted at the center.

The centrifugal force from the rotation is then used to fill the mold perfectly. Here the hot metal is poured along the axis to first fill along the walls with the centrifugal force and towards center due to gravity.

A core is inserted whenever a hollow structure is needed to compose compensating the gravitational forces (dominant) at the center. This process is used to produce large Axis symmetrical products such as flywheel and gear blanks.

3) Centrifuging



This casting process is used to cast metal under high pressure for relatively small mold/die / final product or asymmetrical products in group.

The process is done in group to obtain overall symmetry of the casting. Hot metal is poured into the die along the central axis through central spur using centrifugal force through radial in gates.

It is an easy method for obtaining unidirectional solidification at economical costs.

9(b). Which centrifugal casting method is used for which applications? (6M)

Centrifugal casting provides high material soundness and is the metal casting process of choice for jet engine compressor cases, petrochemical furnace tubes, many military and defense components, and other applications requiring high reliability.

In centrifugal casting molten metal is poured into a spinning die, which can rotate on a vertical axis (vertical centrifugal casting) or horizontal axis (horizontal centrifugal casting) depending on the configuration of the part. Ring and cylinder type shapes are made by vertical centrifugal casting while tubular shapes are made by horizontal centrifugal casting. Either process can be used to produce multiple parts from a single casting.

High centrifugal force applied to the molten metal in the spinning die causes less dense material such as oxides and impurities to "float" to the inner diameter (I.D.) where they concentrate and are removed by machining. Solidification is managed directionally under pressure from the outer diameter (O.D.) to the I.D., avoiding mid-wall shrink. This results in a defect-free structure without cavities or gas pockets.

With the world's largest and most diverse inventory of centrifugal dies, MetalTek minimizes both upfront tooling costs and product lead times for our customers. We can produce horizontal centrifugal castings with O.D. up to 60" (1,524 mm), length up to 432" (10,973 mm), and weight up to 135,000 lbs. (61,235 kg). Vertical centrifugal castings are available with O.D. up to 180" (4,572 mm) and weight up to 34,000 lbs. (15,422 kg).

10(a). How submerged arc welding process takes place? (7M)

Submerged Arc Welding (SAW) is a joining process that involves the formation of an electric arc between a continuously fed electrode and the workpiece to be welded. A blanket of powdered flux surrounds and covers the arc and, when molten, provides electrical conduction between the metal to be joined and the electrode. It also generates a protective gas shield and a slag, all of which protects the weld zone. The make-up of the process can be viewed by reference to Figure 1 below

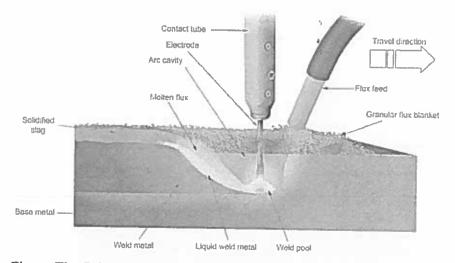


Figure The Submerged Arc Welding Process

As can be seen from Figure, the arc is "submerged" beneath a blanket of flux and is, therefore, not usually visi:ble during the welding operation itself. These facts make the process advantageous from a health and safety viewpoint as there is no arc to promote "arc eye" and very little fume.

There are two welding consumables involved in the process, the electrode and the flux. The electrode can be a solid wire, a cored wire, or a strip. The flux, made from a variety of minerals and compounds, can be rather complex and can be produced in a number of forms.

The general arrangement of the power source and controls, wire feed and flux dispensing are shown in Figure 2.

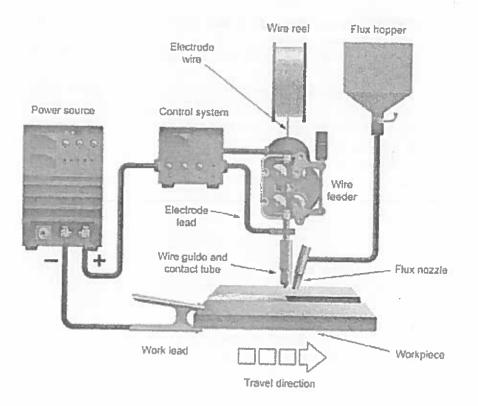


Figure 2. General Arrangement of the Submerged Arc Process Submerged arc welding is viewed as a high productivity process and is usually automated/mechanized in its form. The simplest application of the process uses a single wire.

Selecting the correct wire diameter for a welded joint depends on many factors and the size of the available power source usually limits the diameter of the wire that can be used. While most power sources for this process are 1,000 amps, smaller power source may be used. A 3/32-in.-dia. wire through to a 5/32-in.-dia. wire will run in the 300 to 900 amps range using direct current and with the electrode positive (DC+)

This welding process is typically suited to the longitudinal and circumferential butt welds required in the manufacture of pressure vessels and for joining plating and stiffeners in shipyards. Welding is positionally restricted and is normally carried out in the flat or horizontal positions because of the highly fluid weld pool, the molten slag, and the need to maintain a flux covering over the arc.

As with all welding processes the selection of the consumables (wire and flux) and other parameters such as amps, volts and travel speed are intended to give a weld deposit that satisfies the objectives of the designer. In the case of this welding process, since the arc is submerged, the welding operator cannot see the molten weld pool and must, therefore, very accurately set the welding parameters and location of the welding nozzle within the joint.

10(b). Describe its applications and advantages? (5M)

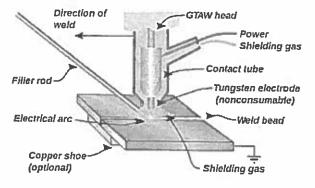
Submerged arc welding has many advantages but there are also restrictions, some of these are listed below Advantages

- High deposition rates and high arc on times when fully automated.
- Minimal welding fume, no weld spatter and no visible arc
- Unused flux can be recovered
- If metallurgically acceptable, single pass welds can be made in relatively thick plates.

here are many more applications of this welding process, other than its use with a single wire, and the fluxes used can be quite complex in their design and production. These items may be covered in later articles.



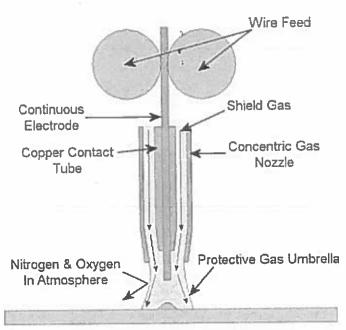
11(a). Explain the TIG systems in arc welding. Give its applications of each? (6M)



tungsten Inert Gas (TIG) welding uses the heat generated by an electric arc struck between a nonconsumable tungsten electrode and the workpiece to fuse metal in the joint area and produce a molten weld pool. The arc area is shrouded in an inert or reducing gas shield to protect the weld pool and the nonconsumable electrode. The process may be operated autogenously, that is, without filler, or filler may be added by feeding a consumable wire or rod into the established weld pool. TIG produces very high quality welds across a wide range of materials with thicknesses up to about 8 or 10mm. It is particularly well suited to sheet material.

The success of this welding process hinges on various factors such as the choice of <u>shielding gas</u>, welding wire, tungsten electrode and the welding technique.

11(b). Explain the MIG systems in arc welding. Give its applications of each? (6M)



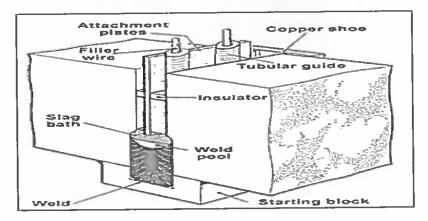
MIG/MAG welding is a versatile technique suitable for both thin sheet and thick section components. An arc is struck between the end of a wire electrode and the workpiece, melting both of them to form a weld pool. The wire serves as both heat source (via the arc at the wire tip) and filler metal for the <u>welding joint</u>. The wire is fed through a copper contact tube (contact tip) which conducts welding current into the wire. The weld pool is protected from the surrounding atmosphere by a shielding gas fed through a nozzle surrounding the wire. Shielding gas selection depends on the material being welded and the application. The wire is fed from a reel by a motor drive, and the welder moves the welding torch along the joint line. Wires may be solid (simple drawn wires), or cored (composites formed from a metal sheath with a powdered flux or metal filling). Consumables are generally competitively priced compared with those for other processes. The process offers high productivity, as the wire is continuously fed.

12(a). Describe the electro slag welding process? (5M)

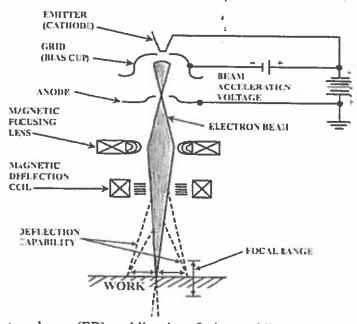
Electroslag Welding is a <u>welding process</u>, in which the heat is generated by an electric current passing between the consumable electrode (<u>filler metal</u>) and the work piece through a molten slag covering the weld surface.

Prior to welding the gap between the two work pieces is filled with a welding flux. Electroslag Welding is initiated by an <u>arc</u> between the electrode and the work piece (or starting plate). Heat, generated by the arc, melts the fluxing powder and forms molten slag. The slag, having low electric conductivity, is maintained in liquid state due to heat produced by the electric current.

The slag reaches a temperature of about 3500°F (1930°C). This temperature is sufficient for melting the consumable electrode and work piece edges. Metal droplets fall to the weld pool and join the work pieces. Electroslag Welding is used mainly for steels.



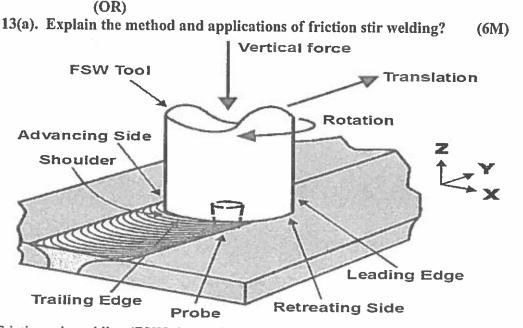
12(b). Describe the electron beam welding process? (7M)



Electron beam (EB) welding is a fusion welding process whereby electrons are generated by an electron gun and accelerated to high speeds using electrical fields. This high speed stream of electrons is tightly focused using magnetic fields and applied to the materials to be joined. The beam of electrons creates kinetic heat as it impacts with the workpieces, causing them to melt and bond together.

Electron beam welding is performed in a vacuum environment as the presence of gas can cause the beam to scatter. Due it being a vacuum process and because of the high voltages used, this welding method is heavily automated and computer controlled. As a result, specialised fixtures and CNC tables are used to move the workpieces inside the welding vacuum chamber.

Recent developments in electron beam welding machine technology have realised a local method of electron beam welding, whereby the electron beam gun is enclosed in a vacuum box on the side of the material to be joined, rather than placing the entire workpiece inside a vacuum chamber.



Friction stir welding (FSW) is a solid-state joining process that uses a non-consumable tool to join two facing workpieces without melting the workpiece material. Heat is generated by friction between the

rotating tool and the workpiece material, which leads to a softened region near the FSW tool. While the tool is traversed along the joint line, it mechanically intermixes the two pieces of metal, and forges the hot and softened metal by the mechanical pressure, which is applied by the tool, much like joining clay, or dough It is primarily used on wrought or extruded <u>aluminium</u> and particularly for structures which need very high weld strength. FSW is capable of joining aluminium alloys, copper alloys, titanium alloys, mild steel, stainless steel and magnesium alloys. More recently, it was successfully used in welding of polymers.^[2] In addition, joining of dissimilar metals, such as <u>aluminium</u> to magnesium alloys, has been recently achieved by FSW

13(b). Explain any two destructive testing of welds? (6M)

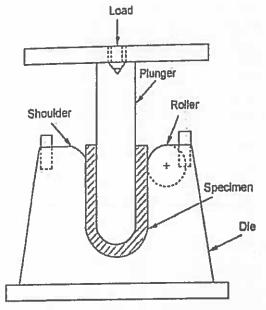
1. Acid Etch Test

This physical weld testing is employed to ascertain the soundness of the weld. The acid attacks the edge of the defects in base metal or weld metal and identifies the weld defects. In the condition of the defect, the boundary becomes accentuated between base and weld metals and can define the defect clearly which is otherwise not visible to the naked eye. This test is performed along the cross-section of the weld joint.

The acid solutions used here are hydrochloric acid, ammonium persulfate, nitric acid, or iodine, and potassium iodide for the etching of carbon and low alloy steel.

2. Guided Bend Test

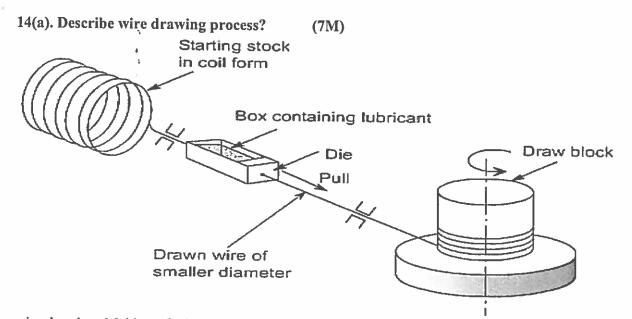
These guided bend tests are used to determine the quality of the weld metal at the root and face of the welded joint. They also judge the fusion and degree of penetration to the base metal along with the efficiency of the weld. The testing of this type can be done in a jig. The required specimens for testing are machined from the already welded plates, the thickness of these specimens should be within the capacity of our jig for bending. The specimen for testing is placed upon the supports of the die that is the lower part of the jig. The hydraulic jack's plunger forced the specimen into it and assured the shape of the die seen.



The requirement of this test is fulfilled by bending the specimens at 180 degrees and now accepted as passable. No, any crack more than 3.2mm in any dimension should be visible on the surface. Face bend tests are made in the jig while facing the weld in tension means outside of the bend. Now the root bend test is made in the jig with the face of the weld in tension as on the outside of the bend. The guided bend tests are shown in the figure.

Notes:

- T-Test plate thickness
- A hardened roll may be utilized on shoulders if needed
- Specific dimension for 3/7 of the plate
- Every shown dimension is in inches.

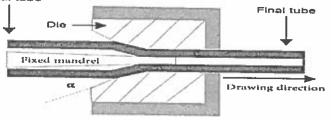


wire drawing, Making of wire, generally from a rod or bar. The wire-drawing process consists of pointing the rod, threading the pointed end through a die, and attaching the end to a drawing block. The block, made to revolve by an electric motor, pulls the lubricated rod through the die, reducing it in diameter and increasing its length. Fine wire is made by a multiple-block machine, because the reduction cannot be performed in a single draft.

14(b). Describe the tube drawing process? (5M)

Tube drawing is a metalworking process used to create a tube with a smaller diameter by pulling, or drawing, a larger diameter tube through a die. There are five methods of tube drawing that are commonly used. These methods are fixed plug drawing, floating plug drawing, tethered plug drawing, rod drawing and tube sinking.

This process is a cold-working process, meaning that the metal tubing is not heated prior to being shaped in the tube drawing process. This gives the finished product added strength because the metal tubing is not affected by thermal expansion during the process. In addition, this process produces tubing with more precise measurements than other methods of production.



(OR)

15(a). Enumerate the typical applications of cold drawing? (4M)

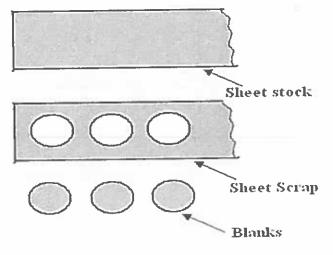
These include pump parts as well as valve stems, linear guide rails or sprockets. Furthermore, they can be gears, keyways, splines and spindles. Turbine parts, curtain wall and facade constructions are also possible. They even include x-ray equipment, louvers, and couplings.

15(b). Explain the blanking and piercing process with a neat sketch? (8M)

Blanking

Punching or blanking is a process in which the punch removes a portion of material from the larger piece or a strip of sheet metal. If the small removed piece is the useful part and the rest is scrap, the operation is called blanking The piece cut out is called as blank and may be further processed. Blanks are often cut out of a sheet or strip.

Blanking wastes certain amount of material. When designing a sheet metal blanking process the geometry of blanks should be nestled as efficiently as possible to minimize the material waste.



BLANKING

Piercing

It is a process by which a hole is cut (or torn) in metal. It is different from punching in that piercing does not generate a slug. Instead, the metal is pushed back to form a jagged flange on the back side of the hole. A pierced hole looks somewhat like a bullet hole in a sheet of metal.

Size of the component is generally larger in piercing than blanking.



ielReiulinslittleopregnology(stilonomme)/IOAex(QuelityManagementsSystem(GMS))

		:	Semeste	r End Sur	oplement	arv Exan	nination, Ap	ril/Mav	2022	, U.F.	14
Degree		B. Tech. (U		Program		EEE	initiation) esp			2021 -	2022
Course	Code	20EE305		Fest Dura			Max. Marks		iemester		
Course		POWER G					MON: MOIND	10 0	GIIICƏLCI		II.
Part A (S	Short An	swer Quest	ions 5 x 1	2 = 10 Mai	rke)						
No		ns (1 throug		c - 10 Mai	ingj				Learning Outo	come (e)	Do
1		luclear fissio							20EE30		L
2	Define t	hree part tar	iff.						20EE30		L
3		e GMR and							20EE30		L
4		y the trans		lines bas	sed on vo	ltage			20EE30		L
5	List the	various met	hods for in	norovina s	strina effici	ency			20EE30		L
	ona An	swer Quest	ions 5 x 1	$12 = 60 M_{\odot}$	arks)	onoy			ZULLJU	0.0	L
No.		ns (6 throug						Marks	Learning Outo	como (e)	Do
				in a therr	nal power	niant wi	th neat block	Mana	Learning Out	50me (3)	00
6	diagram	i. Also m	ention ti	he impor	tance of	hoiler	accessories	12M	20EE30	5 1	L
	econom	izer and cor	idenser w	ith neat di	adrams.	001101	00000000000	14.171	202200	0.1	1
					-3	OR					
7	Sketch	the layout of	of Hydro	power pla	int and bri		ain the main				
1	compon	ents and op	eration of	hydro pov	ver station	i i i i i i i i i i i i i i i i i i i		12M	20EE30	5.1	L
				- ,							
		ating station	has the f	following d	aily load c	ycle					
	Time	1 (1_6	6—10	10—12	12—16	16—20	20-24				
8	(Hour Load			122.016				4014			
0	(MW	1 40	50	60	50	70	40	12M	20EE30	15.2	L
			e and fine	l d (i) mavio	l	l Ind (ii) un	its generated				
	per da	y (iii) averag	e load an	d (iv) load	factor	ina (ii) an	ilis generaleu				
9	of 60% Find : i. the ii. the iii. the run iv the	daily energ reserve cap maximum e ning all the	pacity fac y produce pacity of the energy the time. nergy the	tor of 48% ne plant at could be at could be	and a pla e produced e produced	ant use fa d daily if i daily if	a load factor actor of 80% . the plant was the plant was	12M	20EE30)5.2	ι
				_							
10	Uerive transmi	the calculat ssion lines.	ion of ca	pacitance	tor 2-wire	e system	in overhead	l 12M	20EE30)5.3	L
	aaaann	oolori in 163.				OR					
	Find the	e inductanci	e per oba	ase oer kr	n of doub		3-phase line				
	shown i	n Figure. Th	ie conduc	tors are tr	ansposed	and are	of radius 0.75)			
			d _ ~ ~	479			4				
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14				3m				4014			
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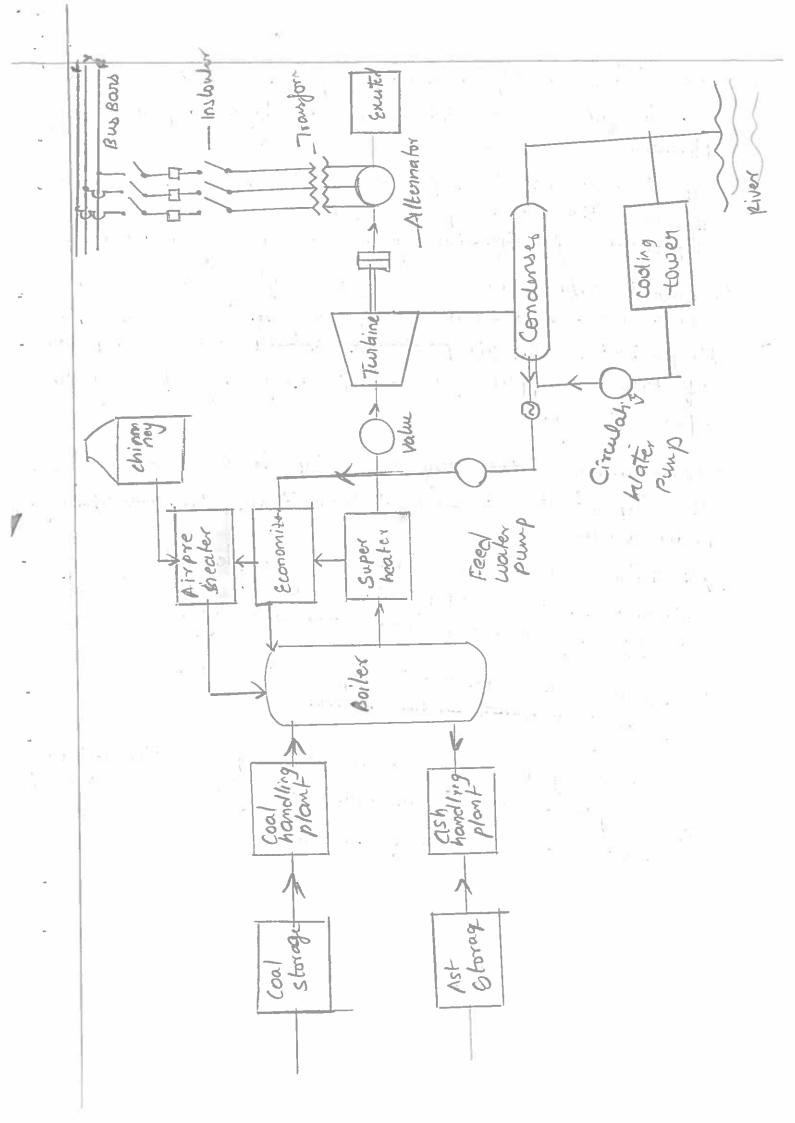
					-
12	Using nominal π method, derive an expression for sending end voltage and current for a medium transmission line.	12M	20EE305.4	L2	
13	A 3-phase, 50-Hz overhead transmission line 100 km long has the following constants : Resistance/km/phase = 0.1Ω Inductive reactance/km/phase = 0.2Ω Capacitive susceptance/km/phase = 0.04×10^{-4} siemen Determine (i) the sending end current (ii) sending end voltage (iii) sending end power factor and (iv) transmission efficiency when supplying a balanced load of 10,000 kW at 66 kV, pf. 0.8 lagging. Use nominal T method.	12M	20EE305.4	L3	
14	A 3-phase transmission line is being supported by three disc insulators. The potentials across top unit (i.e., near to the tower) and middle unit are 8 kV and 11 kV respectively. Calculate (i) the ratio of capacitance between pin and earth to the self-capacitance of each unit (ii)the line voltage and (iii) string efficiency.	12M	20EE305.5	L2	
	OR	6M	20EE305.5	L2	(
15 (a)	Explain any two types of insulators with neat sketch.	UW	20220000	1	
15 (b)	Derive the expression for the Sag in horizontal plane with equal level supports.	6M	20EE305.5	L2	

G. Kelyowi Controller of Examinations NSRIT (A) Visakhapatnam

the free

SEPPLE MENTARY EXAMINATION KEYS PARTA SHORT 1. Recall Nuclean-fission. A. The breaking up of nuclei of heavy atoms into two nearly equal parts with release of huge amount of onergy is known as nucleon fission. 2. Define three post lasiff. A. When the mate of electrical energy is changed on the "Consumer and the basis of maximum demands of the units consumed, it is called a two A When the total charge to be made from the consumer is split into three post · fixed aborge · Semi charge fixed. Charge . running charge. it is known as three post tarilf. 3. Compose GMR and GHD GMR : 1) It is also called Self GMD. (Ps) 2) It is The fall form of GMR is Geometrical mean radius GMD: 1) It is also Called as Mutual GMD is the geometrical mean of the distances form one conductor to the other and there must be tetween the largest and Smallest Such distance. 2) In fact, mutual - GMD Simply snepsi

4. Classify the transmission lines based on voltage. A. Classification Voltage Range (KV) Extra super Beyond 132KV voltage Cable. Extra high From 33 KV to 66 kv tensim Cable From 22 EV to 33 EV Super tension Cable From IKV Collkr High tension Cable upto to IKV Low tension Cable 5. List the Various methods for improving string mel Efficiency A.I) By using longer cross-arms 2) By grading the insulators 3) By using a guard ring PART-B 6. Describe the different blocks in a thermal power plant with next block diagram. Also mention the impollance of boiler accessories economizer and Condenser with next diagram.



Boiler: The heat of Combustion of Coal in the boiler is to utilis -ed to Convert water into steam at high temperate and. Penessure.

Super heater: The Steam produced in the boiler is wet and is passed through a Super heater where it is dried and Superheated by the flue gases on their way to chimney. Economiser: An economiser is essentially a feed water heater and derives heat from the flue gases for this purpose. Air preheater : An air preheater increases the temperature of the air Supplied for Coal Burning by deriving heat trom flue gases.

Steam twibine: The day and Superbeater Steam from the Superheater is fed to the Steam -twibine through main value. Alternator: The Steam -twibine is Coupled to an alternator. The Steam -twibine is Converts mechanical a alternator. The Stea alternator Converts mechanical a chergy of twibine into electrical energy.

the Eurbine

Cooling avaging-ement: In order to improve of the plant, the Steam exhausted from is condensed by means of Condenser.

7. layout of Hydro power plant.

sunge lant Reservoir pam value house Parelunec Penstock Power house River

if Dam: A dam is a barrier which Stores water and Coneates water head. (ii) Spillways. There are times when the rivers flow exceeds the storage Capacity of the reservoir. (iii) Head warks: The head warks consists of the diversion us) tread www. ine neur www. Internet the for an intake They generally Structures at the head of an intake They generally in clude booms and rack for diverting floating debries Stuices for by -paring debris and sediments and. Values for diverting for Controlling the flows of water to the twibines. (W) Surge tank: Open conduits leading water to the turb -ne require no protection. However, not However, when closed conduits are used, protection becomes Acts net (1) Penstocks: Penstocks are open or closed conduit which cevoy water to the twitines. (ui) coater twitine : latater twitishes are used to Converts the energy of falling water into meeranical energy. and turbine was coupled to an alternator and Converte the mechanical energy to electrical energy,

8. Given. data i'j It is clean from the load cure that manimum demand on Power Station is 70 MW. and occurs during the period 16-20 hours. : Massimum demand = 70 MW. (ii) units generated / day: Asrea Cin kuch) under the load Come = 10 [40×6+50×4 +60 ×2+ 50×4 +70×4 +40×4] = 10⁸ [240+200+120+200+280+160] kwn $= 12 \times 10^{5} \text{ kuch}.$ (in) Average load = Units generated / day $= \frac{12 \times 10^{5}}{50,000}$ (W) Load Jactos: Anonage load 50,000 = 0.714=71.4% 70×103 Max. demand 9 Given data Capacitance for 2 wire System in Ownhead 10. transiminer. Consider a single phase own head transmission line consisting of two Poorallel Conductors A and B spaced d-metres apart in air · Suppose that radius of ears conductor is r metors. Let their mespective change be + Q and - Q coulombs

$$V_{A} = \int_{-\infty}^{\infty} \frac{Q}{2\pi^{2} \cdot E_{0}} dx + \int_{-\infty}^{\infty} \frac{-Q}{2\pi^{2} \cdot E_{0}} dx$$

$$= \frac{Q}{2\pi^{2} \cdot E_{0}} \left[\log_{e} \frac{c_{0}}{1} - \log_{e} \frac{c_{0}}{2} \right] \text{volks} = \frac{Q}{2\pi^{2} \cdot E_{0}} \log_{e} \frac{d}{1} \text{ volk}$$
Similarly, P.d. between londuator B and nead-radiinfinite

$$V_{B} = \int_{-\infty}^{\infty} \frac{-Q}{2\pi^{2} \cdot E_{0}} dx + \int_{-\infty}^{\infty} \frac{Q}{2\pi^{2} \cdot E_{0}} dx$$

$$= \frac{-Q}{2\pi^{2} \cdot E_{0}} \left[\log_{e} \frac{c_{0}}{7} - \log_{e} \frac{c_{0}}{4} \right] = \frac{-Q}{2\pi^{2} \cdot E_{0}} \log_{e} \frac{d}{7} \text{ volk}$$
Both these potential one constitutes the same neartheal place

$$V_{AB} = 2V_{A} = \frac{2Q}{2\pi^{2} \cdot E_{0}} \log_{e} \frac{d}{7} \text{ volk}$$

$$C_{AB} = \frac{Q}{2\pi^{2} \cdot E_{0}} \log_{e} \frac{d}{7} = \frac{P}{m}$$

$$\frac{C_{AB}}{\log_{e} \frac{d}{7}} = \frac{T}{E_{0}} \frac{E_{0}}{\log_{e} \frac{d}{7}} = F/m$$
Solution.
Given data.
 $E_{HE} = Of Conductor = 0.975 \times 0.7488 = 0.584 \text{ cm}$

Distance a to $b = \sqrt{3^2 + (0.75)^2} = 3.1 \text{ m}$

Distance
$$a \cdot ba^{i} = \sqrt{3^{2} + (4i+5)^{2}} = 5\cdot62 m$$

Distance $a \cdot ba^{i} = \sqrt{6^{2} + 4^{2}} = 4 \cdot 21m$
Equivalent Self 6+10 Bf one phase is
 $D_{S} = \sqrt[3]{P_{SI} \times D_{S2} \times D_{S3}}$
 $D_{S1} = \sqrt[4]{D_{aa} \times D_{aa'} \times D_{bla'} \times D_{ala}}$
 $a \rightarrow 4m \rightarrow 4^{cl}$
 $b \rightarrow 5 \cdot 5m \rightarrow 4^{cl}$
 $c \rightarrow 4m \rightarrow 4^{cl}$
 $b \rightarrow 5 \cdot 5m \rightarrow 4^{cl}$
 $c \rightarrow 4m \rightarrow 4^{cl}$
 $c \rightarrow 4^{c}$
 $c \rightarrow 4^{cl}$
 $c \rightarrow 4^{c}$
 $c \rightarrow 4^{c}$
 $c \rightarrow 4^{c}$

= 4.17M = DBC

$$DCA = 4\sqrt{DCBX} PCal X Dcla X Dclal}$$

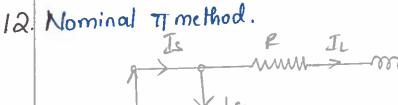
$$= 4\sqrt{6 X 4 X 4 X 6} = 4.9 m$$

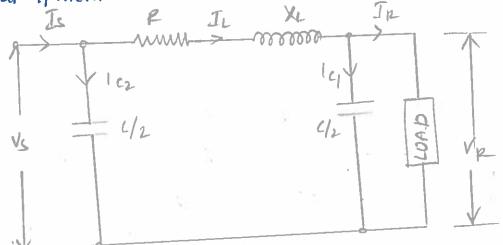
$$Dm^{2} \sqrt{4.17} X 4.17 X 4.9 = 4.4 m$$

$$Induction & [phase lm = 10^{-7} x 2 \log e 4.4] 0.195H$$

$$= 6.23 X 10^{-7} H = 0.625 X 10^{-3} mH$$

Inductance / phase / Km = 0. 623 × 10-3× 1000 = 0.623 mH.





In this method, capacitance of each conductor is divided into two halves : one half being lumped at the sending end and other half the receivery and.

Let
$$I_{R} = Load$$
 Current por phase
 $R = resistance$ por phase
 $X_{L} = inductive$ preactance per phase
 $C = Capacitanen$ per phase
 $Cos \neq p = receiving$ end power factor

Vs = sending and voltage per phase The phasor diagram for the Circuit VR = VR FD be Load current Te = 1 × (cos \$ pe - jsin \$ pe) Charging wovent at Tej=jω (c/2) Vp=jTfcVp load endis 1x Line current IL = TR + Tr. sending end voltage Vs = Ve + I2 = Ve + IL (P-+ jal) charging current at the sending ends $\overline{ic_2} = j\omega(Ch)\overline{v_3} = j\overline{\eta}fc\overline{v_3}$: sendors Is = IL + IL & Givon data. 13. RIL X1/2 X12 S P/2

Total substance /phase
$$f = 0.1 \times 100 = 10 \pi$$
.
Total substance /phase $\chi = 0.2 \times 100 = 20 \pi$.
Capacitius supceptance $\chi = 0.04 \times 100^{-4} \times 100 = 4 \times 10^{-9}$.
Receiving and Voltage/phase $\chi = 0.04 \times 10^{-4} \times 100 = 4 \times 10^{-9}$.
Load Current $Ip = 10,000 \times 10^{3} = 38105 \vee$
Load Current $Ip = 10,000 \times 10^{3} = 1098$
 $\sqrt{3} \times 66 \times 10^{3} \times 0.3$
Cosope = 0.8; Sin $p_{L} = 0.6$
Impedance per phan $\Xi = P+j \chi_{L} = 10+j20$
(i) Taking receiving end voltage as the reference phasor
Receiving end voltage $\nabla p = Vp + j0 = 38,105 \vee$
Load Current $Ip = Ip (\cos p_{R} - j\sin q_{R}) = 109 (0.8-j0.6)$
Voltage across C , $\overline{V_{1}} = \overline{V_{R}} + \overline{T_{R}} + \overline{Z_{12}} = 39,105 + (87 \cdot 1-j6)$
Charging Grownel $\overline{I_{R}} = 39,105 + 436 + 597 + 545$
Charging Grownel $\overline{I_{S}} = \overline{I_{R}} + \overline{I_{S}} = 39,105 + (39,195 + 545)$
Sending and Current $\overline{I_{S}} = \overline{I_{R}} + 1C = (87 \cdot 2) - j65 \cdot 4)$
 $+ (-0.2 \times 8 + j15 \cdot 6)$
 $= 87 \cdot 0 - j49 \cdot 8 = 100 L - 29' \cdot 49/A$
 ε . Dending end Current = 100A

(ii) Sending end voltage,
$$V_{i} = V_{i} + V_{i} \neq L_{i} = (39.195+1545)$$

 $+ (87.0 - j.49.8) (6+j.10)$
 $= 39,195 + j.545 + 434.9 + j.870 - j.249 +$
 $+ 49.8$
 $= 40128 + j.1170 = 40145 U'40'V$
 $a, Line value of sending end voltage
 $a = 40145 \times \sqrt{2} = 695 23 \text{ KeV}$
(iii) Referening to phason diagram
 $\Theta_{1} = 200 \text{ gle}$ between V_{P} and $V_{S} = 1^{a}40'$
 $\Theta_{2} = 200 \text{ gle}$ V_{P} and $V_{S} = 1^{a}40'$
 $\Theta_{2} = 200 \text{ gle}$ V_{P} and $V_{S} = 1^{a}40'$
 $\Theta_{2} = 200 \text{ gle}$ V_{V} and $V_{S} = 1^{a}40'$
 $\Theta_{2} = 200 \text{ gle}$ V_{V} and $V_{S} = 1^{a}40'$
 $\Theta_{2} = 200 \text{ gle}$ V_{V} and $V_{S} = 1^{a}40'$
 V_{P} and $V_{S} = 1^{a}40'$
 $V_{I} = 3127$
 \therefore sending end power -factor, $U_{S} = U_{S} = 1277$
 \therefore sending end power -factor, $U_{S} = U_{S} = 1277$
 $= 0.853 \text{ lag}$
(U) Sending end power = $3V_{S}$ (set $\Theta_{S} = 3X + 0,145 \times 100 \times 0.833$
 $= 102.73105 \text{ W} = 10173.105 \text{ KW}$
Power delivered $= 10,1000 \text{ KW}$
 $a, Transmission Efficient $= \frac{10,1000}{10773,105} \times 100 = 97.347$ M
14. Given data
 $3phate$
 $V_{I} = 8 \text{ KV} = 8 \text{ K}10^{3}$
 $V_{2} = 11 \text{ KV} = 11 \text{ K}10^{3}$
 $V_{2} = 11 \text{ KV} = 11 \text{ K}10^{3}$
 $V_{2} = 11 \text{ KV} = 3 \text{ dics}.$$$

+091mula $V_2 = (1 + 1/m) v_1 =$

.

(a) Pin type insulations: The part section of a pin type 15 insulator, the Pin type insulator is serviced pole . The is a goloove on the Goove to upor end of the insulator for ductor - Shed housing the Conductor. The De Conduct Cementing -tot passes through this groove and is bound by the annetalet wire of the Same material as the Condutor Galvonised Pin type insulators are Steel pin used for transmission and distribution of electric power at Finsulato Voltage upto 33 kv **(b)** Supension type insulators. insulator increases rapidly as the conducting voltage is increased. Therefore this type of ingulator is not economical beyon ds 33 Ev. The Consists of the content of the consists of the content of the consists of the content of the consists of the consis The Consists of a number of porcelair dies Connected in services by metal links in the form of string. The Conductor is Suppended at the bottom end of this storing while . the other end of the String of is secured to the cross-orm of tower. Advantage Suppension type insulation are cheaper. is if any one dics is damaged the whole String does not become use less because the st isforctory to supply the goreaber demand . I lessibily to the line

12(2)
$$\frac{1}{2}$$
 $\frac{1}{2}$ $\frac{1}{2}$

Nedimfall Selvanetayene Relietins Interon Peenholoev (Antonomous). 10/4/29/00/2019 Manapement System

NSRIT

	Semester End Suppleme	ntary Examinati	on April/A	day 2022		
Degree	B. Tech. (U. G.) Program	ECE	on Abrin	Academic Year	2021 - 2022	
Course	–	3 Hrs. Max. M	Marks 70		611	
Course	Digital System Design					
Part A (Short Answer Questions 5 x 2 = 10 Marks)					
No.	Questions (1 through 5)			Learning Outcome (s) DoK	
1	(11x1y)8 =(12c9)16 find x and y values			20EC305.1	L1	
2	Implement F= AB+ B'C using logic gates.			20EC305.2	L1	
3	Differentiate between RAM and ROM.	-100		20EC305.3	L1	
4 5	State the difference between Latch and Flip- What is the basic use of EDA tools?	пор.		20EC305.4 20EC305.5	L1 L1	
	Long Answer Questions 5 x 12 = 60 Marks)			2020303.3	LI	
No.	Questions (6 through 15)		Marks	Learning Outcome	(s) DoK	
	Convert The Following. i. (A6.CE)16 = ()10 ii. (1266.756)8 = ()10 iii.				
6 (a)	(10100011.11001)2 = ()10 iv (179.897) 10:		6M	20EC305.1	L2	
	(()				
6 (b)	Represent AND, OR, NOT, NAND gate using	y NOR gate	6M	20EC305.1	L2	
		OR				
7 (a)	Determine the single error correct code for the	ne information	6M	20EC305.1	L2	
	code 10111 for odd parity (i) Convert the following binary 1010011 into	orav code				
7 (b)	(ii)Convert the following gray code 101011 in		6M	20EC305.1	L2	
	equivalent binary					
8 (a)	Demonstrate by means of truth tables the Bo Associative law and distributive law	polean	6M	20EC305.2	L2	
	Simplify the following Boolean expression to	a minimum				
8 (b)	number of literals. i. F=(BC1+A1D)(6M	20EC305.2	L2	
	ii. F=WYZ+XY+XZ1+YZ.					
	Obtain the simplified supression is sup of a	OR Industa using K				
9 (a)	Obtain the simplified expression in sum of paramethod $F(A, B, C, D) = \Sigma(0, 1, 4, 5, 7)$		6M	20EC305.2	L2	
0.45	Obtain the simplified expression in POS (pro					
9 (b)	F(w, x,y,z)= Σ (1,2,4,7,12,14,15) using K-ma		6M	20EC305.2	L2	
10 (a)	State the applications of demultiplexer and	implement 1 to 4	6M	20EC305.3	L3	
10 (b)	De-Mux. Construct 4*16 decoder using two 3*8 deco	dom	6M	2050205.2	L3	
10(0)	Construct 4 To decoder using two 5 8 deco	OR	VIV	20EC305.3	LJ	
11 (a)	Build a 2 Bit Magnitude Comparator using g		6M	20EC305.3	L3	
11 (b)	Develop the following Boolean functions usi	ng PLA with 3	6M	20EC305.3	L3	
11(0)	AND gates. F1 (ABC) = Σ (3,5,7), F2 = Σ (4,5)	5,7).	UM	2020303.3	LJ	
	Mention the drawback of IV fin floor and av	nloin in whet we		2050205 4	10	
12 (a)	Mention the drawback of JK flip- flop and ex the drawback is eliminated in Master -Slave		6M	20EC305.4	L2	
	Draw the logic diagram of a SR flip flop usin		6M	20EC305.4	L2	
12 (b)	and illustrate its operation using truth table a					
	expression.		(Controller of Eka	Chui	
		OR		Nepit //	minations	
1000	en an state and state			NSRIT (A Visakhapatr	y set ex o i e i e	1
		8 2		roaniapatr	lam	

13 (a)	Show the circuit diagram of 4-bit Johnson counter using D-flip flop and explain its operation with the help of bit pattern	6M	20EC305.4	L2
13 (b)	Draw the block diagram of asynchronous sequential circuit	6M	20EC305.4	L2
14 (a)	Contrast the program structure of VHDL and Explain the significance of entity and architecture	6M	20EC305.5	L4
14 (b)	Inspect a VHDL code for 4:1 MUX using Case Statement OR	6M	20EC305.5	L4
15 (a)	Analyze the dataflow design style of VHDL programming with suitable example	6M	20EC305.5	L4
15 (b)	List out the detail about packages and libraries used in VHDL	6M =	20EC305.5	L4

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G.Kolyouri Controller of Examinations NSRIT (A) Visakhapatnam

5) What is the basic use of EDA tools, (2M) Sol:- EAD expands to Electronic Design Automation and these tools are used for Synthesis, Implemention and simulation of Electronic Concurts on the Software itself.

Part B (Long Answer Questions
$$5X12 = 60 \text{ M}$$
)
6(a) Convert the following (6H)
(i) (A6.CE)₁₆ = ()₁₀ (0.5 M)
Sol:- A 6.CE₁₆ = AX16¹ + 6X16[°] + CX16⁻¹ + EX16⁻¹
= 10X16¹ + 6X16[°] + C(12)X16⁻¹ + (14) EX16⁻¹
= 160 + 6 + 0.75 + 0.05468
= 166.80468₍₁₀₎
...(A6.CE)₁₆ = (166.80468)₁₀

(ii)
$$(1266.756)_{q} = ()_{10} - (999)(1.5 M)$$

Set: $(266.756)_{q} = ()_{10} - (999)(1.5 M)$
 $5et: (266.756)_{q} = ()_{10} - (099)(1.5 M)$
 $+ 5 \times 9^{2} + 6 \times 8^{2}$
 $= 512 + 128 + 48 + 6 + 0.875 + 0.07912 + 0.0117$
 $= (6.944.9649)_{10}$
(iii) $(10100011.11001)_{2} = ()_{10} - (1.5 M)$
Set: $(266.756)_{9} = (694.9645)_{10}$
(iii) $(10100011.1100)_{2} = ()_{10} - (1.5 M)$
Set: (287033007733355)
 $= 128703264 + 123^{2} + 023^{2} + 023^{2} + 023^{2} + 023^{2} + 023^{2} + 023^{2} + 023^{2} + 023^{2} + 023^{2} + 023^{2} + 023^{2} + 023^{2} + 023^{2} + 023^{2} + 023^{2} + 12$

: 0.897,= E5AICIL - 179.897,= B0.E5AICIL

6(b) Represent AND, OR, NOT, NAND gale Using NON July. Soli- AND Gate $A = D = \overline{A} = \overline{A} = \overline{B}$ $Y = A \cdot B = \overline{B} = \overline{B} = \overline{A} \cdot \overline{B}$ (1.5 M) $B = \overline{D} = \overline{A} \cdot \overline{B} = \overline{A} \cdot \overline{B}$ OR Gate Y=A+B A DATB D Y=A+B (1.5M) B D D A+B (1.5M) NOT Gake X_EDa-x+x = X NAND Gales $Y = \overline{AB}$ $A = \overline{D} = \overline{\overline{A} + \overline{B}}$ $\overline{A + \overline{B}}$ $\overline{Y - \overline{A} + \overline{B}}$ (1.5 M) $B = \overline{D} = \overline{D} = \overline{D} = \overline{D} = \overline{A + \overline{B}}$ (1.5 M)FCa) Retermine the stongle error correct code for the information Code 10111 for odd party (GM) iol- Guiven Information Code = 10111 No of Information bits x=5 No of practity bits P=? 2° 7 x+p+1 Let p=3 2375+3+1 [87,9] Not Natisfied et P=0, 2° 7, 5+0+1 \$ 17,6 Not Let P=4 Sat is feed 247,5+4+1 Lat P=1, 2'3,5+1+1 [167,10] Saturfut 277 Not Satisfeed · No of poulty bills = 4 Let p=2, 22 > 5-+2-+1 47,8 Nost blesfie .

A. .. Irlains

(2) S(A) Demonstrate by mound of touth bables the Boolean elisiocallance date and Destructure have - (6M)

Taur: A(B+C) - AB+AC

- 30 A J - 1					[AB+AC]
4 5 -	RIC	A(R+C)	A B	AC	1.31
ABL		1 1	0	6	C
0 0 0	0	C	0	0	- C
0 0 1		0			-12
010		- Č	<i>i</i>	0	
	- r		0	C	- Q
0 1 1			0	0	0
I D O	0	- C2		i	
(0)			0	0	1
110			1	™eer?	
		1	1	1	
+ + +	1 4				

Sible Simplify the following Boaldan organization to a menimum. Number of leterals (i) F = (BC-1A'D) (AB-1CD')

(11) F = WYZ + XY + XZ' + YZ (3.4 M) = Y(WZ + X) + XZ' + YZ [.: A) + BL = (A + B) (A+C) = V(WZ + X) + XZ' + YZ [.: A) + BL = (A + B) (A+C)

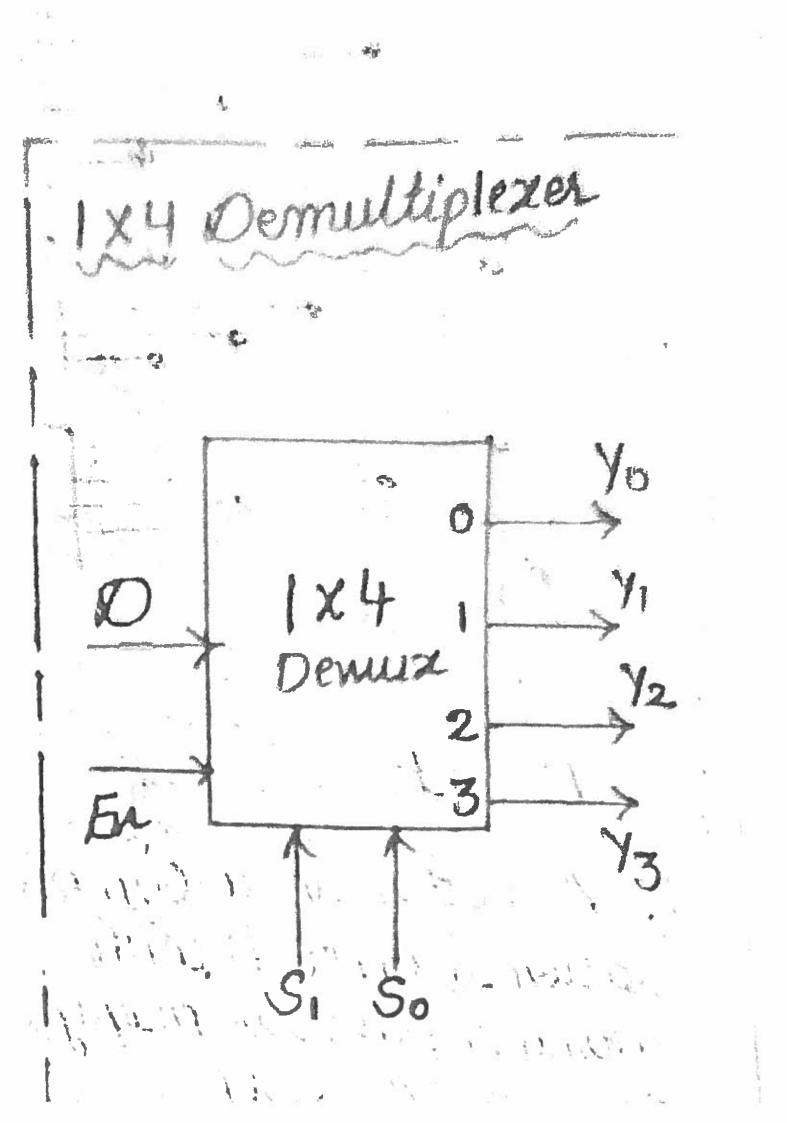
$$= Y[(x+w)(x+z_1) + x_2 + x_w + w_2)] + x_2' + 12$$

$$= Y [(X + Xz + Xw + wz)] + Xz' + Yz$$

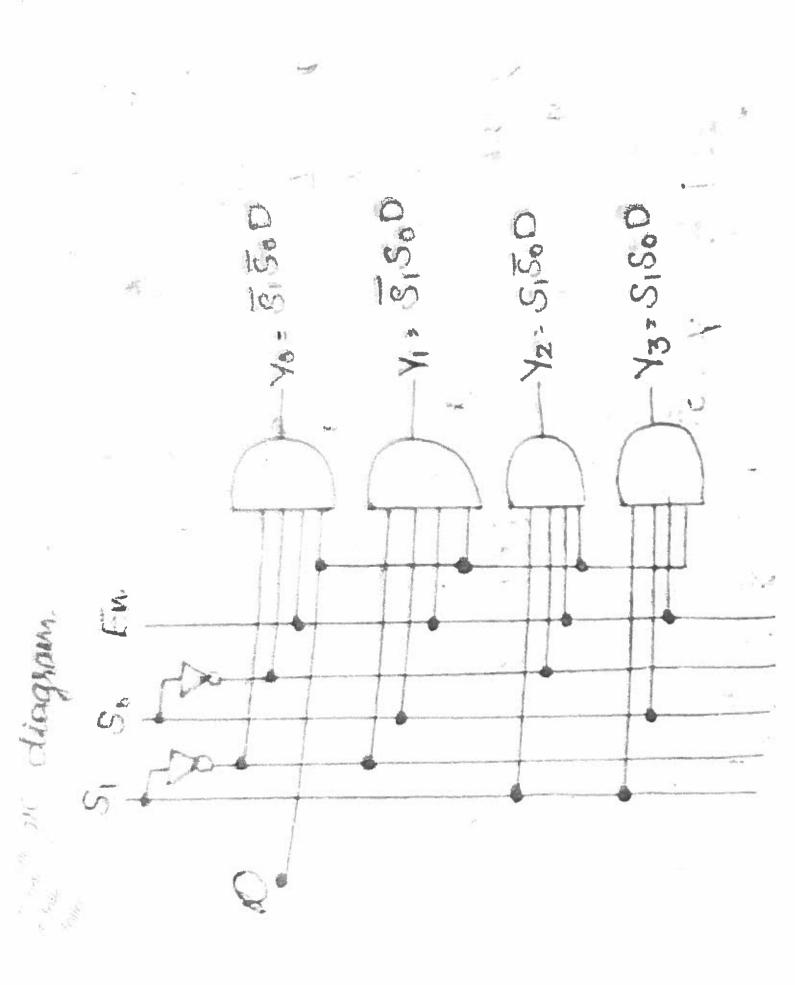
 $= \frac{x7(1+w)+7w2+x2'+72}{x7+7w2+x2'+72}$

(.s.m).

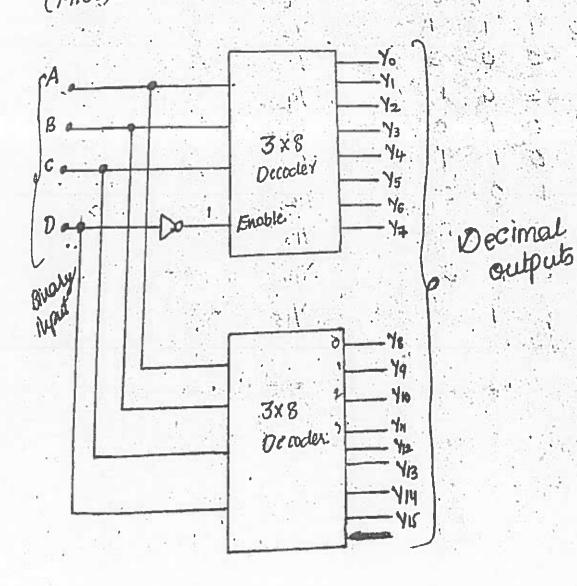
3) derea of Application is Communication System. 1) It is weed to connect a dingle dounce la number. 10 (a) State the application of demutuplexes and - GM 3) Secural to Parallel Conventer Luplement 1 to 4 De-mux Set: - Applications: destendation.



un an X4 clum 1.00 5 ole'. 00 3 1 Jhe 7

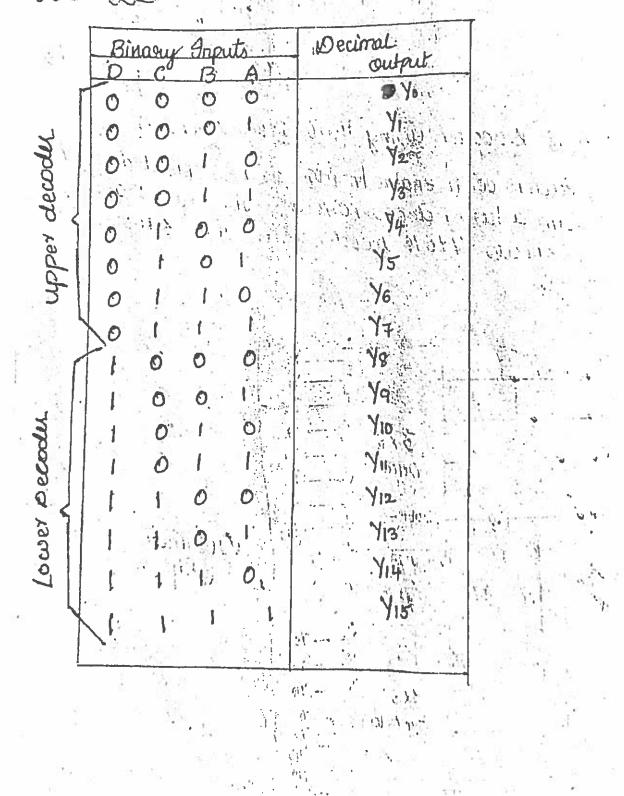


10(b) 4 to 16 Decoder using Two 3 to 8 Decoders (GM) Decoders with enable inputs can be connected together to form a larger decoder circuit. The circuit arrongement shows 4 to 16 decoder using two 3 to 8 decoder (74138).

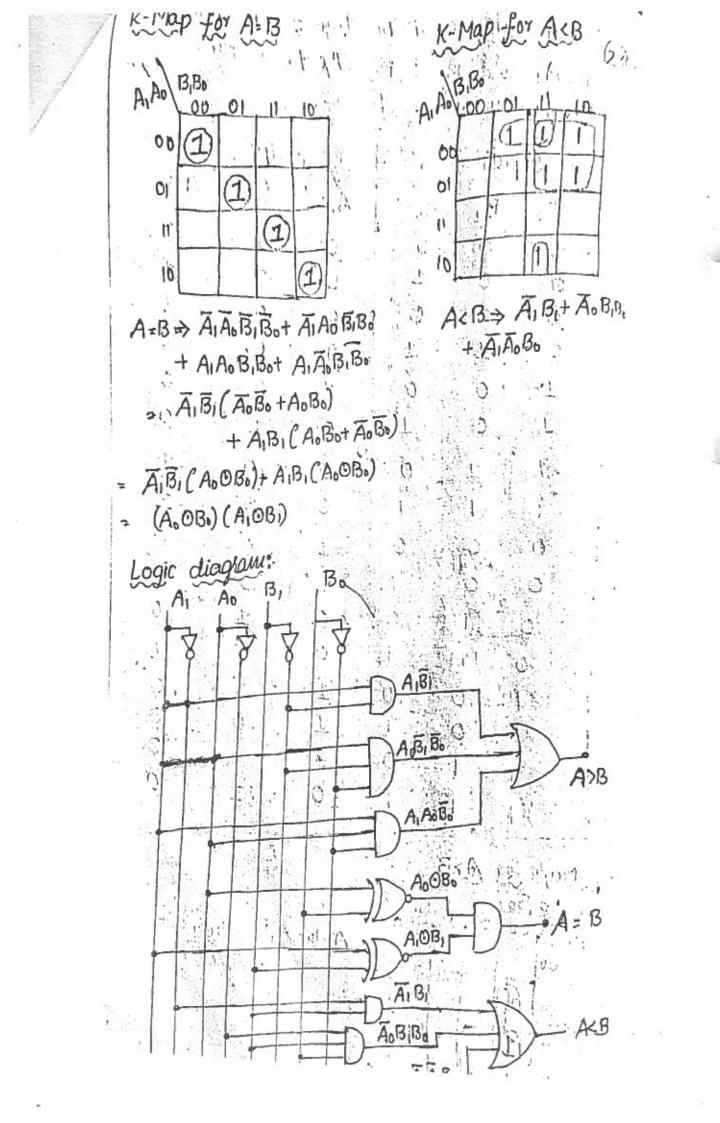


Through an inverter E to the upper decoder (Bby, (3)) Through an inverter E to the upper decoder (Bby, (3)) to Y=) and directly to E to the lower decoder. (Ys to Y15). Jhus when D=0 'Llow, the upper decoder is enabled and lower decoder is disabled. When D=1 (HIGH), the lower decoder is enabled and upper decoder is disabled.

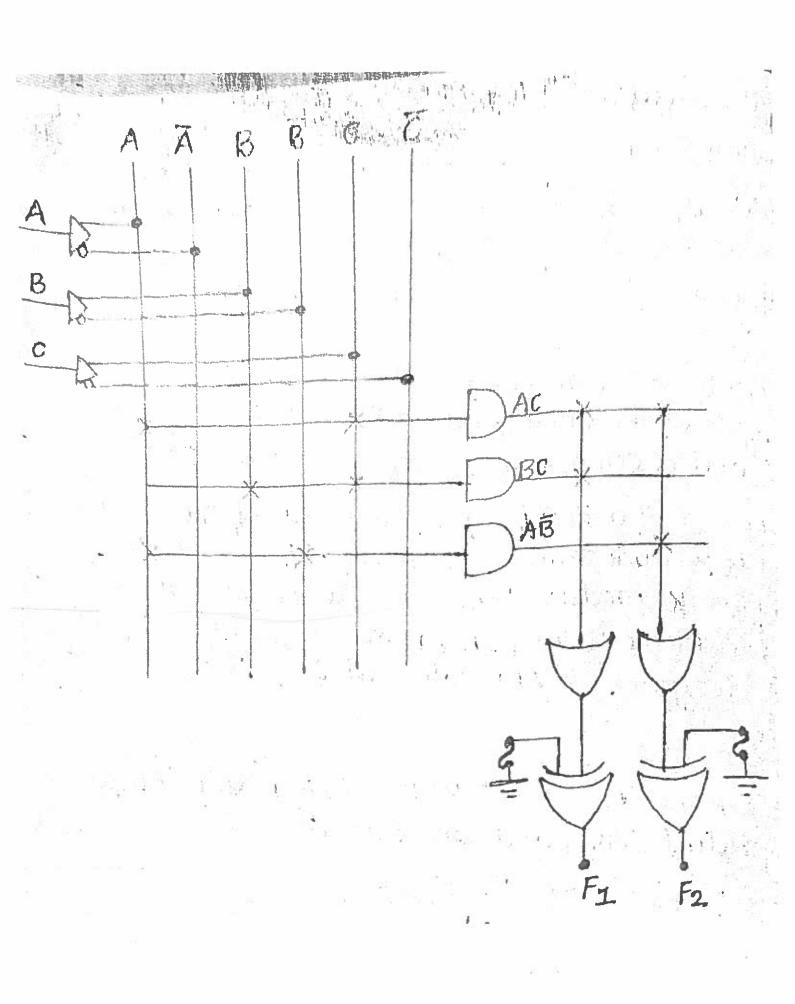
Truth Jable.



be A = A1 A0 & B: 13, B0. The truth table for 2-bit community ||(a)|GH 2-bit comparator is as follows. Outputs 1. 4 Inputs ALB A=B A>B. Ay A. B₁ B. 0 1 0 0 Ð 0 0 1 OH! 1 0 0 0 0 1 Ö O I 1:3: 0 Qui 0 J. + • t. . 1 0 1 1 0 6 divit ha chu 0 1. 0 0 17-Os: 1 1 0 0 0,. 1 1 0. .0.,0 0 O. 1 Line V ŀ 0 1 0 0 0 1 1 0 1 0 0 0 0 Ģ O 1 1. 0 1 Ø 0 C) 1 0 0 0 0 1 Ir. \mathbf{O} 1 5 1 K-Map for A>B ь÷ B,Bo A, Ao 101 10 14 0D A>B= A,B, + A0B, Bo+ 90 A, AOB. ABIBOL 01 ŧ/ 0



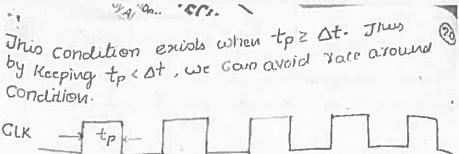
$$(b) \begin{array}{l} (b) \\ (c) \quad \text{Supplement the circuit coith a PLA having 3 inpuls,} \\ (c) \quad \text{Sproduct terms and two outputs} \\ (c) \quad \text{Set: } K \cdot Map for F_1 \\ (c) \quad \text{Set: } K \cdot Map for F_1 \\ (c) \quad \text{Set: } K \cdot Map for F_2 \\ (c) \quad \text{True form} \\ (c) \quad \text{Set: } K \cdot Map for F_2 \\ (c) \quad \text{True form} \\ (c) \quad \text{Set: } K \cdot Map for F_2 \\ (c) \quad \text{True form} \\ (c) \quad \text{Set: } K \cdot Map for F_2 \\ (c) \quad \text{True form} \\ (c) \quad \text{Set: } K \cdot Map for F_2 \\ (c) \quad \text{True form} \\ (c) \quad \text{Set: } K \cdot Map for F_2 \\ (c) \quad \text{True form} \\ (c) \quad \text{Set: } K \cdot Map for F_2 \\ (c) \quad \text{Set: } K$$

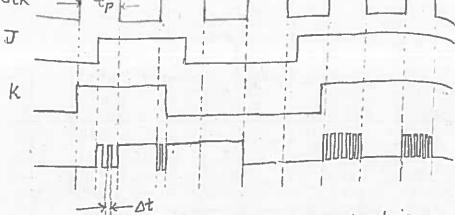


12(0)

Pace-Around Condition:-In JK Flip-flop, consider the assignment of excitations In JK Flip-flop, consider the assignment of excitations J:k: 1. If the width of the clock pulse "tp' is too long the state of the flip-flop cuill keep on changing from the state of the flip-flop cuill keep on changing from the state of the flip-flop cuill keep on changing from the state of the flip-flop cuill keep on changing from the state of the flip-flop cuill keep on changing from the state of the flip-flop cuill be unreviain. This of the clock fulse, its state cuill be unreviain. This Phenomonon is known on flace ground condition.

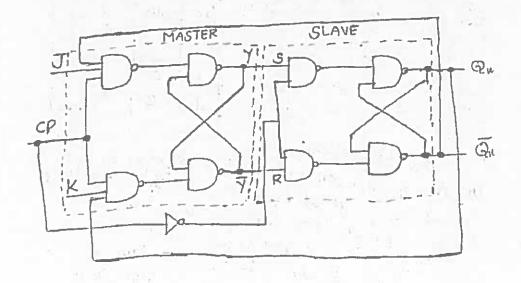
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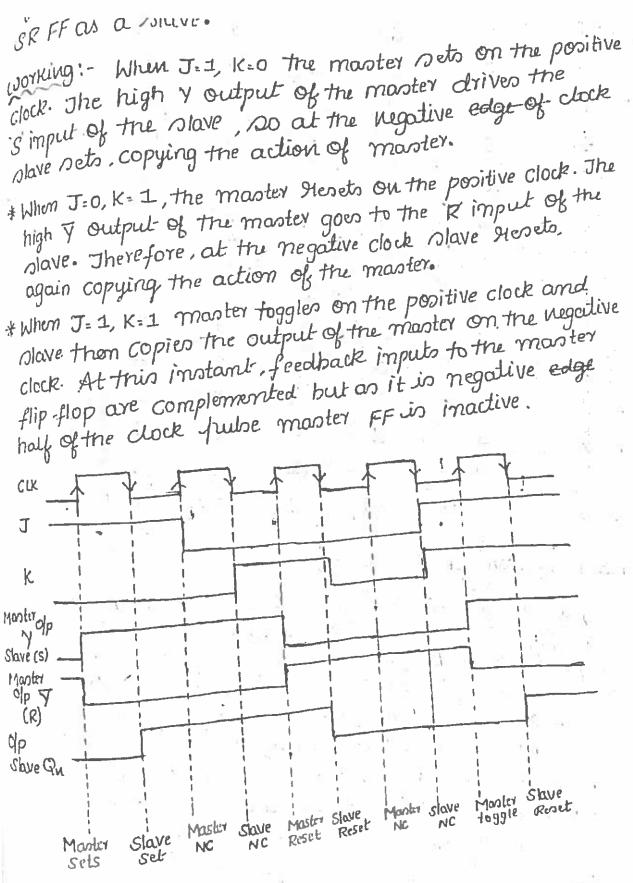




A more producal method for overcoming this difficulty is the use of Master-Slave Configuration

Master-Slave JK Flip-flop:-





S

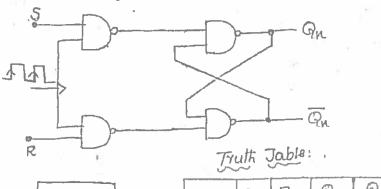
-	1		_	_				. 2				1				
- Annual sectors								0								
7	0	NC	=	NG	0	NC	0	NC	Н	NC	F	NC	H	NG	0	NC
Qu.	0	0	Ч	- 	0	g	4	Ч	0	0					स •	Н
×	0	0	0	0	Н	Ч	4	-1	0	0	0	0	H	-	+	- 4
Þ	0	0	0	0	0	0	0	0	4	Ч	Ч	H	Ц.	Н	4	H
Crk	4	\rightarrow	4	->	4	->	. 4	. →	4	*	Ł	4	4	4.3	•	>

12(b) SR Flip-Flop:

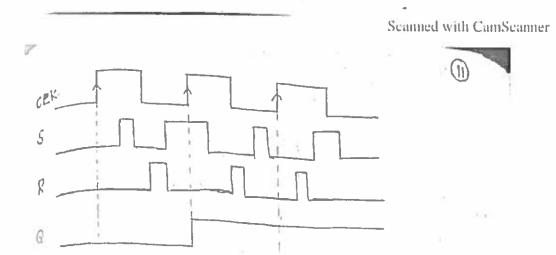
(611)

The following figure shows the positive edge triggened clocked SR FF. The circuit is similar to the SR latch except enable signal is replaced by the clock Pulse CCLK). The circuit output suspends to the S'd R'imputs only at the positive edges of the clock R'imputs only at the positive edges of the clock Pulse. At any other inclants of time, the SR flip-flop Pulse. At any other inclants of time, the SR flip-flop

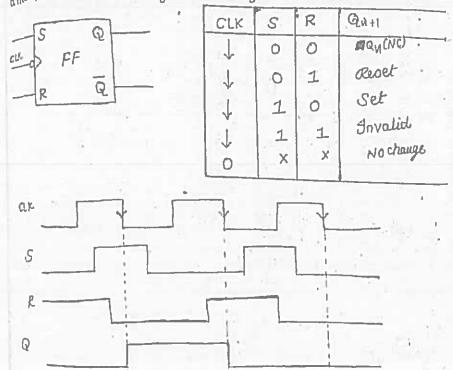
Fig: Positive edge Jriggered Flip-flop



	S	Q]	CLK	S	R	Qn.	Qutt	State
GLK	↓ F	F	(3)	1	0	0	0	0	NO
X	P2	র		1	0	0	1	1	chauge
	L			T	0	ļ	0	0	Resel
GIK	S	R	Quil	1	0	1	1	0	
1	0	0	No Chauge	T	1	0	0	1	Set
1	0	1	Reset	1		0	1	1	
\uparrow	1	0	Set	Ť			0	X	And
Ť	1	12.	Invalid	1	L	2	1	X	Invalia
0	X	× -	NO chaige	D	Х	X	0	0	NO.
	·			0	х	X	1	1	charge

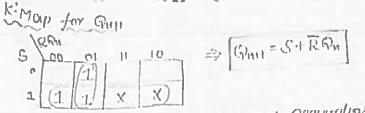


In negative edge higgered SR FF the circuit output suspends the negative edges of the clock fulse. The logic symbol and the truth table of the -ve edge triggered SR flip - flop.



Characteristic Equation of SR Elip-flop: The characteristic equation of a flip-flop is the Equation expressing the next state of a flip-flop in terms of its present state and present excitations.

Scanned with CamScanner



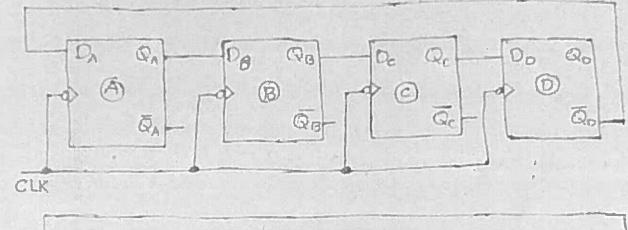
Excitation Jable: For the design of Dequarties circuit we phould know the excitation takke of flip-flops. The excitation takke of the flip-flop corn be obtained from its truth take. It inclicates the impulse required to be applied to the flip-flop to take it from present

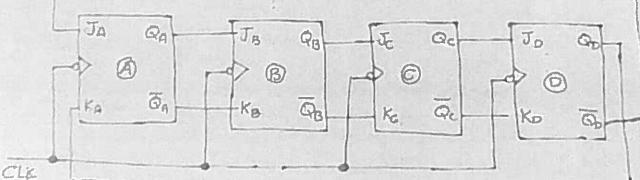
State to word state. Excitation Juble of SR flip-flop is as follows:

PS	NS	Requir	ed input
Qn	Quit	S	R
0	0	0	X
0	1	1	D
1	0	0	-
1	-1	X,	0,

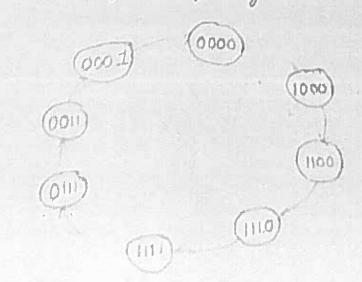
13(a)

JOHNSON COUNTER (THISTED RING COUNTER) (M) ba This counter is obtained from a ring counter. by providing feedback from the inverted output of the last FF to the D input of the first FF: # The Q input of each stage to connected to the is imput of the mexit stage. but the Q output of the last stage to connected to the Winput of the last stage to connected to the Winput of the last stage. Inst the Winput of the last stage to connected to the Winput of the last stage, therefore the Warne is twisted Counter.





The state diagram of a johnson counter.



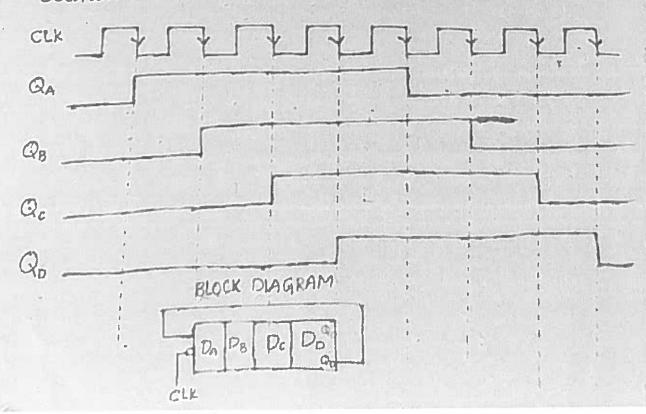
Sequence Jable.

QB Q. QD After CLX pulsos QA . Ø 0. I ユ 'n 上 1.

(59)

Let the initially all the FF's are Reset, i.e., the state of the counter is 0000. After each clock pulse, the level of Q, is shifted to QB, the level of QB is shifted to Qc. The level of Qc is shifted to QO, the level of Q is shifted to QA.

* A Johnson Counter of N-FF's cuill have 2n unique states and count upts 2n pulsos. So it is a MOD-2n Counter. It is more economical than suing Counter but less economical than supple counter



13(6) Asynchronous counteus:-(CM) Jhree-bit Ripple up. counter woing Negative edge + The three-bit up-counter consists of 3 FF's and this counter counts in the order 0,1,2,3,4,5,6,7,0 The external CLK pulse is applied to one of the FF and its output is considered as the LSB. logic 1 logic 1 g legic 1 JB CLK Qn KA a) Logic diagram The tirning diagram for the 3 bit up counter Wing - ye edge triggered FP's is as follows:

3 3 0 3 6 D QA QB Qc pecimal CLOCK pube Qc QB QA To. excossis 1 0 Initially 0 0 0 1 18t (1) 1 0 0 2nd (4) 2 0 0 1 3rd (4) 3 1 1 0 4 4th (1) 1 0 0 5 5th (4). 0 1 1 6 1 0, 6th (1) 1 7 1. 1 1 -jt (1). 0 0 8th (4) 0 0

14) a) Structure of VHDL Program :

(143)

Every VHDL program consists of at least one entity/architecture pair. In a large design, you will typically write many entity/architecture pairs and connect them together to form a complete circuit. An entity declaration describes the circuit as it appears from the "outside" - from the perspective of its input and output interfaces. The second part of a minimal VHDL design description is the architecture declaration.

ENTITY :

1. 1. 1.

ENTITY is the list with specifications of all input and output pins of the circuit. Its syntax is shown below :

ENTITY name IS PORT (port_name : signal_mode signal_type; port_name : signal_mode signal_type; ...); END name;

The mode of the signal used may be IN, OUT, INOUT or BUFFER. IN and OUT are unidirectional pins, while INOUT is bidirectional. BUFFER, is used when the output signal is used internally in the design. The type of the signal may be BIT, STD_LOGIC, INTEGER, etc. The name of the entity should be not use the VHDL reserved words.

Let us consider the NAND gate, Its ENTITY may be declared as :

ENTITY nand_Two IS

PORT (a, b : IN BIT;

x : OUT BIT

);

END nand_Two;

ARCHITECTURE :

The ARCHITECTURE is the description of how the circuit of design works. Its syntax is as follows,

ARCHITECTURE architecture_name OF entity_name IS

[declarations]

BEGIN

(code)

END architecture_name;

As shown in above syntax, architecture has two parts, 1) A declarative part where signals and constants are declared and 2) The code part. The name of architecture is any name except VHDL reserved words.

14 b)Design of 4 to 1 Multiplexer using CASE Statement (Behavior Modeling Style)

(би)

1 1 1

library IEEE; use IEEE.STD_LOGIC_1164.all;

entity multiplexer_case is
 port(
 din : in STD_LOGIC_VECTOR(3 downto 0);
 sel : in STD_LOGIC_VECTOR(1 downto 0);
 dout : out STD_LOGIC
);
end multiplexer_case;

architecture multiplexer_case_arc of multiplexer_case is begin

mux : process (din,sel) is begin case sel is

```
when "00" => dout <= din(3);
when "01" => dout <= din(2);
when "10" => dout <= din(1);
when others => dout <= din(0);
end case;
end process mux;
```

end multiplexer_case_arc;

6.2.3

15 .b)Explain about Package in VHDL?

• Groups of procedures and functions that are related can be aggregated into a module that is called package.

(GM)

• A package can be shared across many VHDL models (type definitions, functions, procedures, etc)

A package can also contains user defined data types and constants.

• When working with a large design project consisting of many small VHDL programs, it is convenient to have common procedures and functions in separate packages.

Package declaration

It contains interface or specifications of the functions and procedures that are available in the package. It gives idea about – list of functions and procedures, the input parameters, the type of input parameters,

the output value, the type of the output value etc.

Package declaration Syntax:

packagepackage-name is

package-item-declarations_ these may be:

subprogram declarations / type declarations / subtype declarations / constant declarations /

signal declarations / variable declarations / file declarations / alias declarations / component declaration / attribute declarations / attribute specifications /disconnection specifications

end packagepackage-name;

Package Body

It basically contains the code that implements the subprograms

Package Body Syntax:

package body package-name is

subprogram bodies / complete constant declarations / subprogram declarations / type and subtype declarations / file and alias declarations / use clauses

end package bodypackage-name;

Library in VHDL

Design units such as packages, architectures and Entities can be compiled into a library. Libraries are generally implemented as directories on a computing system.

Library is referenced by its logical name and there exists a mapping of logical name into the physical name of the system (which is the directory names on the computing system).

The mapping is maintained by the system.

Just like variables and signals, library must be declared before we can use it.

Syntax for Library declaration:

LIBRARY logical_library_name_1, logical_library_name2, ...;

In the VHDL language, the libraries STD and WORK are implicitly declared in the source code. Library STD contains the standard packages with VHDL distribution.

The WORK library refers to the current working directory.

There are other different libraries used for different tools used.

Other libraries such as math and other misc. are often supplied with the tool

15) a)Explain Data flow modeling in VHDL. (G M) In this modeling style the flow of data is monitored from input to output inside an entity. Data Flow Modeling Style works on Concurrent Executions. Concurrent Signal Assignment (<=), With Select Assignment, When Else Assignment

Generate Expressions are generally used in Data Flow Modeling Style. Programs:

VHDL code for AND-OR-INVERTER (AOI) in dataflow modeling style?

Library IEEE; Use IEEE.STD_LOGIC_1164.all; Entity AOI is Port(A,B,C,D: std_logic; Y: out_logic); End AOI; Architecture dataflow1 of AOI is Begin Y=not((A and B) or (c and D));

End dataflow1;

nall Safyana para kajo nali llegi Technology (Autoponous) (CAL Gueiny Menagemen System (CMS)

										115	RII
				ter End Supple	-					0004	0000
Degree		B. Tech.		Program		E (AI & ML) & C			lemic Year	2021 - 3	
Course C	Code	20CS305		Test Duration	3 Hrs.	Max. Mark	ks 70	Sem	ester		1
Course		COMPU	TER ORG	ANIZATION							
Part A (S No.	Questic	ons (1 thro	ugh 5)	x 2 = 10 Marks)		2		I	_earning Outo		DoK
1				y used binary co	de?				20CS30		L1
2		binary ad							20CS30		L1
3		control m							20CS30 20CS30		L1 L2
4		you hand							20CS30 20CS30		L2
5		OM an au			-1				200000	0.0	L. <i>6</i> .
		iswer Que ons (6 thro		x 12 = 60 Marks	21		N	Aarks	Learning Ou	tcome (s)	Do
No. 6 (a)				sentation of nur	hers in the	ALU	4.	6M	20CS3		Ľ
6 (a) 6 (b)				into octal, hexad		BCD.		6M	20CS3		L
7 (a)	Disting	uish hetw	een error	detection and co				6M	20CS30	05.1	L
				veen encoder an			ne role		20CS3		Ľ
7 (b)	of thes	e compon	ents in th	e design of comp	outers			6M	20053	UD. I	L
8 (a)	operat	ions.	·	n? Discuss the				5M	20CS3	05.2	Ľ
8 (b)				es of an instruction in the second seco			n case	7M	20CS3	05.2	L
	Mbat :	e register	transfor !	anguage? Desc		•	ised in	<u>ः</u>			
9 (a)		s register er transfer.		anyuayer Desc	une me ng	aio ayrribula u	1360-111	5M	20CS3	05.2	L
9 (b)				four memory re	ference inst	ructions.		7M	20CS3	105.2	L
10 (a)	organi	zation.		memory in a				8M	20053	305.3	L
10 (b)	What Discus desigr	ss relative	mean by e advanta	Complex Instr ges and disadv	uction Set antages of	Computer (such instruct	CISC)? ion set	4M	20053	305.3	L
	_				O	R				205.2	
11 (a)				er instruction for		alata at a	Harris .	5M	20CS:	305.3	l
				y Addressing	modes? Ex	oplain the fo	Dilowing				
11 (b)	addre	ssing moo	ies:	de ii) Immediate	Addressin	a mode iii) a	Relativo	7M	20CS	305.3	l
				et Addressing m		y mode my r	(CIQUVE		2		
		0	*	-					10000	100E 4	
12 (a)				ion algorithm.	۳۰ _ المام مام	ann ia a be	ordusere	6M	2005	\$305.4	I
12 (b)	imple		that pe	units and the erforms addition				6M	2003	305.4	
	-					R					
13 (a)				hich explains	multiplication	on of TWO	signed	6M	2003	\$305.4	
10 (0)		itude fixed			الأمسيانا	والأفار المحرولة الم	4 hi				
13 (b)				adding and sul numbers signed				6M	2003	5305.4	
14 (a)				e of cache men	nory and w	rite about dir	rect and	7M	200	5305.5	
	45500	iative ma				- 11		<i>1</i> 28.7	יירחני ו	\$305.5	
14 (b)	Distin	iguish bet	ween syn	chronous and as		s data transfer)R		5N	200	2003.3	
15 (a)	With	the help a	f a block i	diagram explain			roller.	6M	200	S305.5	
	What	is virtual	l memory	? With a neat b				6N	200	S305.5	
15 (b)	mem	ory addre	ss transla	tion.							~
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Visakhapatnam



N S RAJU INSTITUTE OF TECHNOLOGY

(AUTONOMOUS)

SONTYAM , ANANDAPURAM, VISAKHAPATNAM – 531 173

ANSWER KEY AND SCHEME OF EVALUATION

III Semester End Supplementary Exam

April 2022

Computer Organization (20CS305) (Common to CSE, CSE(AI &ML),CSE(DS))

Part A (Short Answer Question 5X2 =10 Marks)

1. What is the most commonly used binary code?

ASCII-1 M+ Explain purpose of ASCII - 1M

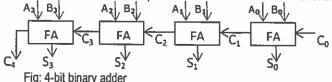
ASCII (American Standard Code for Information Interchange) is the most commonly used binary code. In an ASCII file, each alphabetic, numeric, or special character is represented with a 7-bit binary number (a string of seven 0s or 1s). ASCII was soon expanded to an 8-bit system that has 256 code points.

ASCII codes represent text in computers, telecommunications equipment, and other devices. Most modern character-encoding schemes are based on ASCII, although they support many additional characters.

2. What is binary adder?

Definition of binary adder with explanation and circuit-2 M

A binary adder is a digital circuit that produces the arithmetic sum of two binary numbers. It can be constructed with full adders connected in cascading form with the output carry from each full adder connected to the input carry of the next full adder in chain. Addition of n-bit numbers requires a chain of n full adders or chain of 1 half-adder and n-1 full adders.



3. What is control memory address?

Control memory address purpose-2 M

The control memory address register specifies the address of the microinstruction. The microinstruction contains a control word that specifies one or more microoperations for the data processor. The next address is computed in the next address generator also called microprogram sequencer. Typical functions of a microprogram sequencer are incrementing the control address register by one, loading into the control address register an address from control memory, transferring an external address, or loading an initial address to start the control operations.

4. How do you handle divide overflow?

Handling for divide overflow - 2 M

In some computers it is the responsibility of the programmers to check if DVF if set after each divide instruction. They can branch to a subroutine that takes a corrective measure. In some older computers divide overflow stopped the computer. But it is time consuming. The procedure in most computers is to provide an interrupt request when DVF is set. The interrupt causes the computer to suspend the current program and branch to a service routine to take a corrective measure. The most common corrective measure is to remove the program and type an error message explaining the reason why the program could not be completed. It is then the responsibility of the user who wrote the program to rescale the data or take any other corrective measure. The best way to avoid a divide overflow is to use floating point data.

1

5. Is CDROM an auxiliary memory?

CDROM is auxiliary memory and explanation of auxiliary memory-2M

highest-capacity and slowest-access storage in a computer system. In these memories, programs and data are kept for long-term storage or when not in immediate use. Examples of auxiliary memories are magnetic tapes, magnetic disks, floppy CD-ROM

Part B (Long Answer Questions 5X12 = 60 Marks)

6 (a) Explain floating point representation of numbers in the ALU

Scientific notation with example-3M+Normalized notation with example-3M

Floating Point Notation is a way to represent very large or very small numbers precisely using scientific notation in binary.

The floating-point notation has two types of notation

Scientific notation and Normalized notation

Scientific notation – Method of representing binary numbers into a x b^e form where a is real number, b is base and e is exponent. Scientific notation is further converted into floating-point notation because floating-point notation only accepts scientific notation. For example:-x = 376.4423 (its not scientific notation)

Number in scientific = 36.4423 x 10¹ or 3.64423 x 10²

To represent the number -(0.000000000000000000000101)₂ in floating point

Keeping only one digit to the left of the decimal point we get

Actual number	->	-	(0.00000000000000000000000000000000000
Scientific notation	->	-	1.01×2^{-24}

Normalized notation- To make the fixed part of the representation uniform, both the scientific method (for the decimal system) and the floating-point method (for the binary system) use only one non-zero digit on the left of the decimal point. This is called normalization. In the decimal system this digit can be 1 to 9, while in the binary system it can only be 1. In the following, d is a non-zero digit, x is a digit, and y is either 0 or 1.

Decimal	 ±	d.xxxxxxxxxxxxxx	Note: <i>d</i> is 1 to 9 and each <i>x</i> is 0 to 9
Binary	 ±	1.ууууууууууууууу	Note: each y is 0 or 1

0 cannot be represented or normalized because the representation set starts from 0.1, so how can we normalize zero. It's not possible. If the most significant bit of mantissa is a non zero, then such a representation is called normalized floating-point.

2

6 (b) Convert the number 7425 into octal, hexadecimal and BCD.
Octal conversion-2 M + Hexadecimal conversion of 7425 - 2 Marks + BCD Conversion-2M
Base for octal is 8, hexadecimal is 16.
Divide 7425 by the base number keeping notice of the quotient and the remainder.
Continue dividing the quotient by base until you get a quotient of zero.
Then just write out the remainders in the reverse order to get the equivalent of decimal number 7425 (i) Octal of 7425 is obtained by dividing 7425 with 8
7425 / 8 = 928 with remainder 1
928 / 8 = 116 with remainder 0
116 / 8 = 14 with remainder 4
14 / 8 = 1 with remainder 6

1 / 8 = 0 with remainder 1 Octal of 7425 is (16401)₈

(ii) Hexadecimal of 7425 is obtained by dividing 7425 with 16

7425 / 16 = 464 with remainder 1 464 / 16 = 29 with remainder 0 29 / 16 = 1 with remainder 13 (D in hexadecimal) 1 / 16 = 0 with remainder 1 Hexadecimal of 7425 is $(1D01)_{16}$ (iii) BCD of 7425 is 0111 0100 0010 0101

7 (a) Distinguish between error detection and correction codes

chor detection and endriconection unterences along with party bit, GRG etc-low

Error Detection is the detection of errors caused by noise or other impairments during transmission from the transmitter to the receiver. Error correction is the detection of errors and reconstruction of the original, error-free data.

Parity Bit: A parity bit is a bit that is added to a group of source bits to ensure that the number of set bits (i.e., bits with value 1) in the outcome is even or odd. It is a very simple scheme that can be used to detect single or any other odd number (i.e., three, five, etc.) of errors in the output. An even number of flipped bits will make the parity bit appear correct even though the data is erroneous.

For example, if each of a series of m-bit "words" has a parity bit added, showing whether there were an odd or even number of ones in that word, any word with a single error in it will be detected. It will not be known where in the word the error is, however. If, in addition, after each stream of n words a parity sum is sent, each bit of which shows whether there were an odd or even number of ones at that bit-position sent in the most recent group, the exact position of the error can be determined and the error corrected.

Any error-correcting code can be used for error detection. A code with minimum Hamming distance, d, can detect up to d – 1 errors in a code word. Using minimum-distance-based error-correcting codes for error detection can be suitable if a strict limit on the minimum number of errors to be detected is desired.

Codes with minimum Hamming distance d = 2 are degenerate cases of error-correcting codes, and can be used to detect single errors. The parity bit is an example of a single-error-detecting code.

7 (b) State the differences between encoder and multiplexer. Mention the role of these components in the design of computers Atleast 3 differences between encoder and multiplexer-4 M+Role of encoder and multiplexer in computer design-2M

Encoder	Multiplexer
An encoder is a combinational circuit that converts binary	
information from "n" input lines to a max of 2" unique output	several digital input signals and send those inputs to the final
lines. $n \rightarrow 2^n$	output. 2 " \rightarrow 1
The encoder does not have any selection input lines	A multiplexer of 2 ⁿ inputs allows n selection lines to select the
	number of input lines to send the final output
There is no need of select input line and limited data can be	The select lines determine which input is connected to the
sent thereby	output and also increase the amount of data that can be sent
	over a network within a certain time.

An encoder is a device that can be used to change a signal or data into a specific code. The encoder is used for compression of data. The encoder will convert the information from one format to another format i.e like electrical signals to counters. The feedback signal of the encoder will determine the position, count, speed and direction. The control devices are used to send the command to a particular function.

Examples of encoders are priority encoder, decimal to BCD encoder, octal to binary encoder. The code may be used for any purposes like, for compression of information required for transmission and storage.

In telephone networks, multiplexers are used where multiple audio signals are integrated on a single line of transmission. Multiplexers are also used as parallel to serial data converters.

A common application of multiplexing occurs when several embedded system devices share a single transmission line or bus line while communicating with the device. Each device in succession has a brief time to send and receive the data. 8 (a) What is a micro-operation? Discuss the four different types of microoperation.

Microoperation-1 M+Four types of microoperation-Arithmetic, Logical, Shift and Register transfer microoperation -1 Mark for each(4M)

Microoperation: The operations executed on data stored in registers is called microoperation.

Arithmetic microoperations- Perform arithmetic operations on numeric data in registers

The basic arithmetic microoperations are: addition, subtraction, increment and decrement.

 $R3 \leftarrow R1 + R2$. It states that contents of registers R1 and R2 are added and sum is transferred to register R3.

Subtraction Microoperation: R3 \leftarrow R1 – R2 or : R3 \leftarrow R1 + R2 + 1

Logical microoperations- Perform bit by bit operations on nonnumeric data in registers

Logic microoperations are: Clear, Exclusive Or, AND, OR, NAND etc

Shift microoperations are microoperations that are used for serial transfer of data. It shifts data either to left or right Example : Logical Shift, Arithmetic Shift, Circular Shift (either left or right)

Register Transfer : The symbolic notation used to describe the micro-operation transfers amongst registers is called Register transfer language. Example : R3 - R1 and Read : DR - M[AR], Write : M[AR] - R1

8 (b) Explain the different phases of an instruction cycle. What happens in case an instruction has memory operands? Fetch, Decode and Execute -> Instruction cycle-4M+If instruction has memory operands explain what happens-3M Different phases of instruction cycle are as follows:

1. Fetch the instruction from memory

2. Decode the instruction

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4. Execute the instruction.

Upon completion of step 4, the control goes back to step 1 to fetch, decode and execute the next instruction. This process

continues indefinitely unless a HALT instruction is encountered.

Fetch and Decode : Initially, the program counter PC is loaded with the address of the first instruction in the program. The

sequence counter SC is cleared to 0, providing a decoded timing signal T₀. After each clock pulse, SC is incremented by one,

so that the timing signals go through a sequence To, T1, T2, and so on. The microoperations for the fetch and decode phases can

be specified by the following register transfer statements.

T₀: AR←PC

 T_1 : IR \leftarrow M[AR], PC \leftarrow PC+1

 T_2 : D₀...D₇ \leftarrow Decode IR(12-14), AR \leftarrow IR(0-11), I \leftarrow IR(15)

Decoder output D₇, is equal to 1 if the operation code is equal to binary 111. If D₇ = 0 then it is a memory-reference instruction and if I =1 then it means it's a indirect address which is achieved using $AR \leftarrow M[AR]$.

This address is used during the memory read operation. The word at the address given by AR is read from memory and placed on the common bus. The LD input of AR is then enabled to receive the indirect address that resided in the 12 least significant bits of the memory word. The three instruction types are subdivided into four separate paths. The selected operation is activated with the clock transition associated with timing signal T₃. This can be symbolized as follows:

 $D_7 IT_3: AR \leftarrow M[AR]$

D₇ I'T₃: Nothing

When a memory-reference instruction with I = 0 is encountered, there is no need to do anything. The sequence counter SC must

be incremented when D₇'T₃, = 1, so that the execution of the memory-reference instruction can be continued with timing variable

T4.

9 (a) What is a register transfer language? Describe the basic symbols used in register transfer

Register transfer language – 2 M + Basic symbols used in register transfer – 3 M

Register Transfer Language: Shortly known as RTL. The symbolic notation used to describe the micro-operation transfers amongst registers is called Register transfer language.

Basic symbols used in register transfer are as follows :

Basic Symbols for Register Transfers

Symbol	Description	Examples			
Letters (and numerals)	Denotes a register	MAR, R2			
Parentheses () Arrow - Comma	Denotes a part of a register Denotes transfer of information Separates two microoperations	$R2(0-7), R2(L)$ $R2 \leftarrow R1$ $R2 \leftarrow R1, R1 \leftarrow R2$			

Example : $R2 \leftarrow R1$ means transfer the content of register R1 to register R2 Read : DR \leftarrow M[AR], Write : M[AR] \leftarrow R1

9 (b) Discuss with examples any four memory reference instructions

Mention any four memory-reference instructions among AND, ADD, LDA, STA, BUN, BSA and ISZ – 3M+ Example for 4 memory reference instructions-1 M for each example Total—4M

For example: AND-Performs AND logic operation on bits in AC and memory word given by the effective address

LDA-Load to AC-Transfers the memory word specified by the effective address to AC.

STA- Store Accumulator-Store the content of AC into the memory word specified by effective address. It is represented by D₃. AND requires operation decoder D₀ and is symbolically represented as

 $D_0T_4: DR \leftarrow M[AR]$

 D_0T_5 : AC \leftarrow AC \land M[AR], SC \leftarrow 0

ADD requires operation decoder D1 and is symbolically represented as

 $D_1T_4: DR \leftarrow M[AR]$

 $D_1T_5:AC \leftarrow AC + M[AR], E \leftarrow C_{out}, SC \leftarrow 0$

LDA requires operation decoder D_2 and is symbolically represented as

 D_2T_4 : DR \leftarrow M[AR]

U215. AU←UR, OU←U

STA requires just one microoperation since the output of AC is applied to the bus and the data input of memory is connected to the bus.

 $D_3T_4:M[AR] \leftarrow AC, SC \leftarrow 0$

10 (a) Explain about control memory in a micro programmed control organization.

Control memory definition -2M + Purpose of control memory-3M+Diagram of microprogrammed control organization-2M The control memory address register specifies the address of the microinstruction. The microinstruction contains a control word that specifies one or more microoperations for the data processor. The next address is computed in the next address generator also called microprogram sequencer.

Each word in control memory contains within it a microinstruction. The microinstruction specifies one or more microoperations for the system. A sequence of microinstructions constitutes a microprogram. Since alterations of the microprogram are not needed once the control unit is in operation, the control memory can be a read-only memory (ROM).

The content of the words in ROM are fixed and cannot be altered by simple programming since no writing capability is available in the ROM. ROM words are made permanent during the hardware production of the unit. The use of a microprogram involves placing all control variables in words of ROM for use by the control unit through successive read operations. The content of the word in ROM at a given address specifies a microinstruction. A more advanced development known as dynamic microprogramming permits a microprogram to be loaded initially from an auxiliary memory such as a magnetic disk. Control units that use dynamic microprogramming employ a writable control memory. This type of memory can be used for writing (to change the microprogram) but is used mostly for reading. A memory that is part of a control unit is referred to as a control memory. A computer that employs a microprogrammed control unit will have two separate memories: a main memory and a control memory. The main memory is available to the user for storing the programs. The contents of main memory may alter when the data are manipulated and every time that the program is changed. The user's program in main memory consists of machine instructions and data. In contrast, the control memory holds a fixed microprogram that cannot be altered by the occasional user. The microprogram consists of microinstructions that specify various internal control signals for execution of register microoperations. Each machine instruction initiates a series of microinstructions in control memory. These microinstructions generate the microoperations to fetch the instruction from main memory; to evaluate the effective address, to execute the operation specified by the instruction, and to return control to the fetch phase in order to repeat the cycle for the next instruction. The control memory is assumed to be a ROM, within which all control information is permanently stored.





10 (b) What do you mean by Complex Instruction Set Computer (CISC)? Discuss relative advantages and disadvantages of such instruction set design

CISC-1 M + Advantages and disadvantages of CISC-3M

Complex Instruction Set Computer (CISC) is a computer architecture in which single instructions can execute several low-level operations or are capable of multi-step operations or addressing modes within single instructions.

Example of CISC architectures are System/360, PDP-11, Motorola 6800.

CISC works by combining simple instructions into a single complex one, thereby optimizing the instructions per program and reducing the number of instructions that a particular program has.

Advantages of CISC:

- Microprogramming requires assembly language that is easier to implement.
- Instructions that manipulate operands in memory
- The CISC architecture reduces the amount of work that the compilers have to do because the instructions are already
 high-level
- Fewer instructions needed to implement a task

Disadvantages of CISC:

- The number of general-purpose registers that can be fitted into the processor is less because decoding instructions require more transistors
- As CPU does more work in a single instruction, the clock speed tends to be slightly slower than a RISC-based CPU
- The code requires several clock cycles to execute a single instruction despite having a minimal code size. This can
 decrease system efficiency.
- To simplify software hardware needs to be complex

11 (a) Explain the basic computer instruction formats

Computer instruction format-1 M + Types-Zero,One-address,Two-Address,Three-Address and RISC with examples-4M Computer instruction format is defined as standard machine instruction format that can be directly decoded and executed by the CPU. It defines the layout of the instruction.

Instruction includes a set of operation codes and operands that manage with the operation codes. Instruction format contains fields including opcode, operands, and addressing mode.

Based on number of address, instructions are classified as zero-address, one-address, two-address, three-address and RISC. Example : X = (A+B) * (C+D)

Zero Address In	· · ·		One Ad	dress In	struction	Ì	Two Addres	s Instruction
PUSH A	TOS←A		LOAD	Α	AC←M	[A]	MOV R1,A	R1←M[A]
PUSH B	TOS←B		ADD	8	AC← A	C+ M[B]	ADD R1,B	R1←R1+M[B]
ADD	TOS← (A + B)		STORE	Т	M[T]←/	AC	MOV R2,C	R1←M[C]
PUSH C ==	TOS←C		LOAD	С	AC←M	[C]	ADD R2,D	R2←R2+M[D]
PUSH D	TOS←D		ADD	D	AC← A	C+ M[D]	MUL R1,R2	R1←R1*R2
ADD	TOS←(C + D)		MUL	Т	AC← A	C * M[T]	MOV X,R1	M[X]←R1
MUL	$TOS \leftarrow (C + D) X (A + B)$)	STORE		→ [X]M	AC		
POP	M[X] ← TOS							
Three Address			RISC Ir	structio	п			
ADD R1, A, B	R1 ← M[A]+M[B]		LOAD	R1, A		R1←M[A]		
ADD R2, C, D	R2 ←M[C]+M[D]		LOAD	R2, B		R2←M[B]		
MUL X,R1,R2	M[X] ← R1 * R2		LOAD	R3, C		R3←M[C]		
	6 8		LOAD	R4, D		R4←M[D]		
			ADD	R1,R1,	R2	R1←R1+R2		
			ADD	R3,R3,	R2	R3←R3+R2		
			MUL	R1,R1,	R3	R1←R1*R3		
			STORE	X,R1		M[X]		

11 (b) What do you mean by addressing modes? Explain the following addressing modes

(i) Index addressing mode (ii) Immediate Addressing mode (iii) Relative Addressing mode (iv) Direct addressing mode Addressing Mode-1 M + Index, immediate, relative, direct addressing modes-1.5 M for each addressing mode with examples Addressing mode refers to the way in which the operand of an instruction is specified. The addressing mode specifies a rule for interpreting or modifying the address field of the instruction before the operand is actually executed.

Immediate addressing mode(Symbol #): In this mode data is present in the address field of the instruction itself. Example : MOV AL,35H

Index addressing mode: In this mode data index register is added to the address part of the instruction to obtain the effective address.

Example:LD ADR(X) AC←M[XR + ADR]

Relative addressing mode: In this mode program counter is added to the address field of the instruction to obtain the effective address. It is used to implement intra segment transfer of control.

ExampleLD \$ADR AC←M[PC + ADR]

Direct addressing mode:

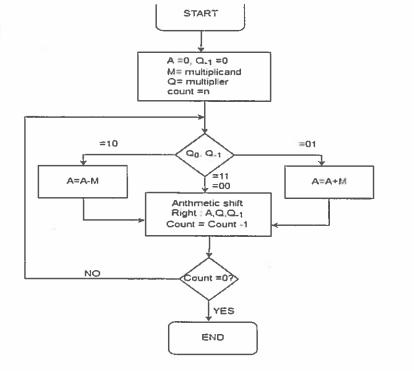
The effective address of the data is part of the instruction itself. Here only one memory reference is used to access the data as the address specifies which register or memory word contains the operand.

Instruction Memory



Example: ADD AL,[0123]. This instruction adds the contents of location 123 to AL and stores the result in AL

12 (a) Discuss Booth multiplication algorithm Booth algorithm flowchart-4M+Example-2M



Here arithmetic shift right (ashr) is performed on A and multiplier Q including the appended Q_{-1} . Example : Perform multiplication of 7 and 3

	chomin mana				
Α	Q	Q_1	M	SC	Remarks
0000	0011	0	0111	4	Initital values
1001	0011	0	0111	3	Subtract M { A←A – M }
1100	1001	1	0111		Perform Arithmetic shift right(ashr)
1110	0100	1	0111	2	Perform Arithmetic shift right(ashr)
0101	0100	1	0111	1	Addition {A←A + M }
0010	1010	0	0111		Perform Arithmetic shift right(ashr)
0001	0101	0	0111	0	Final product is available in A and Q
					14

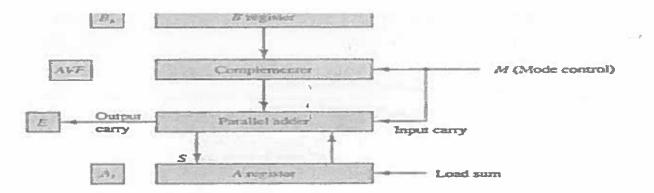
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12 (b) Explain the functional units and their data flow in a hardware implementation th	at performs addition and subtraction of
signed-magnitude numbers	

Table for addition	& subtraction of signed	d magnitude fixed	point-2MHardware for	signed mag	nitude addition	and subtraction-4M
	J					

		Subtract Magnitudes			
Operation	Add Magnitudes	When $A > B$	When $A < B$	When $A = B$	
(+A) + (+B)	+(A + B)				
(+A) + (-B)		+(A - B)	-(B-A)	+(A - B)	
(-A) + (+B)		-(A - B)	+(B - A)	+(A - B)	
(-A) + (-B)	-(A + B)	• •		· /	
(+A) - (+B)		+(A - B)	-(B - A)	+(A - B)	
(+A) - (-B)	+(A + B)	. ,			
(-A) - (+B)	-(A + B)				
(-A) - (-B)		-(A - B)	+(B-A)	+(A - B)	

Table-Addition and subtraction of signed-magnitude-numbers



13 (a) Draw a flowchart which explains multiplication of two signed magnitude fixed point numbers Flowchart for multiplication of 2 signed magnitude fixed point numbers—6M

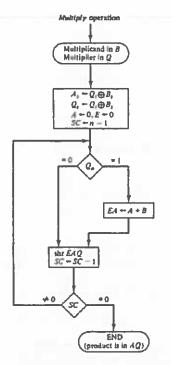
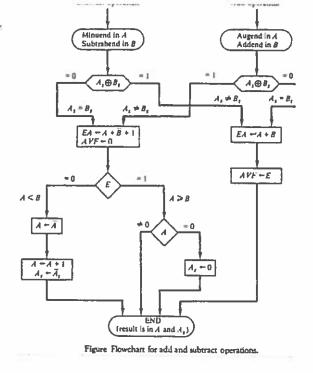


Figure Flowchart for multiply operation.

13 (b) Draw a flowchart for adding and subtracting two fixed point binary numbers where negative numbers in signed 1's complement representation

8

Flowchart for addition and subtraction of two fixed point binary numbers in signed 1's complement-6M

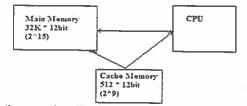


14 (a) What is the significance of cache memory and write about direct and associative mapping techniques Significance of cache memory—1M + Direct Mapping —3M + Associative mapping—3M

Cache memory consumes less access time compared to main memory. Cache is faster than main memory and hence CPU first refers to any data if available in cache, if available it is called a hit and improves the performance of the system. The transformation of data from main memory to cache memory is referred to as a mapping process. There are three types of mapping:

- i) Associative mapping
- ii) Direct mapping
- iii) Set-associative mapping

To help understand the mapping procedure, we have the following example:



Associative mapping: The fastest and most flexible cache organization uses an associative memory. The associative memory stores both the address and data of the memory word. This permits any location in cache to store any word from main memory The address value of 15 bits is shown as a five-digit octal number and its corresponding 12- bit word is shown as a four-digit octal number.

CPU address (15 bits) Argument register	
address	data
01000	3450
01777	6710
22345	1234
L	

9

If the address is found, the corresponding 12- bits data is read and sent to the CPU. If not, the main memory is accessed for the word. If the cache is full, an address-data pair must be displaced to make room for a pair that is needed and not presently in the cache.

Direct Mapping

Associative memory is expensive compared to RAM. In general case, there are 2^k words in cache memory and 2ⁿ words in main memory (in our case, k=9, n=15). The n bit memory address is divided into twofields: k-bits for the index and n-k bits for the tag field.

Each word in cache consists of the data word and its associated tag. When a new word is first brought into the cache, the tags are stored alongside the data bits. When the CPU generates a memory request, the index field is used for the address to access the cache.

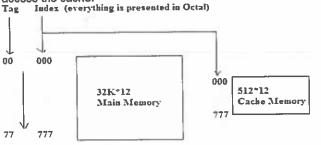


Fig. Addressing relationships between main and cache memories

Memory Address	Memory Data	Index Address	Tag	g Data	
00000	1220	000	00	1220	
		111	01	2222	
00777 01000	2340 3450				
01111	2222	777	02	6710	
01777 02000	4560 5670	(b) Cache memory			
02777	6710				
	J	Ļ			
(a) Mai	n memory				

Fig. Direct mapping cache organization

The tag field of the CPU address is compared with the tag in the word read from the cache. If the two tags match, then it's a hit and the desired data word is in cache. Otherwise it's a miss and required word is read from main memory. It is then stored in the cache together with the new tag, replacing the previous value.

10

Consider the numerical example shown above. The word at address zero is presently stored in the cache(index=000, tag=00, data=1220). Let us assume that CPU wants to access the word at address 01000. The index address is 000, so it is used to access the cache. The two tags are then compared. The cache tag is 00 but the address tag is 01, which does not produce a match. Hence, the main memory is accessed and the data word 3450 is transferred to the CPU. The cache word at index address 000 is then replaced with a tag of 01 and data of 3450.

14 (b) Distinguish between synchronous and asynchronous data transfer

Synchronous data transfer—1 M + Asynchronous data transfer—1 M+ Differences between them—3 M

In synchronous data transfer the sending and receiving units are enabled with the same clock signal

Synchronous data transfer	Asynchronous data transfer
The master performs a sequence of instructions for data	None of the actions are synchronized with a common clock,
transfer in a predefined order	hence no predefined order
As there is common clock there is no need of control signal	The asynchronous data transfer between the two units require
between source and destination while transferring data	control signals be transmitted between the communicating
	units to indicate when to send data
The master does not expect any acknowledgement when data	In handshaking method there is possibility of knowing whether
is sent by the master to the slave	sent data is received or not at the receiver end via
	acknowledgement

15 (a) With the help of a block diagram, explain the concept of DMA controller DMA controller—1M + Block diagram of DMA controller with explanation—5M

DMA Controller:

1

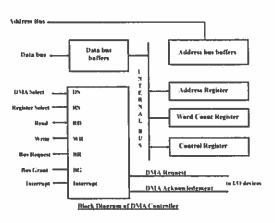
The DMA controller needs the usual circuits of an interface to communicate with the CPU and I/O device. The DMA controller has three registers :Address Register, Word Count Register and Control Register

Address Register :- Address Register contains an address to specify the desired location in memory.

Word Count Register(WC) :- WC holds the number of words to be transferred. The register is incremented/decremented by one after each word transfer and internally tested for zero.

Control Register - Control Register specifies the mode of transfer

The unit communicates with the CPU via the data bus and control lines. The registers in the DMA are selected by the CPU through the address bus by enabling the DS (DMA select) and RS (Register select) inputs. The RD (read) and WR (write) inputs are bidirectional.



When the BG (Bus Grant) input is 0, the CPU can communicate with the DMA registers through the data bus to read from or write to the DMA registers. When BG =1, the DMA can communicate directly with the memory by specifying an address in the address bus and activating the RD or WR control.

15 (b) What is virtual memory? With a neat block diagram explain the virtual memory address translation

Virtual memory—1M + Virtual memory address translation with neat diagram—5M

A computer can address more memory than the amount physically installed on the system. This extra memory is actually called virtual memory and it is a section of a hard disk that's set up to emulate the computer's RAM.

A virtual memory system provides a mechanism for translating program-generated addresses into correct main memory locations. This is done dynamically, while programs are being executed in the CPU. The translation or mapping is handled automatically by the hardware by means of a mapping table.

An address used by a programmer will be called a virtual address, and the set of such addresses the address space. An address in main memory is called a location or physical address. The set of such locations is called the memory space.

For example, consider a computer with a main-memory capacity of $32K (2^5 X 2^{10})$ words .15 bits are needed to specify a physical address in memory . Suppose that the computer has available auxiliary memory for storing $2^{20} = 1024K$ words. Thus auxiliary memory has a capacity for storing information equivalent to the capacity of 32 main memories. For this example address space N= 1024K and memory space M = 32K.

In a multiprogram computer system, programs and data are transferred to and from auxiliary memory and main memory based on demands imposed by the CPU. Suppose that program 1 is currently being executed in the CPU.

In a virtual memory system the address field of the instruction code has a sufficient number of bits to specify all virtual addresses. In our example, the address field of an instruction code will consist of 20 bits but physical memory addresses must be specified with only 15 bits. Thus CPU will reference instructions and data with a 20-bit address, but the information at this address must be taken from physical memory because access to auxiliary storage for individual words will be prohibitively long.

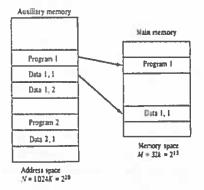


Figure Relation between address and memory space in a virtual memory system.

A table is needed to map a virtual address of 20 bits to a physical address of 15 bits.

The mapping table may be stored in two ways:

1) Stored in a separate memory (additional memory) -Here an additional memory unit is required as well as one extra memory access time

2) Main memory --Here the table takes space from main memory and two accesses to memory are required with the program running at half speed.

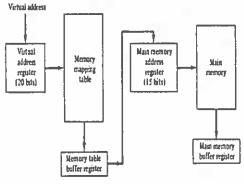


Figure Memory table for mapping a virtual address.

The physical memory is broken down into groups of equal size called blocks, which may range from 64 to 4096 words each. The term "page frame" is sometimes used to denote a block. The term page refers to groups of address space of the same

size. Consider a computer with an address space of 8K and a memory space of 4K. If we split each into groups of 1K words we obtain eight pages and four blocks as shown in Fig.

At any given time, up to four pages of address space may reside in main memory in any one of the four blocks.

Block 0
Block J
Block 2
Block 3
Memory spac $M = 4K = 2^{12}$

1 1

Address space and memory space split into groups of 1K words.

Address mapping using pages and tables

The table implementation of the address mapping is simplified if the information in the address space. And the memory space is each divided into groups of fixed size.

The physical memory is broken down into groups of equal size called blocks, which may range from 64 to 4096 words each. The term page refers to groups of address space of the same size.

Also, consider a computer with an address space of 8K and a memory space of 4K.

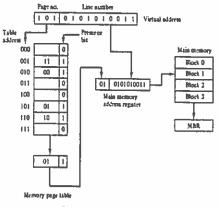


Figure Memory table in a paged system.

Consider a computer with an address space of 8K and a memory space of 4K.

If we split each into groups of 1K words we obtain eight pages and four blocks as shown in the figure.

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At any given time, up to four pages of address space may reside in main memory in any one of the four blocks.

Nadimpalli Satyanarayana Raju Institute of Technology (Autonomous). IQAC: Quality Management System (QMS)

Semester End Regular/Supplementary Examination, Dec./Jan., 2022 - 2023

NSRIT

Course	Power Generatio	n and Transmissic					
Course Code	50EE302	Test Duration	3 Hrs.	Max. Marks	02	Semester	111
Degree	B. Tech. (U. G.)	Program	333			Academic Year	2022 - 2023

Part A (Short Answer Questions 5 x 2 = 10 Marks)

17	S0EE302'2	(augure fund fur adur la analysis and fur	
P 1		Mention any two methods of improving string efficiency.	9
11	20EE305.4	Recall surge impedance loading (SIL).	Þ
เา	20EE306.3	Define GMR and GMD.	3
17	20EE306.2	Define the diversity factor.	7
11	20EE306.1	List any two disadvantages of nuclear power plant.	1
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Part B (Long Answer Questions 5 x 12 = 60 Marks)

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F3	20EE306.5	M8	Derive an expression for sag of a line supported between two	(-) =
		1410	supports of the same tower height.	(e) ç
27	S0EE302'2	Mħ	A 132 kV transmission line has the following data: Wt. of conductor = 680 kg/km; Length of span = 260 m Ultimate strength = 3100 kg; Safety factor = 2. Calculate the height above ground at which the conductor should be supported. Ground clearance required is 10 m.	(q) g

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N S RAJU INSTITUTE OF TECHNOLOGY

(AUTONOMOUS)

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ANSWER KEY AND SCHEME OF EVALUATION

Part-A:	n Herstein († 1975) 17. júlio – Maria
91: Two adisadvantages -	
92: Definition - IM Formula - IM.	
Q3: GMR - IM	gala da Maria de Como de Anglia. A companya da Maria
04: SIL Definition - 2M	
05: Two methods - IM	for each.
Part-B	ine i se Roman se
Q6: Figure : 4M Components List : 4M	e ^{lle} e (j. x.). el
Explanation: 4M.	
Q7: Figure : 4 M Component List: 4 M Explanation : 4 M.	

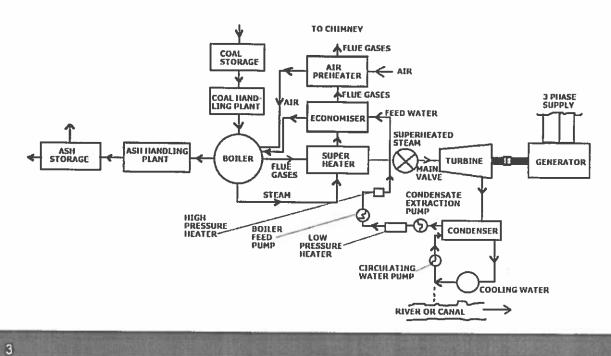
98: U) 3M ; (ii) 3M ; (iii) 3M ; (iv). Curve: 3M.	2M
	6 Pypes 2M for each types.	
Q10:	Basic explanation : 4M.	
	Unsymmetrical: 4M. Derivation	
	Paransposition: 4M. Derivation.	
	Enplanation of Law: 4M.	
	Descivation : 8M.	
912:	Short 7/m Line: 2M	
	Medium 7/m Line : 6M.	
	Long T/m Line : 2M. DC T/m line : 2M Derivation : 6M.	
5	DC T/m line : 2M	
6:	Figures : 3M.	
	Derivation: 3M.	
Q14: Ea	ach richter	
Q 15 a:	Figure: 4M Derivation: 4M.	
b:	O I F G I I F O F HILL	

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1.	 List any two disadvantages of nuclear power plant. Any two of the below a) Strong pressure vessel is required due to the use of high-pressure water system b) Formation of low temperature steam c) Use of expensive cladding material for prevention of corrosion d) 4.High losses from heat exchanger e) 5.High power consumption by auxiliaries f) 6.Requires more elaborate safety devices 	(2m)
2.	Define the diversity factor. It shows the diversity of load connected to a power station Diversity factor = Sum of individual maximum demands / Maximum demand on power system	(2m)
3.	Define GMR and GMD. GMR stands for Geometrical Mean Radius. It is also called the self GMD (Geometrical Mean Dista GMR = 0.7788R Where R is the radius of the conductor	(1m)
	GMD stands for Geometrical Mean Distance GMD represents the geometrical mean distance conductor to the other. GMD for a different arrangement of conductors has different values.	e from one (1m)
4.	Recall surge impedance loading (SIL). SIL is defined as the maximum load (at unity power factor) that can be delivered by the transmission the loads terminate with a value equal to surge impedance (Zs) of the line.	on line when (2m)
5.	Mention any two methods of improving string efficiency. Any of the following By Using Insulators with Larger Discs or by Providing Each Insulator Unit with a Metal Cap By Using Longer Cross-Arms By Capacitance Grading By Static Shielding	(2m)
6	Evolution the construction and working principle of the male successful at	

6. Explain the construction and working principle of thermal power plant.

5 V.



A Thermal Power plant converts the heat energy of coal into electrical energy. Coal is burnt in boiler which converts water into steam.

The expansion of steam in turbine produces mechanical power which drives the Alternator coupled to the turbine. Thermal power plants contribute maximum to the generation of power for any country.

In thermal generating stations coal, oil, natural gas etc. are employed as primary courses of energy. Main Components

- Coal handling plant
- Pulverizing plant
- Boiler
- Turbine
- Condenser
- Cooling towers and ponds
- Feed water heater
- Economizer
- Air preheater

Coal Handling Plant

- Coal is transported to power station by rail or road and stored in coal storage plant and then pulverized.
- The function of coal handling plant is automatic feeding of coal to the boiler furnace.
- A thermal power plant burns enormous amounts of coal.
- A 200MW plant may require around 2000 tons of coal daily.

Pulverizing Plant

- In modern thermal power plant, coal is pulverized i.e. ground to dust like size and carried to the furnace in
 a stream of hot air. Pulverizing is a means of exposing a large surface area tot the action of oxygen and
 consequently helping combustion.
- Pulverizing process consists 3 stages classified as:
 - 1.Feeding
 - 2.Drying.
 - 3.Grinding

Boiler

The function of boiler is to generate steam at desired pressure and temperature by transferring heat
produced by burning of fuel in a furnace to change water into steam.

Turbine

- In thermal power plants generally 3 turbines are used to increase the efficiency.
- High pressure turbine
- Intermediate pressure turbine
- Low pressure turbine

Condenser

 The surface condenser is a shell and tube heat exchanger where cooling water flows through tubes and exhaust steam fed into the shell surrounds the tubes, as a result, steam condense outside the tubes.

Cooling Towers and Ponds

- A condenser needs huge quantity of water to condense the steam.
- Most plants use cooled cooling system where warm water coming from condenser is cooled and reused.
- Cooling tower is a steel or concrete hyperbolic structure with the height of 150m.

Feed water heater

- Feed water heating improves overall plant efficiency
- Thermal stresses due to cold water entering the boiler drum are avoided.
- Quality of steam produced by the boiler is increased.

Economizer

4

Flue gases coming out of the boiler carry lot of heat. An economizer extracts a part of this heat from flue
gases and uses it for heating fee water.

Saving coal consumption and higher boiler efficiency.

Air Preheater

- The function of air preheaters is to preheat the air before entering to the furnace by utilizing some of the • energy left in the flue gases before exhausting them to the atmosphere.
- After flue gases leave economizer, some further heat can be extracted from them and used to heat • incoming heat. Cooling of flue gases by 20-degree centigrade increases the plant efficiency by 1%.

Ash Handling plant

- The ash from the boiler is collected in two forms •
- Bottom ash (Slurry): It's a waste which is dumped into ash pond
- Fly ash: Fly ash is separated from flue gases in ESP. .

Water Handling Plant

- Water in a power plant is used for .
- Production of steam for rotating turbine
- Cooling purpose For cooling of various equipment.
- Water is recycled and used for various purpose: Raw water - For cooling purposes - Steam - Condenser - Raw water

About 4 cubic meter water is lost/day/MW.

Electrostatic precipitator (ESP)

- An ESP electrically charges the ash particles and imparts a strong electric field in the flue gas to collect • and remove them. ESP is comprised of a series of parallel, vertical metallic plates (collecting electrodes) forming lanes through which the flu gases pass.
- 7. Explain the layout, classification, and operation of hydro power plant.

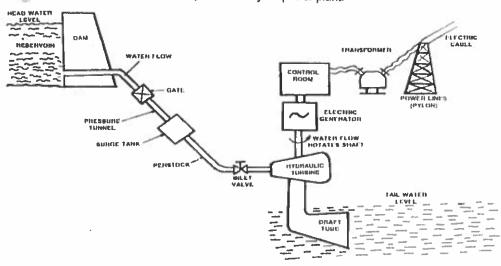


Fig.Layout of Hydro electric Power plant

(4m)

In hydroelectric power station potential and kinetic energy of stored water is converted into electrical energy. Working Principle

Hydroelectric power is the power obtained from the energy falling water whereas hydroelectric power • plant is the power utilizing the potential energy of water at a high level for the generation of the electrical energy.

Main Components

- ÷. Reservoir
- Dam
- Trace rack
- Turbine
- Forebay

- Surge tank
- Penstock
- Spillway
- Turbine
- Powerhouse

Dam

- Develops a reservoir to store water
- Builds up head for power generation

Spillway

To safeguard the dam when water level in the reservoir rises

Intake

Contains trash racks to filter out debris which may damage the turbine

Forebay

 The forebay is used as "regulating reservoir" storing water temporarily during light load and providing the same for initial increases on account of increasing load during which water in canal is being accelerated. A forebay is an enlarged body of water just above the intake which is used to store water temporarily to meet the hourly load fluctuations. A forebay is not required if plant is located just at the base of the dam but, if the plants are situated away from the storage reservoir, a forebay is must.

Surge tank

- Surge tank is a small reservoir in which the water level rises or falls to reduce the pressure swings so that they are not transmitted to the penstock.
- When the load demand is reduced on the power station then, it causes rise in water level in the surge tank which produces a retarding head and reduces the velocity of water in the penstock and hence avoiding the undesirable phenomenon called "water hammer"
- When the load on the plant is increased, the governor causes the turbine to open the gates in order to allow more water to flow through the penstock to supply the increased load and there is a tendency to cause a vacuum or a negative pressure in the penstock.

Penstock

 Penstock is a closed conduit which connects the forebay or surge tank to the scroll case of the turbine. In case of high head plants, a single penstock is provided.

Valves and Gates

 Gates are used in low head plants at the entrance to the turbine casing to shut-off the flow and provide for unwatering the turbine for inspection and repairs. Valves are used at the entrance to the turbine casing if a long or medium length penstocks is used in the hydro power plant.

Trash Racks

Trash racks are used to prevent the ingress of floating and other material to the turbine. These are built
up from long, flat bars set vertically or nearly so and spaced in accordance with the minimum width of
water passage through the turbine.

Tail race

After the useful work is done by water, it is discharged to the tail race.

Draft tube

6

 It is an airtight pipe of suitable diameter attached to the runner outlet and conducting water down from the wheel and discharging it under the surface of the water in the tail race. With the help of draft tube operating head on the turbine is increased resulting in increase in output and efficiency

Prime Movers or water turbines

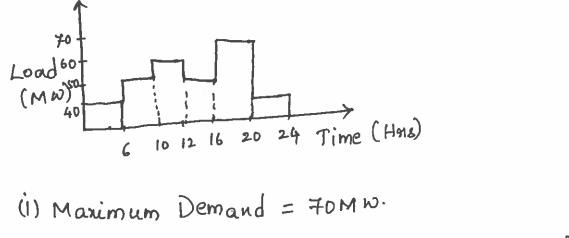
 In hydroelectric plants water turbines serve the purpose of prime mover which converts the kinetic energy of water into mechanical energy which is further utilized to drive the alternators generating electrical energy.

8. A generating station has the following daily load cycle:

(4m)

Time (Hrs)	0-6	6-10	10-12	12-16	16 - 20	20-24
Load (MW)	40	50	60	50	70	40

Draw the load curve and find (i) maximum demand (ii) units generated per day (iii) average load and (iv) load factor.



(ii) Unite Generated per day = $(6 \times 40) + (4 \times 50) + (2 \times 60)$ + $(4 \times 50) + (4 \times 70) + (4 \times 40)$ = 1200 MWh. (iii) Average Load = $\frac{\text{Unite generated}}{\text{Hns in a day}}$ = $\frac{1200 \text{ MWh}}{24} = \frac{1200 \times 10^6}{24}$ = $50 \times 10^6 = 50 \text{ MW}.$ (iv) Load Factor = $\frac{\text{Average Demand} \times 100}{\text{Maximum Demand}} = \frac{50 \text{ MW}}{70 \text{ MW}}$

7

= 71.4.

9. Explain the types of Tariff methods.

Simple Tariff: When there is a fixed rate per unit of energy consumed, it is the Simple Tariff. In this type of tariff, the price charged per unit is constant. It does not vary with increase or decrease with the number of Units consumed.

Flat Rate Tariff: When different types of Consumers are charged at different uniform per unit rates, it is the Flat Rate Tariff. The rate for each type of consumer is arrived at by taking its load factor, diversity factor into consideration. The Bill will be Total units Consumed x Rate/Unit.

Block Rate Tariff: When a given block of energy is charged at a specific rate and the succeeding blocks of energy are charged progressively at reduced rates. Then the Tariff is called the Block Rate Tariff.

If the number of units generated increases, then the cost of generation per-unit-automatically decreases. For the first 30 units may be charged at the rate of 60paisa per unit, the next 29 units at the rate of 55paise per unit and the remaining additional units may be charged at the rate of 30 paise per unit. This type of tariff is being majorly used for residential and small commercial consumers.

Two Part Tariff: When the rate of electric energy is charged based on maximum demand of the consumer and the units consumed, it is called the Two-part Tariff.

In two-part tariff, the total charge to be made from the consumer is spilt into two components, fixed charges and running charges. The fixed charges depend upon the maximum demand of the consumer, while the running charges depend upon the no. of units consumed by the consumer.

Total Charges = Rs. (B kW+ C kWh)

B = Charges per kW of maximum demand

C = Charges per kWh of energy consumed

This type of Tariff is mostly applicable to Industrial Consumers.

Maximum Demand Tariff: It is similar to Two Part Tariff with the only difference that the maximum demand is actually measured by installing maximum demand meter in the premises of the consumer.

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This type of tariff is mostly applied to big consumer, as a separate maximum demand meter is required.

Three-part tariff: When the total charge to be made from the consumer split into three parts, fixed charges, Semi fixed charges and running Charges.

Total Charge = Rs. (A + B kW + C kWh)

A = Fixed Charges during each billing period

B = Charge per kW of Maximum demand

C = Charge per kWh of energy consumed

Power Factor Tariff:

The Tariff in which Power factor of the Consumer's load is taken into consideration is known as Power Factor Tariff.

The following are the important types of Power Factor tariff

- (i) kVA maximum demand tariff: It is a modified form of two-part tariff. The fixed charges are made on the basis of maximum demand in kVA and not in kW. As kVA is inversely proportional to power factor, a consumer having low power factor has to contribute more towards the fixed charge.
- Sliding Scale Tariff: This is also known as average PF tariff. In this case an average power factor say 0.8 lagging is taken as reference. If the pf of the consumer falls below this factor, suitable additional charges are made. If the PF is above the reference, a discount is allowed to the consumer.
 kW and kVAR Tariff: In this type both active power (kW) and reactive power (kVAR) supplied are
 - kW and kVAR Tariff: In this type both active power (kW) and reactive power (kVAR) supplied are charged separately. A consumer having low PF will draw more reactive Power and hence shall have to pay more charges.
- 10. Derive the equation for inductance of a three-phase overhead line.

Inductance in Three Phase Transmission Line:

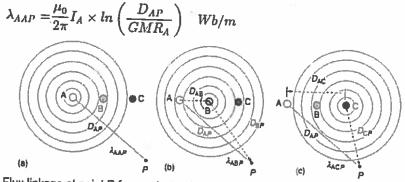
In the three-phase transmission line, three conductors are parallel to each other. The direction of the current is same through each of the conductors.

Let us consider conductor A produces magnetic flux ϕ_{A} , Conductor B produces magnetic flux ϕ_{B} , And conductor C produces magnetic flux ϕ_{C} . When they carry the current of the same magnitude "I", they are in flux linkage with each other.

Now, let us consider a point P near three conductors. So, flux linkage at point P due to current through conductor A is,

$$\lambda_{AP} = \lambda_{AAP} + \lambda_{ABP} + \lambda_{ACP}$$

Flux linkage at point P for conductor A due to current through conductor A =



Flux linkage at point P for conductor A due to current through conductor B =

$$\lambda_{ABP} = \frac{\mu_0}{2\pi} I_B \times \ln\left(\frac{D_{BP}}{D_{AB}}\right) \quad Wb/m$$

Flux linkage at point P for conductor A due to current through conductor C =

$$\lambda_{ACP} = \frac{\mu_0}{2\pi} I_C \times ln \left(\frac{D_{CP}}{D_{AC}} \right) \quad Wb/m$$

Therefore, flux linkage at point P for conductor A,

$$\Rightarrow \lambda_{AP} = \frac{\mu_0}{2\pi} \left[I_A \times ln \left(\frac{1}{GMR_A} \right) + I_B \times ln \left(\frac{1}{D_{AB}} \right) + I_C \times ln \left(\frac{1}{D_{AC}} \right) \right]$$

$$+ \frac{\mu_0}{2\pi} \left[I_A \times ln(D_{AP}) + I_B \times ln(D_{BP}) + I_C \times ln(D_{CP}) \right] \quad Wb/m$$
(2m)

As, $D_{AP} = D_{BP} = D_{CP}$ and $I_A + I_B + I_C = 0$ in balanced system, then we can write that $I_A = -I_B - I_C$

$$\therefore \frac{\mu_0}{2\pi} [I_A \times \ln(D_{AP}) + I_B \times \ln(D_{BP}) + I_C \times \ln(D_{CP})]$$

$$= \frac{\mu_0}{2\pi} [-I_B \times \ln(D_{AP}) - I_C \times \ln(D_{AP}) + I_B \times \ln(D_{BP}) + I_C \times \ln(D_{CP})] = 0$$

$$\Rightarrow \lambda_{AP} = \frac{\mu_0}{2\pi} \left[I_A \times \ln\left(\frac{1}{GMR_A}\right) + I_B \times \ln\left(\frac{1}{D_{AB}}\right) + I_C \times \ln\left(\frac{1}{D_{AC}}\right) \right] + 0 = \lambda_A (say)$$
So, $\lambda_A = \frac{\mu_0}{2\pi} \left[I_A \times \ln\left(\frac{1}{GMR_A}\right) + I_B \times \ln\left(\frac{1}{D_{AB}}\right) + I_C \times \ln\left(\frac{1}{D_{AC}}\right) \right]$
Similarly, $\lambda_B = \frac{\mu_0}{2\pi} \left[I_A \times \ln\left(\frac{1}{D_{BA}}\right) + I_B \times \ln\left(\frac{1}{GMR_B}\right) + I_C \times \ln\left(\frac{1}{D_{BC}}\right) \right]$

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and,
$$\lambda_{C} = \frac{\mu_{0}}{2\pi} \left[I_{A} \times ln\left(\frac{1}{D_{CA}}\right) + I_{B} \times ln\left(\frac{1}{D_{CB}}\right) + I_{C} \times ln\left(\frac{1}{GMR_{C}}\right) \right]$$

For a balanced system,

$$\begin{split} D_{AB} &= D_{BC} = D_{CA} = D\\ I_A + I_B + I_C &= 0 \end{split}$$

In balanced system, then we can write that, $\, I_A = -I_B - I_C \,$

$$\begin{split} \lambda_A = & \frac{\mu_0}{2\pi} \left[I_A \times \ln\left(\frac{1}{GMR_A}\right) + I_B \times \ln\left(\frac{1}{D}\right) + I_C \times \ln\left(\frac{1}{D}\right) \right] \\ = & \frac{\mu_0}{2\pi} \left[I_A \times \ln\left(\frac{1}{GMR_A}\right) + (I_B + I_C) \times \ln\left(\frac{1}{D}\right) \right] \\ = & \frac{\mu_0}{2\pi} \left[I_A \times \ln\left(\frac{1}{GMR_A}\right) + (-I_A) \times \ln\left(\frac{1}{D}\right) \right] \\ = & \frac{\mu_0}{2\pi} I_A \times \ln\left(\frac{D}{GMR_A}\right) \quad Wb/m \end{split}$$

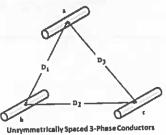
$$\lambda_B = \frac{\mu_0}{2\pi} I_B \times \ln\left(\frac{D}{GMR_B}\right) \quad Wb/m$$
$$\lambda_C = \frac{\mu_0}{2\pi} I_C \times \ln\left(\frac{D}{GMR_C}\right) \quad Wb/m$$

So, inductance per metre per phase, $L_{\text{phase}} = \frac{\mu_0}{2\pi} \times \ln\left(\frac{D}{GMR_{\text{phase}}}\right) H/m$

Consider a 3-phase overhead transmission line with phase conductors a, b, and c of radius r each spaced unsymmetrically such that the distance between the three conductors is D1, D2, and D3 as shown in the figure below. Let I_a , I_b , and I_c be the currents flowing in the conductors a, b, and c respectively.

Assume that the whole system is balanced which leads to the flow of equal currents in the conductors i.e., Ia = $I_b = I_c = I(say)$. The currents I_a , I_b , and I_c are displaced at 120° apart from each other. If I_a is taken as the reference phasor then the currents are given by,

 $I_a = I(1-j0)$ $I_b = I(-0.5 - j0.866)$ and $I_c = I(-0.5 + j0.866)$



(2m)

(2m)

We know that the flux linkage of any conductor, in a group of conductors, is due to its own current and currents in the other conductors. Therefore,

flux linkage of the conductor 'a' is due to its own current and currents in the conductor's 'b' and 'c', and it is given by,

$$\lambda_a = 2 \times 10^{-7} \left[I_a \ln \frac{1}{r'} + I_b \ln \frac{1}{D_1} + I_c \ln \frac{1}{D_3} \right]$$

Similarly, flux linkages of conductors 'b' and 'c' is given by,

$$\begin{split} \lambda_b &= 2 \times 10^{-7} \left[I_b \ln \frac{1}{r'} + I_a \ln \frac{1}{D_1} + I_c \ln \frac{1}{D_2} \right] \\ \lambda_c &= 2 \times 10^{-7} \left[I_c \ln \frac{1}{r'} + I_a \ln \frac{1}{D_3} + I_b \ln \frac{1}{D_2} \right] \end{split}$$

Where, r' = Geometric mean radius (GMR) of conductor = 0.07788 × r. Substituting the values of Ia, Ib, and Ic in the above equation we get,



$$\begin{split} \lambda_{a} &= 2 \times 10^{-7} \left[l \, \ln \frac{1}{r'} + l(-0.5 - j0.866) \ln \frac{1}{D_{1}} + l(-0.5 + j0.866) \ln \frac{1}{D_{3}} \right] \\ \lambda_{a} &= 2 \times 10^{-7} \, l \left[\ln \frac{1}{r'} + \ln \sqrt{D_{1}} + \ln \sqrt{D_{3}} + j0.866 \ln D_{1} + j0.866 \ln \frac{1}{D_{3}} \right] \\ \lambda_{a} &= 2 \times 10^{-7} \, l \left[\ln \frac{1}{r'} + \ln \sqrt{D_{1}D_{3}} + j \frac{\sqrt{3}}{2} \ln D_{1} + j \frac{\sqrt{3}}{2} \ln \frac{1}{D_{3}} \right] \\ \lambda_{a} &= 2 \times 10^{-7} \, l \left[\ln \frac{1}{r'} + \ln \sqrt{D_{1}D_{3}} + j \frac{\sqrt{3}}{2} \ln D_{1} + j \frac{\sqrt{3}}{2} \ln \frac{1}{D_{3}} \right] \end{split}$$

We know that the inductance La is given by,

$$L_{a} = \frac{u}{l_{a}}$$

$$L_{a} = \frac{2 \times 10^{-7}}{l} \left[\ln \frac{1}{r'} + \ln \sqrt{D_{1}D_{3}} + j\sqrt{3} \ln \sqrt{\left(\frac{D_{1}}{D_{3}}\right)} \right]$$

$$\therefore L_{a} = 2 \times 10^{-7} \left[\ln \frac{1}{r'} + \ln \sqrt{D_{1}D_{3}} + j\sqrt{3} \ln \sqrt{\left(\frac{D_{1}}{D_{3}}\right)} \right] H/n$$

λ.,

Similarly, the inductance due to the conductor's 'b' and 'c' can be calculated and they are given by,

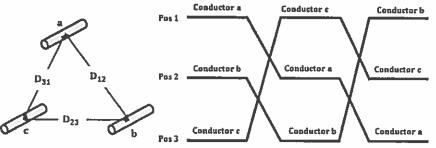
$$\begin{split} & L_b = 2 \times 10^{-7} \left[\ln \frac{1}{r'} + \ln \sqrt{D_1 D_2} + j \sqrt{3} \ln \sqrt{\left(\frac{D_2}{D_1}\right)} \right] H/m \\ & L_c = 2 \times 10^{-7} \left[\ln \frac{1}{r'} + \ln \sqrt{D_2 D_3} + j \sqrt{3} \ln \sqrt{\left(\frac{D_3}{D_2}\right)} \right] H/m \end{split}$$

When conductors are unsymmetrically spaced in a 3-phase line, the flux linkage and inductance per phase are not identical, and knowing the inductance of each phase becomes complicated and it results in an unbalanced circuit.

The equilibrium in the circuit can be retained by shifting the places of the conductor at every period such that the conductors take the initial place of every conductor at the same distances.

Such an arrangement of the conductor's position obtained by shifting their places is called transposition. By transposition, we can obtain almost the same inductance between the conductors. Let us see the expression for inductance when the 3-phase line is transposed.

Inductance of Unsymmetrically Spaced 3-Phase Transmission Line When Transposed: (2m) The simple configuration of a 3-phase conductor is shown in the figure. As the conductors are transposed, so their position in the transposition cycle would be as shown in the figure below.



Unsymmetrically Spaced 3-Phase Conductors with Transposition Cycle

When the conductors are connected in parallel, it results in low inductance keeping the distance between phases as small as possible. For deriving the inductance per phase of 3-phase conductors placed unsymmetrically, we need to determine the flux linkage of each conductor for every position it occupies in a transposition cycle. After that, we can determine the average flux linkages. So, the flux linkage of phase 'a' in position 1 is given as,

$$\lambda_{a1} = 2 \times 10^{-7} \left[I_a \ln \frac{1}{D_s} + I_b \ln \frac{1}{D_{12}} + I_c \ln \frac{1}{D_{31}} \right]$$

When 'a' is in position 2, 'b' in position 3 and 'c' in position 1,

$$\lambda_{a3} = 2 \times 10^{-7} \left[I_a \ln \frac{1}{D_s} + I_b \ln \frac{1}{D_{31}} + I_c \ln \frac{1}{D_{23}} \right]$$

The average value of flux linkage of a single-phase 'a' is,

$$\lambda_a = \frac{\lambda_{a1} + \lambda_{a2} + \lambda_{a3}}{3}$$

$$\lambda_{a} = \frac{2 \times 10^{-7}}{3} \left[3I_{a} \ln \frac{1}{D_{s}} + I_{b} \ln \frac{1}{D_{12}D_{23}D_{31}} + I_{c} \ln \frac{1}{D_{12}D_{23}D_{31}} \right]$$

We know that

$$l_{a} + l_{b} + l_{c} = 0$$

$$l_{b} + l_{c} = -l_{a}$$

$$\lambda_{a} = \frac{2 \times 10^{-7}}{3} \left[3l_{a} \ln \frac{1}{D_{s}} - l_{a} \ln \frac{1}{D_{12}D_{23}D_{31}} \right]$$

$$\lambda_{a} = 2 \times 10^{-7} l_{a} \ln \sqrt{\frac{\sqrt{D_{12}D_{23}D_{31}}}{D_{s}}}$$

Therefore, the average inductance per phase is,

$$L_a = 2 \times 10^{-7} \ln \frac{D_{eq}}{D_s} H/m$$

Where,
$$D_{eq} = \sqrt[2]{D_{12}D_{23}D_{31}}$$

Where,

Deg = Geometric mean of three distances of unsymmetrical line

D_s = GMR (geometric mean radius) of the conductor.

The above expression is the inductance per phase of a 3-phase transmission line with unsymmetrical spacing but lines are transposed. Nowadays the transposition of conductors is made at switching stations to balance inductance.

(2m)

If the conductors are equispaced,

D1 = D2 = D3 = D

$$L = 2 \times 10^{-7} \ln \frac{\sqrt[3]{d^3}}{r'} \Rightarrow L = 2 \times 10^{-7} \ln \frac{d}{r'}$$

11. Derive the equation for capacitance of a two-wire overhead line.

Electric Field Intensity due infinite line charge:

Consider a long wire having q coulomb/m as shown in figure

Using Gauss's law, Field Intensity (E) at a point P, which is r metre from the conductor, can be calculated as

$$\phi \overline{D} \cdot \overline{ds} = Q$$

The Flux density at point P, considering the cylindrical shell of radius r and length I can be calculated using Gauss's Law

$$D \cdot 2\Pi r \cdot l = ql \Rightarrow D = \frac{q}{2\pi r} coulomb/m^2$$

And Electric Field intensity E is given by

$$E = \frac{D}{\varepsilon_0}$$

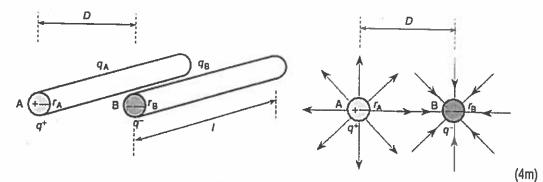
Capacitance and Capacitive Reactance:

Capacitance exists among transmission line conductors due to their potential difference. To evaluate the capacitance between conductors in a surrounding medium with permittivity ε , it is necessary to determine the voltage between the conductors, and the electric field strength of the surrounding. (4m)

Capacitance of a Single-Phase Line with Two Wires:

Consider a two-wire single-phase line with conductors A and B with the same radius r, separated by a distance $D > r_A$ and r_B . The conductors are energized by a voltage source such that conductor A has a charge q+ and conductor B a charge q- as shown in Fig.

The charge on each conductor generates independent electric fields. Charge q+ on conductor A generates a voltage V_{AB-A} between both conductors. Similarly, charge q- on conductor B generates a voltage V_{AB-B} between conductors.



Electric field produced from a two-wire single-phase system.

 V_{AB-A} is calculated by integrating the electric field intensity, due to the charge on conductor A, on conductor B from r_A to D

$$V_{AB-A} = \int_{r_A}^{D} E_A \, \mathrm{d}x = \frac{q}{2\pi\varepsilon_0} \ln\left[\frac{D}{r_A}\right]$$

VAB-B is calculated by integrating the electric field intensity due to the charge on conductor B from D to rB

$$V_{AB-B} = \int_{D}^{r_{B}} E_{B} dx = \frac{-q}{2\pi\varepsilon_{0}} \ln\left[\frac{r_{B}}{D}\right]$$

The total voltage is the sum of the generated voltages V_{AB-A} and V_{AB-B}

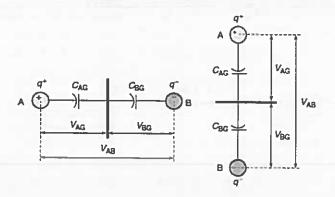
$$V_{AB} = V_{AB-A} + V_{AB-B} = \frac{q}{2\pi\varepsilon_0} \ln\left[\frac{D}{r_A}\right] - \frac{q}{2\pi\varepsilon_0} \ln\left[\frac{r_B}{D}\right] = \frac{q}{2\pi\varepsilon_0} \ln\left[\frac{D^2}{r_A r_B}\right]$$

If the conductors have the same radius, $r_A=r_B=r$, then the voltage between conductors V_{AB} , and the capacitance between conductors C_{AB} , for a 1-m line length are

$$V_{AB} = \frac{q}{\pi \varepsilon_0} \ln \left[\frac{D}{r} \right]$$
(V)
$$C_{AB} = \frac{\pi \varepsilon_0}{\ln \left[\frac{D}{r} \right]}$$
(F/m)

The voltage between each conductor and ground (G) is one-half of the voltage between the two conductors. Therefore, the capacitance from either line to ground is twice the capacitance between lines (4m)

$$V_{AG} = V_{BG} = \frac{V_{AB}}{2} (V)$$
$$C_{AG} = \frac{q}{V_{AG}} = \frac{2\pi\varepsilon_0}{\ln\left[\frac{D}{r}\right]} (F/m)$$



12. Classify the types of transmission lines with model representations.

Classification of Transmission Lines - Short, Medium & Long Transmission Lines:

- A. AC transmission line, and
- B. DC transmission line

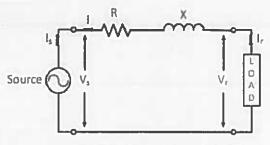
Depending upon the operating voltage and length, the overhead ac transmission lines are classified as,

- 1. Short transmission lines
- 2. Medium transmission lines
- 3. Long transmission lines

Short Transmission Line:

- Overhead transmission line is less than 50km
- Operating voltages of less than 20kV.

In these lines, the effect of capacitance is neglected due to smaller length and low operating voltage. Hence, the resistance and inductance effects of the line are considered while determining the performance of the short transmission line as shown below.



Equivalent circuit model of a short line

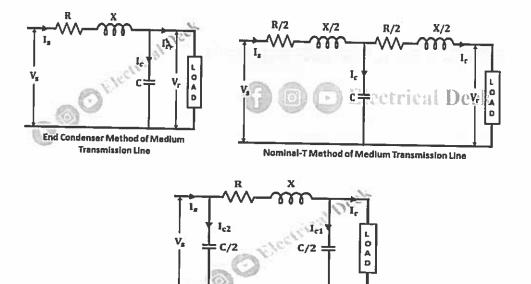
Medium Transmission Line:

Length of the overhead transmission line is in the range of 50-150km

✓ the operating voltage is greater than 20kV.

Based on the location of the capacitance at different places, the medium transmission lines have different configurations. These configurations show the different ways in which the effect of capacitance is taken into consideration. The three configurations based on the location of capacitance are,

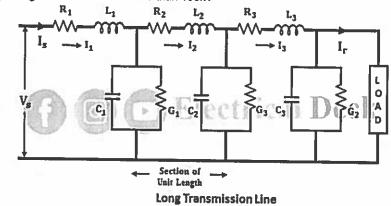
- i. End condenser representation of medium transmission lines
- ii. Nominal-T representation of medium transmission lines
- iii. Nominal-π representation of medium transmission lines.



Nominal-π Method of Medium Transmission Line

Long Transmission Line:

- ✓ Overhead transmission lines whose length is more than 150km
- Operating voltage of these lines is more than 100kV



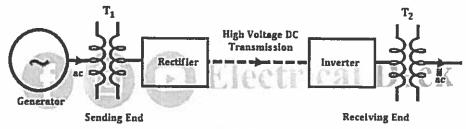
DC Transmission Line

In AC transmission lines, for transmission of power over long distances at higher voltages, the cost of transmission line and loss increases.

Also, the AC long transmission line suffers from problems like stability limits, voltage control, line compensation, interconnection of lines, ground impedance, etc due to an increase in voltage levels and distance.

The various problems associated with long-distance AC transmission have led to the development of HVDC (high voltage direct current) transmission nothing but a dc transmission line.

The use of DC power for long transmission lines has various advantages like no stability problem, absence of charging current, no skin effect, need for reactive compensation, bulk power transfer, economic power transmission, etc.



HVDC Transmission System

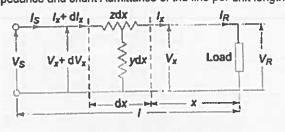
It is inefficient to use a dc transmission system for shorter and medium distances.

13. (a) Derive the expressions for the Performance of long transmission lines using rigorous method with relevant equations.

ABCD parameters in Long Transmission Line-rigorous Solution:

For rigorous solution, consider the Line Impedance and admittance uniformly distributed and not lumped.

Consider a small element of length dx situated at a distance x from receiving end. Let Z and Y denoted respectively the Series Impedance and shunt Admittance of the line per unit length.



Schematic diagram of a long line

The voltage at two ends of the element are denotes as V (towards receiving end) and V + dV (towards sending end) respectively.

Let dx be an elemental section of the line at a distance x from the receiving-end having a series impedance zdx and a shunt admittance ydx. The rise in voltage to neutral over the elemental section in the direction of increasing x is dV_x .

We can write the following differential relationships across the elemental section:

$$dVx = l_x z dx \text{ or } \frac{dV_x}{dx} = z l_x$$
(5.14)

$$dIx = V_x y dx$$
 or $\frac{dI_x}{dx} = y V x$ (5.15)

It may be noticed that the kind of connection (e.g. T or π) assumed for the elemental section, does not affect these first order differential relations. Differentiating Eq. (5.14) with respect to x, we obtain

$$\frac{\mathrm{d}^2 V_s}{\mathrm{d}x^2} = \frac{\mathrm{d}I_s}{\mathrm{d}x} z$$

Substituting the value of dl_x/dx from Eq. (5.15), we get

$$\frac{\mathrm{d}^2 V_x}{\mathrm{d}x^2} = y z V_x \tag{5.16}$$

This is a linear differential equation whose general solution can be written as follows:

$$V_{r} = C_{1}e^{\gamma t} + C_{2}e^{-\gamma t}$$
(5.17)

where

 $\gamma = \sqrt{yx}$ (5.18)

and C1 and C2 are arbitrary constants to be evaluated.

Differentiating Eq. (5.17) with respect to x:

$$\frac{dV_x}{dx} = C_1 \gamma e^{\gamma x} - C_2 \gamma e^{-\gamma x} = z J_x$$

$$I_x = \frac{C_1}{Z_r} e^{\gamma x} - \frac{C_2}{Z_r} e^{-\gamma x} \qquad (5.19)$$

Where

$$Z_c = \left(\frac{z}{y}\right)^{1/2} \tag{5.20}$$

The constants C_1 and C_2 may be evaluated by using the end conditions, i.e. when x = 0, $V_x = V_R$ and $I_x = I_R$. Substituting these values in Eqs. (5.17) and (5.19) gives

$$W_R = C_1 + C_2$$

 $I_R = \frac{1}{Z_r} (C_1 - C_2)$

which upon solving yield

$$C_{1} = \frac{1}{2} (V_{R} + Z_{r}J_{R})$$
$$C_{2} = \frac{1}{2} (V_{R} - Z_{r}J_{R})$$

With C₁ and C₂ as determined above, Eqs. (5.17) and (5.19) yield the solution for V_x and I_x as

$$V_{x} = \left(\frac{V_{R} + Z_{c}I_{R}}{2}\right)e^{\gamma x} + \left(\frac{V_{R} - Z_{c}I_{R}}{2}\right)e^{-\gamma x}$$
$$I_{x} = \left(\frac{V_{R}/Z_{c} + I_{R}}{2}\right)e^{\gamma x} - \left(\frac{V_{R}/Z_{c} - I_{R}}{2}\right)e^{-\gamma x}$$
(5.21)

Here Z_c is called the characteristic impedance of the Long Transmission Line and γ is called the propagation constant.

Knowing V_R , I_R and the parameters of the line, using Eq. (5.21) complex number rms values of V_x and I_x at any distance x along the line can be easily found out.

A more convenient form of expression for voltage and current is obtained by introducing hyperbolic functions. Rearranging Eq. (5.21), we get

$$V_x = V_R \left(\frac{e^{\gamma x} + e^{-\gamma x}}{2} \right) + I_R Z_c \left(\frac{e^{\gamma x} - e^{-\gamma x}}{2} \right)$$
$$I_x = V_R \frac{1}{Z_c} \left(\frac{e^{\gamma x} - e^{-\gamma x}}{2} \right) + I_R \left(\frac{e^{\gamma x} + e^{-\gamma x}}{2} \right)$$

These can be rewritten after introducing hyperbolic functions, as

$$V_x = V_R \cosh \gamma x + I_R Z_c \sinh \gamma x$$

$$I_x = I_R \cosh \gamma x + V_R \frac{1}{Z_c} \sinh \gamma x$$
(5.22)

where x = I, $V_x = V_s$, $I_x = I_s$

$$\therefore \qquad \begin{bmatrix} V_s \\ I_s \end{bmatrix} = \begin{bmatrix} \cosh \gamma l & Z_c \sinh \gamma l \\ \frac{1}{Z_c} \sinh \gamma l & \cosh \gamma l \end{bmatrix} \begin{bmatrix} V_R \\ I_R \end{bmatrix}$$
(5.23)

Here

$$A = D = \cosh \gamma l$$

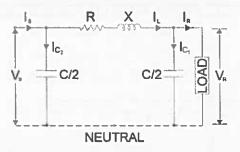
$$B = Z_c \sinh \gamma l$$

$$C = \frac{1}{Z_c} \sinh \gamma l$$

(5.24)

13. (b) Using nominal π method, derive an expression for sending end voltage and current for a medium transmission line.

In Nominal π Method, the shunt capacitance of each line i.e. phase to neutral is divided into two equal parts. One part is lumped at the sending end while the other is lumped at receiving end as shown in figure below.



Notice that, in this method there is no effect of shunt capacitance at sending end on the line voltage drop and hence on voltage regulation but this accounts for the charging current in sending end.

Let

IR = Load Current per phase

R = Resistance per phase

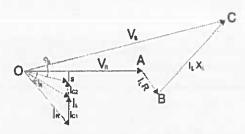
X = Reactance per phase C = Capacitance per phase

Cas QD = Dessiving and asw

Cos Ø R = Receiving end power factor (lagging)

Vs = Sending end voltage

Let us now draw the phasor. Assume receiving end voltage VR as reference and load current IR lagging this voltage by ØR.



Therefore, VR= VR+ j0 and IR= IR<-ØR = IR (CosØR – jSinØR) Charging Current at load end IC1 = j(ω C/2) VR = jmfCVR Line Current IL = IR + IC1 (phasor sum) Sending end voltage Vs = VR + IL (R+jX) Now, Charging current at sending end IC2 = j(ω C/2)Vs = jmfCVs Hence, Sending end current Is = IL + IC2 (phasor sum) Thus, sending end current and voltage is calculated as above and from these parameters the performance of line is evaluated.

14. Explain the different methods used to improve the string efficiency of insulators with necessary diagrams.

Methods of Improving String Efficiency

Method # 1. By Using Insulators with Larger Discs or by Providing Each Insulator Unit with a Metal Cap:

It is clear from the expression of string efficiency that the string efficiency increases with the decrease in value of K (i.e. the ratio of shunt capacitance to mutual capacitance). One method is to design the units such that the mutual capacitance (capacitance of each unit) is much greater than the shunt capacitance (capacitance to earth). This can be achieved by using insulators with larger discs or providing each insulator unit with a metal cap. The ratio K can be made 1/6 to 1/10 by this method.

Method # 2. By Using Longer Cross-Arms:

The ratio of shunt capacitance to mutual capacitance, K can alternatively be reduced by using longer crossarms so that the horizontal distance from line support (pole or tower) is increased thereby decreasing the shunt capacitance. But the limitations of cost and mechanical strength of line supports do not allow the crossarms to be too long and it has been found that in practice it is not possible to obtain the value of K less than 0.1.

Method # 3. By Capacitance Grading:

It is seen that non-uniform distribution of voltage across an insulator string is due to leakage current from the insulator pin to the supporting structure, which cannot be eliminated. However, it is possible that discs of different capacities are used such that the product of their capacitive reactance and the current flowing through the respective unit is same. This can be achieved by grading the mutual capacitance of the insulator units i.e., by having lower units of more capacitance—maximum at the line unit and minimum at the top unit, nearest to the cross-arm. By this method complete equality of voltage across the units of an insulator string can be obtained but this method needs a large number of different-sized insulator units and maintaining spares of all varieties of insulator discs. So this method is not used in practice below 200 kV.

Consider a 4-unit string. Let C be the capacitance of the top unit and let the capacitances of others units are C_2 , C_3 and C_4 , as shown in Fig.

Assume $C_1 = k C$

Applying Kirchhoff's first law to node A we get, $I_2 = I_1 + i_1$ => $\omega C_2 v = \omega C v + \omega C_1 v \text{ or } C_2 = C + K C = C (1 + K) ... (9.11)$ Applying Kirchhoff's first law to node B we get, $I_3 = I_2 + i_2$ or $\omega C_3 v = \omega C_2 v + \omega C_1 \times 2 v$ (or) $C_3 = C_2 + 2 K C = C (1 + K) + 2 K C = C (1 + 3 K) ... (9.12)$ Applying Kirchhoff's first law to node C we get, $I_4 = I_3 + i_3$ or $\omega C_4 v = \omega C_3 v + \omega C_1 \times 3 v$ or $\omega C_4 v = \omega C_3 v + \omega C_1 \times 3 v$ or $C_4 = C_3 + 3 K C = C (1 + 3 K) + 3 K C = C (1 + 6 K) ... (9.13)$

Thus, it will be possible to equalize the potential across all the units, if their capacitances are in the ratio of 1: (1 + K): (1 + 3 K): (1 + 6 K) and so on.

But in practice it is impossible to obtain such units which will have their capacitances in above ratio, although nearby results can be obtained by employing standard insulators for most of the units and employing larger units adjacent to the line.

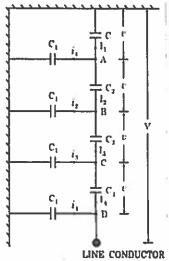
Method # 4. By Static Shielding:

In case of capacitance grading, insulator units of different capacitances are used so that the flow of different currents through the respective units produce equal voltage drop. In static shielding, pin to supporting structure charging currents are exactly cancelled so that the same current flows through the identical insulator units and produce equal voltage drops across each insulator unit.

This arrangement is given is Fig. 9.22. In this method a guard or grading ring, which usually takes the form of a large metal ring surrounding the bottom unit and electrically connected to the metal work at the bottom of this unit, and therefore to the line conductor.

The guard ring screens the lower units, reduces their earth capacitance C_1 and introduces a number of capacitances between the line conductor and the various insulator unit caps. These capacitances are greater for lower units and thus the voltages across them are reduced. With this method also it is impossible to obtain in practice an equal distribution of voltage but considerable improvements are possible.

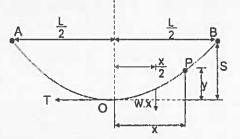
Let the capacitances between the links and the shield be C_x , C_y and C_z respectively as shown in Fig. 9.22, and let v be the potential across each unit.



Since the capacitance of each unit is same, therefore, their charging currents I1, I2, I3, and I4 would be same, let it be I. If C₁ = K C Applying Kirchhoff s first law to node A we get, $|+i_x = |+i_1 \text{ or } i_x = i_1... (9.14)$ Similarly, $i_y = i_2...$ (9.15) and $i_2 = i_3 \dots (9.16)$ Also, $i_1 = \omega C_1 v = \omega K C v ... (9.17)$ $i_2 = 2 \omega C_1 v = 2 \omega K C v \dots (9.18)$ $i_3 = 3 \omega C_1 v = 3 \omega K C v \dots (9.19)$ The potential causing current ix is 3 v (voltage across three units leaving the top one). So, $i_x = \omega C_x \times 3 v = 3 \omega C_x v ... (9.20)$ Comparing Eqs. (9.14), (9.17) and (9.20), we have, $3 \omega C_x v = \omega K C v \text{ or } C_x = K C/3 \dots (9.21)$ The potential causing current y is 2 y and therefore. $l_y = 2 \omega C_y v \dots (9.22)$ Comparing Eqs. (9.15), (9.18) and (9.22) we have, $2 \omega C_y v = 2 \omega K C v \text{ or } C_y = K C \dots (9.23)$ The potential causing current iz is v and therefore, $i_z = \omega C_z v \dots (9.24)$ SHIELD Comparing Eqs. (9.16), (9.19) and (9.24), we have, C. $\omega C_z v = 3 \omega K C v$ or C_z = 3 K C LINE CONDUCTOR In general, if there are n units $i_1 = \omega K C v$ and $i_x = (n - 1) \omega C_x v$ or $C_x = KC/(n-1)$ Similarly, $C_y = 2KC/(n-2)$ and $C_z = 3KC/(n-3)$

- or The capacitance of the pth metal link to the line is given as: $C_p = pKC/(n-p)$
- 15. (a) Derive an expression for sag of a line supported between two supports of the same tower height.
 Calculation of Sag in a Transmission Line:
 Sag calculation for supports is at equal levels

Suppose, AOB is the conductor. A and B are points of supports. Point O is the lowest point and the midpoint.



Let, L = length of the span, i.e. AB

w = weight per unit length of the conductor

T = tension in the conductor.

We have chosen any point on the conductor, say point P.

The distance of point P from the Lowest point O is x.

y is the height from point O to point P.

Equating two moments of two forces about point O as per the figure above we get,

$$Ty = wx \times \frac{x}{2}$$
$$Now, y = \frac{wx^2}{2T},$$

The maximum dip (sag) is represented by the value of y at either of the supports A and B. At support A, x = 1/2 and y = S.

Then $S = \frac{wL^2}{8T}$

15. (b) A 132 kV transmission line has the following data:
Wt. of conductor = 680 kg/km; Length of span = 260 m Ultimate strength = 3100 kg; Safety factor = 2. Calculate the height above ground at which the conductor should be supported. Ground clearance required is 10 m.

Sol: Weight of conductor/mt run,
$$W = \frac{680}{1000} = 0.68 \text{ kg}$$
.
Working Pension, $T = \frac{U1 \text{ trimate Strength}}{Safety factor}$
 $= \frac{3100}{2} = 1550 \text{ kg}$.

Span Length,
$$l = 260 \text{ m}$$
.
 $\therefore \text{ Sag} = \frac{W l^2}{8T} = \frac{0.68 \times 260^2}{8 \times 1550}$

$$= 3.7777$$

$$= 3.7777$$

$$= 10 + 3.7 = 13.777$$

$$= 13.777$$

$$= 13.777$$

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Semester End Regular/Supplementary Examination, Dec./Jan., 2022-2023

제한 이번 Net Market Mar

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Degree	B. Tech. (U. G.) Program	ECE		Provide and the local	Academic Year	2022 - 2023
Course	Code 20EC305	Test Duration	3 Hrs.	Max. Mar	ks 70	and a second sec	· III
Course	Digital Syster	n Design				n haan aa ahaa ahaa ahaa ahaa ahaa ahaa	
Dart A (Chart Annuar Quantia	no E u O n 40 Marta					
rait A (Short Answer Questio	ns 5 x 2 = 10 Warks,	1				
No.	Questions (1 through 5		24			Learning Outcome (s)	DoK
1	Draw and write the trut	h table for EXOR gal	ie.			20EC305.1	L1
2	Give an example for M	in term and Max tern	1 involving	3 Variables	5.	20EC305.2	- 11
3	What is Carry Depende					20EC305.3	L1
4	What is Race round co	ndition?		Total and the second second	Tertereriterer	20EC305.4	L1
5	Expand VHDL.					20EC305.5	L1
Part B (Long Answer Question	ns 5 x 12 = 60 Marks	5)				
No.	Questions (6 through 1			•	Marks	Learning Outcome (s) DoK
6 (a)	Convert the following ((iii) 695 ₁₀ =() ₁₆ .		2023 . 1997 - 1997 - 1997 - 1997 - 1997 - 1997 - 1997 - 1997 - 1997 - 1997 - 1997 - 1997 - 1997 - 1997 - 1997 - 1997 -		6M	20EC305.1	L2
6 (b)	Perform the given sul methods: (11100011) ₂	otraction using 1's a -(01110000)2.	nplement	6M	20EC305.1	L2	
			OR			Construction of the second	
7	Explain basic gates wi	-	12M	20EC305.1	L2		
8	Implement AND gate	OR gate, EXOR g	jate and N	Vot gate	1014	0050005.0	10
U	using NAND gate only	the second second			12M	20EC305.2	L2
9 (a)	Simplify the expressio	n Y = (A+B) (A'+C)/B'	OR +C')		6M	20EC305.2	10
	Minimize the following	function using Karna	ugh man ti	echnique	UIVI	2020303.2	L2
9 (b)	$f(A,B,C,D) = \Sigma_m(5,6,7,1)$	$2,13$)+ $\Sigma_{d}(4,9,14,15)$	ացո ուզի ն	connique	6M	20EC305.2	L2
10 (a)	Design the full adder u	sing two half adders.			6M	20EC305.3	L3
10 (b)	Design a 4-bit BCD ad			11550 - P	6M	20EC305.3	L3
		A REAL PROPERTY AND A REAL PROPERTY AND A REAL PROPERTY.	OR	en ar e ar			EO
11 (a)	Differentiate PAL, PRC Show and implement	DM, and PLA.	usino a Pl	ROM	6M	20EC305.3	L2
l1 (b)	$F(w,x,y,z) = \Sigma_m(1,9,11,$	12.13.15)	uonig u ri	(OIM	6M	20EC305.3	× L3
.,	$G(w,x,y,z) = \Sigma_m(0,1,2,3)$		14,15)		0141	2010303.3	L3
2 (a)	Explain the working of	JK Flip Flop.			6M	20EC305.4	L2
l2 (b)	Explain the Conversion	n of SR Flip Flop to T	Flip Flop. OR	1	6M	20EC305.4	L2
13	Design a 3 bit up coun	ter using D-flip flop.		······	12M	20EC305.4	L3
14	Explain the program significance of entity a	structure of VHDL	and Exp	plain the	12M	20EC305.5	L2
15	List and discuss variou		OR		12M	20EC305.5	4 1

AC 15 00 2024 Question Peper for End Semester Exemination 1 Academic Regulation 2020

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N S RAJU INSTITUTE OF TECHNOLOGY (AUTONOMOUS)

SONTYAM , ANANDAPURAM, VISAKHAPATNAM - 531 173

Degree - B. Tech - ECE ANSWER KEY AND SCHEME OF EVALUATION AC: 2022-2023 coubise Code- 20EC305 DIGITAL SYSTEM DESIGN SEM : III PART-A 1. Draw and write the truth table for EXOR gate (2M) XOR gate Symbol - Im truth table - im 2. Give an example for min and Max term involving 3 Variables (2m) Minterm example - Im Maxterim example - im 3. What is cavery dependency - 2m Cavury dependency - 2m 4. What is race around Condition - 2m Place abound Condition - 2m 5. Expand VHDL VHDL abrevation - 2m PART-B Ga) Convert the following Gm i) $ABC_{16} = (1_{10} - 2m)$ ii) $154_{8} = (1_{10} - 2m)$ iii) $695_{10} = (1_{10} - 2m)$ 6(b) Perform the given subtraction Using is & 2's compliment methods - Gm 15 Compliment method -3m 2's Compliment method - 3m 7) Explain Basic gates with their truth table AND gate - 4m (Symbol - Im, expression - Im, OR gate - 4m explaination - Im, truth table - Im) NOT gate - 4m

8. Implement AND gate, OR gate, EXOR gate, and NOT gate Using NAND gates only. AND gate stealisation - 3m. OR gate stealisation - 3m. EXOR gate stealisation - 4m. NOT gate stealisation - 2m.
q(a) Simplify the expression $\gamma = (A+B)(A+C)(C-C)$
(6) minimize the following using K-112,13) \rightarrow (6m) $f(A,B,C,D) = \sum_{m} (5,6,7,12,13) + \sum_{d} (4,9,14,15) \rightarrow (6m)$
10 a) Design the full adder using two half uator full adder truth table - (2m) design a full adder - (4m)
106) Design a 4 bit BCD adder Circuit -(2m) Truth table for BCD adder circuit -(2m) expressions. (2m)
11(a) Differentiate PAL, PROM, and The Contract
11(b) Show and implement the following guinent
implementations (un) 12(a) explain the working of JK flip flop (6m) Circuit diagram (2m) Explaination (1m) Truth table (1m)
characteristic tubic (m) exclusion
126) explain Conversion Of SR FF to T FF 6m characteristic table - T - IM simplification - IM
excitation table - SR-IM Circuit - 2m
13(a) Design a 3bit up Counter Using D-tt -12m 13(a) Design a 3bit up Counter Using D-tt -12m State table-2m, Circuit diagram -4m, timing diagram - 4m explaination -2m.
19. explain the program stratice of more capating of the
Polognam Structure -2m, entity Syntax - 3m, Architecture Syntax - 3m Significance - 4m
15. list and discuss various data types in VHD1 with explains exam data types list (Bm) each data type explaination -(9m)



N S RAJU INSTITUTE OF TECHNOLOGY (AUTONOMOUS) SONTYAM, ANANDAPURAM, VISAKHAPATNAM – 531 173

ANSWER KEY AND SCHEME OF EVALUATION

PART-A:

Draw and write the truth table for EXOR gate. A

A	В	A⊕B
0	0	0
0	ł	1
1.	0	1
1 3	1	0

2. Give an example for Minterm and Mark term involving 3 Variable.

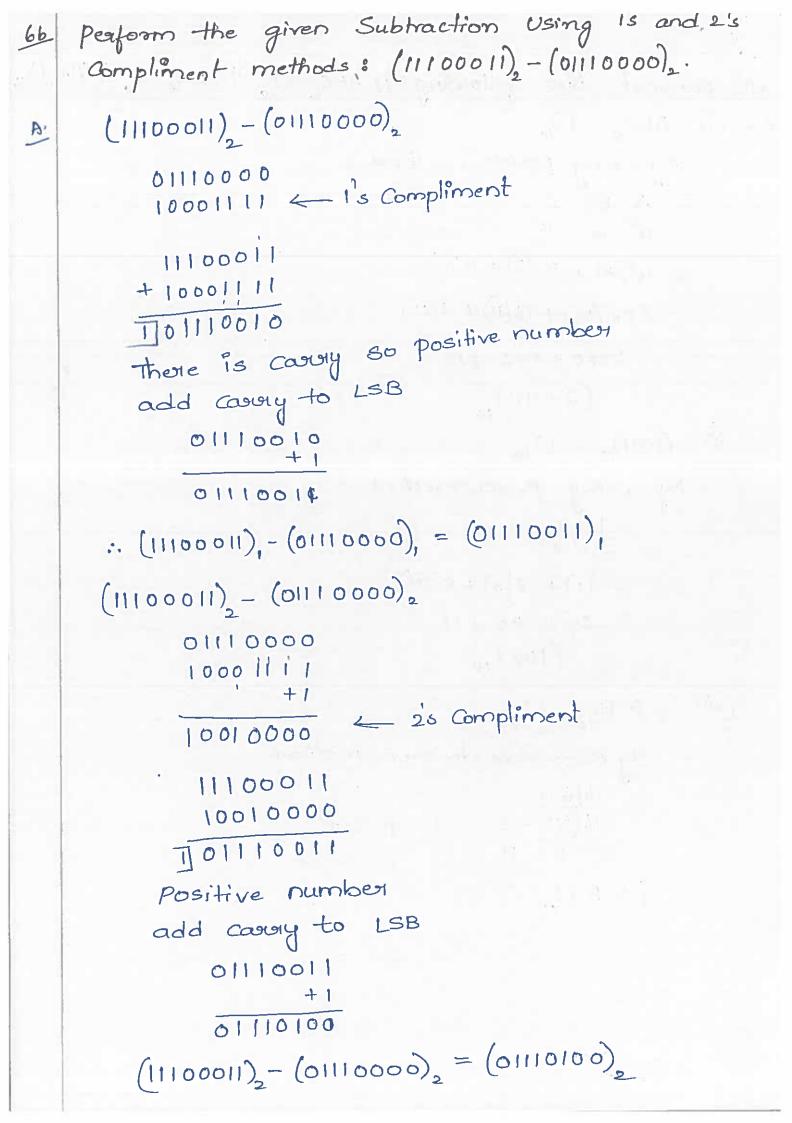
A. Minterm: $f(A, B, c) = (\Xi m(1, 3, 5, 7))$

Y= ABC + ABC + ABC + ABC
max term: f(A,B,C) = TTm(1,3,5,7)
Y = (A+B+Z)(A+B+C)(A+B+C)(A+B+C)(A+B+C)
What is carry dependency?
In Adder circuit next stage circuit will depende on previous stage carry. To generate the next stage outputit has to wait for previous stage carry. This dependency is called as carry dependency. A) What is Race around Condition?
A) Place around Condition in JK flipflop, if J=1 and K=1 and if Clk=1 for a long period of time, then Q output will toggle as long as CLK is high, which makes the output of the flip-flop Unstable or Uncertain. This problem is called slace around Condition.

5) Expand VHDL ? VHDL — VHSIC HDL Very high speed integrated circuit hardware description language.

and the product of the state of

b = b



7. Explain basic gates with their truth tables I AND Gate: A B AB=Y An AND gate is an electrical circuit that combines two signals so that the output is on if both Signals are present. The output of the AND gate is connected to the base driver which is coupled to the bases of transistors, and alternately Switches the transistors at opposite corners of the inverter. logic expression for AND gate is Y = A·B Y = output Variable A, B are input variable Truth table :-A B Y=AB 0 0 0 0 1 0 14.1 0 0 1 A _____ A+B : y 2. OR gate :-

An OR gate is a digital logic gate that gives an output of 1 when any of its inputs are I. Otherwise O. An OR gate performs like two switches in parallel supplying a light, so that when either of the Switches is closed the light is On. OR gates Can have more than two inputs. logic expression for OR gate is Y = A+B Y = output Variable A, B are input Variables

Touth table :-

A	B	Y=A+B	
0	0	0	
0	1		
j F	0		
1	1	ł	

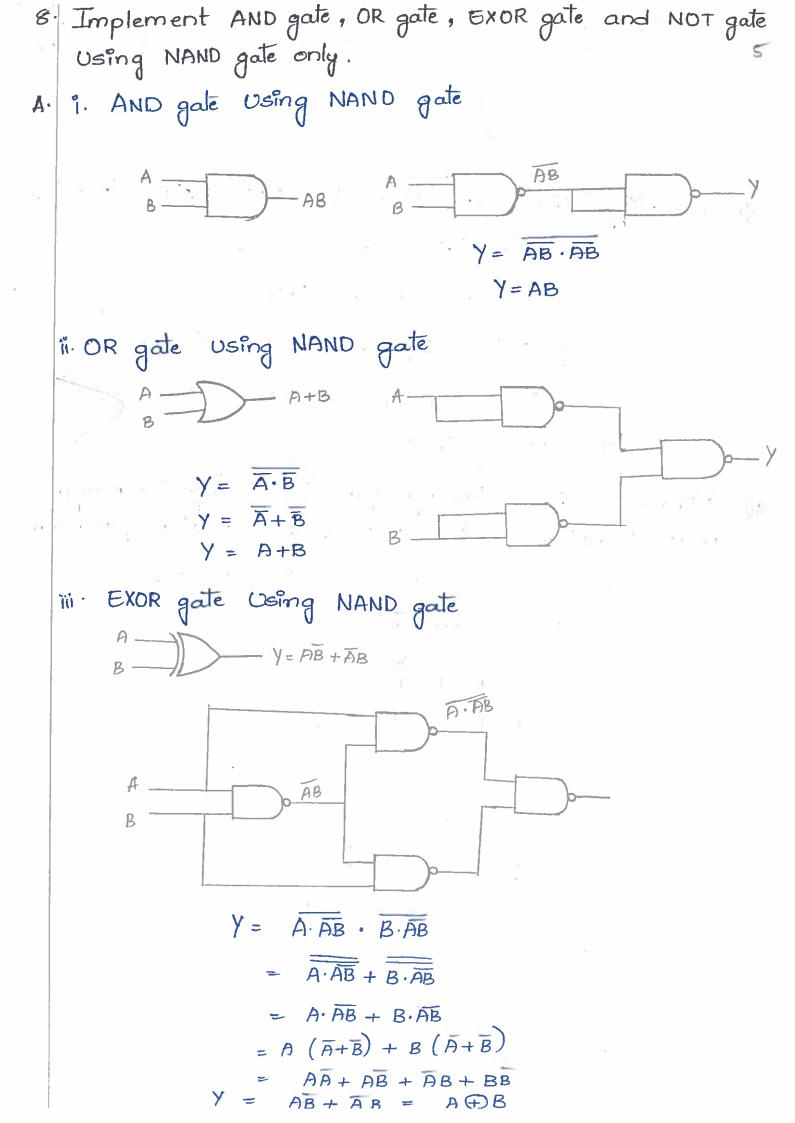
Not gate :- A _____A

The NOT gate is an electronic circuit that produces an invested version of the input at its output. It is also known as an invester. If the input Variable is A, the invested output is known as NOT A. This is also known as \overline{A} or $\overline{A'}$. logic expression for NOT gate is $\overline{A} = \overline{A}$ $Y = \overline{A}$

South and a straight of a public hydrogenet. Additional fills

Truth table :-

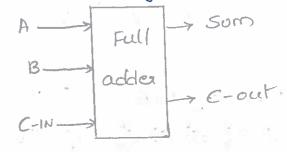




iv Not gate using NAND gate:
A
$$A\overline{P} = \overline{P}$$

AF = \overline{P}
Simplify the expression $Y = (A+B)(A^{1}+c)(B^{1}+c^{1})$.
A $Y = (A+B)(A^{1}+c)(B^{1}+c^{1})$
 $= (AA^{1}+A^{1}B + Ac+Bc)(B^{1}+c^{1})$
 $= AA^{1}B^{1} + A^{1}B^{1} + AB^{1}c + BB^{1}c + AB^{1}c^{1} + BB^{1}c^{1} + AB^{1}c^{1} + BB^{1}c^{1} + BB^$

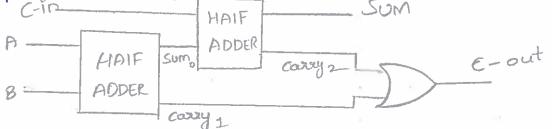
100 Design the full adder Using two half adders. 6 Full adder is the adder that adds three inputs and produces two outputs. The first two inputs are A and B and thrind input is input cavery as C-IN. The autput caving is designated as c-out and the normal output is designated as S which is SUM.



A.

r.	INPUTS		OUTF	ru t
A	B	C-IN	Sum	c-out
0	0	0	0	O
0	0	1	ж. П. П. С.	0
0		0	1	0,4
0		10	0	l
1	0	0		0
le la companya de la companya	0	1	0	t
ľ	l l	0	0	1
ŀ		= 1,	÷	

full Adder Using two half adders: 2 Half address and an OR gate is required to implement a full adder.



With this logic circuit, two bits can be added togethey taking carry from the next lower order of magnitude and sending a cavery to the next higher order of magnitude.

10b Design a 4-bit BCD adder Circuit

BCD adder refers to a 4-bit binary adder that Can add two 4-bit words of BCD format. The output of the addition is a BCD format 4-bit output word, which defines the decimal sum of the addend & augend and a cavity that is created in case this sum exceeds a decimal value of 9.

Decimal	B	ina	ry	9u	m	B	BCD	Su	m	
Value	C_3^*	53	52	S1*	S.	С	53	s,	S	, S.
0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	0	t	0	0	0	0	- 1
2	0	0	0	1	0	0	0	0	- t	0
3	0	0	0	1	1	0	б	٥	I.	1
4	0	0	1	0	0	0	0	Þ	0	0
5	0	0	T	0	1	0	0	L	0	J
6	0	0	E	t	0	0	0	1	1	0
Т	0	0	1	1	l t	0	0	1	1	1
8	0	E	0	0	0	0	1	0	O	0
9	0	1	0	0	1	0	1	0	0	1
10	0	T.	0	E	0	1	0	0	0	0
1)	0	t	0	1	1	1	0	0	0	
12	0	t	t	0	0		0	0	1	0
13	0	1	E	0	1		0	0	1	-
14	0	1	1	T.	0		0	1	0	0
15	0	1	i.	1	1			1		
16							D		0	
	t	0	0	0	0	l	0	t) (2
17	t	0	0	0	1	t	0	ł	1	1
18	t	0	0	I	0	1	t	0	0 0	
โล	t	0	D	ť	t	t	1	0	0	1

A.

1st Condition $\rightarrow C^* = 1$ 2^{2nd} Condition $\rightarrow S_3^* \left(S_2^* + S_1^* \right)$ Correction :- $C^{*} + S_{3}^{*}(S_{2}^{*} + S_{1}^{*})$ $= C^{*} + S_{3}^{*} S_{2}^{*} + S_{3}^{*} S_{1}^{*}$ B3 B2 B, B, B A3 A2 A1 A0 4 - bit binary Adder Č* 53 Sz* S S. 4 - bit binary adden S, č s, S. } BCD Sum S,

11 a.	Differentiate PAI	L, PROM, and PLA.	
	PROM	PAL	PLA.
	1. The Decoder (AND auray) împlements all minterims	1 The AND average implements a limited no. of product terms.	1. The AND avoray implements a limited no poroduct terms.
	2. AND asuray is not programmable	AND avoiay is pologrammable	AND avoray is priogram mable
	3. OR Array is pologrammable	3. OR array is not programable	OR auray is Priogrammable.
	4. Inter Connection aure more.	are medium	Inter Connections are less.
	5. cheap and Simple 6. least tlexible	cheap and simple than PLA. Moderate flexible	to PROM & PAL.
116	Show and imple $F(W, X, Y, Z) = \Sigma_{rr}$	ment of following	function Using PROM
	W 4x X DECODE Y Z Z		
		F(w,x,y,z) $G(w,x,y,z)$	

12a Ezplain the working of Jk flip flop.

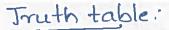
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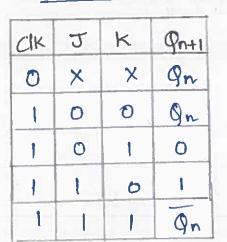
A J-k flip flop is nothing more than S-R flip flop with an added layer of feedback. This feedback Selectively enables one of the two Set / Heset inputs So that they cannot both cavery on active signal to the multivibrator Circuits, thus eliminating the invalid Condition.

1. When J=0, K=0when J=0, k=0 it is memory 9=0 Q =1 3. When JEI, K=0 when J=0, K=121 OP OP 9=0 Q = 0 0 0 $\overline{Q} = 1$ $\overline{Q} = 1$ 9 = 0 0 9 =1 0 $\overline{Q} = 1$ 0 $\overline{Q} = 0$ when J=1, K=14. 0 P

q = 0 P = 0 q = 0 Q = 1 Q =

when J=1, K=1 then the output is inverted or toggled state.

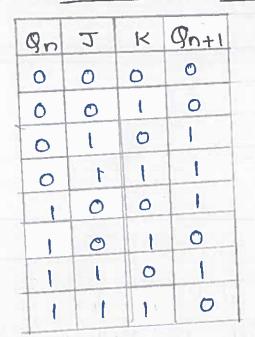




excitation table

9n+1	Qn	J	K
0	0	0	×
0	1	×	1
1	0	1	×
1	1	×	0

Characteristic table:



126) Explain the Conversion of SR flip-flop to T-flip flop. A. Avoilable flip flop - SR Grequired flip flop - T excitation table for available flip-flop - SR

Qn	QnH	S	R	
0	0	٥	×	
ଡ଼	1	t	0	
t	0	0	1	
1	1	×	0	

Characteristic table for required flip-flop-T

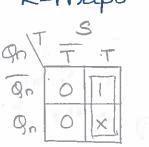
Combine both the tables Of SR and F

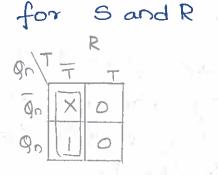
9n	Т	Qn+1	S	R
0	0	0	0	×
0	1	ł	ł	0
1	0	1	0	1
١	J	0	×	0

Draw K-maps

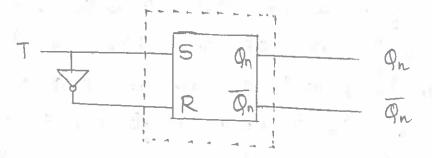
13

Ar





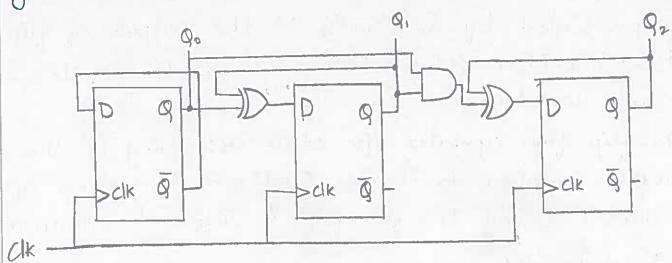
S=T R=T



Design a 3 bit up counter Using D-flip flop. Consider a 3-bit counter with each bit count depresented by Po, 9, 1 9, as the outputs of flip flops FFo, FF1, FF2, despectively. Then the state table would be D-flip flop updates the state according to the input applied to it i.e g=D. So designing up Counter Using D-flip flop is different than a T-flip flop.

state	92	9,	9.	
0	0	0	0	
1	0	0	1	ĺ
2_	0	1	0	1
3	0	T.	t	
4	1	0	0	
4 5	1	0	1	
6	t	1	0	
T	t	t	1	ľ

According to the State table of up-Counter Q_0 is Continuously changing so the input to FF, will be $D_0 = \overline{Q}_0 \cdot Because$ it will toggle the state whenever $Q_1 = 1$, when its previous State $Q_1 \notin Q_0$ are not equal \notin $Q_1 = 0$, when its previous State $Q_1 \notin Q_0$ are equal. That is the same as XOR operation. So $D_1 = Q_1 XOR Q_1$ $Q_2 = 1$, when in its previous state the AND of $Q_1 \notin Q_0$ is not equal to $Q_2 \cdot Q_2 = 0$, when in its previous state th AND of $Q_1 \notin Q_0$ is equal to $Q_2 \cdot SO D_2 = Q_2 XNOR (Q, \notin Q)$ Schematic of Synchronous up-counter using D-H is given below.



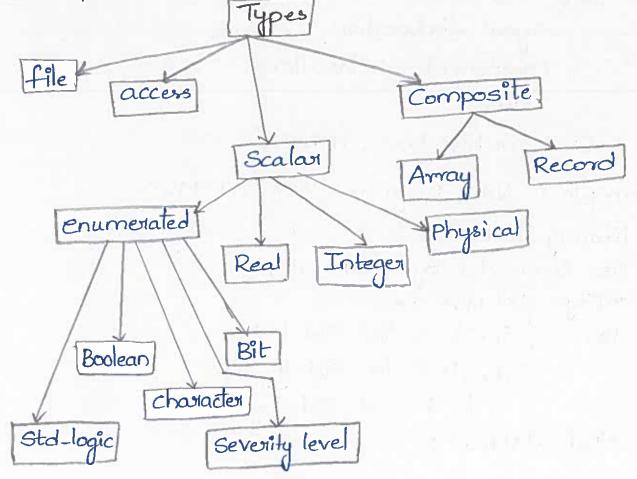
14 Explain the priogram structure of VHDL and explain the significance of entity and architecture. PHOgram Structure Of VHDL Ð library declaration; Entity Entityname 18 Port (port_name : Signal-mode Signal-type; Port-name : Signal-mode Signal-type); END Entity name; Architecture aschitecture - name Of entity-name is (declariations) Begin (Code) END auchitecture - Hame; Signal declaviation Component declasiation Body End architecture _ name ; Example :- VHDL program for 4X1 MUX. library ieee; Use leee. std-logic-1164. all; entity AXI MUX is Port (SI, So : in Std-logic ; A, B, C, D : in Std-logic; Y: out Std-logic); end 4x1 mux;

auchitecture 4x1MUX-1 of 4x1MUX 18 Begin With select 5,5. Y & A when "00", B when "01", C when "10", D when "11", U when others;

end 4x1 MUX=1;

5

A VHDL models Consist of an entity declaration and a architecture Body. The entity defines the interface, the architecture defines the function. The entity declaration names the entity and defines the interface to its environment. List and discuss Various data types in VHDL, with examples.



Scalar types: Scalar types describe objects that can hold, at most, one value at a time. The type itself can Contain multiple values, but an object that is declared to be a Scalar type can hold, at most, one of the Scalar Values at any point in time. 1. Integen types 2. Real types 3. Enumerated types 4. Physical types 5. floating point Enumerated Data types An enumerated data type is a very powerful tool for abstract modelling. A designer can use an enumerated type to represent exactly the values required for a specific operation. 1. Boolean 2. character 3. Bit 4. Std-logic 5. Severity level 1. Boolean This data type is used when we need to convey Some true or false Conditions. Eq: Architecture ---Begin Process (----) if acb then temp & FRue; else temp & False; end if end Process;

Character: This data type is used when we need to use all alpha numeric and special characters. Bit: This data type is used when we need to suppresent binary values ('d' and '1'). Severity level This data type is used in Complex projects when we need to show warnings, everors in suntime, failure in suntime. Std_llogic :-This data types are declared in std-logic-1164. all Package of Ieee library. Integer Data type: Integers are exactly like mathematical integers. All of the normal predefined mathematical function like add, subtract, multiply, and divide apply to integer types Eg:- 1) Type - integer declaration type < word lengt > is stange 0 to 31; 2) Object - integer declaration Constant < loop number >: < integer > <= 345; Real Data type:" Real types are used to declare objects that emulate mathematical steal numbers. They can be used to represent numbers out of the range of integer Values as well as fractional values.

Physical Data types:

physical types are used to represent physical Quantities Buch as distance, Cworent, time and so on A physical type provides for a base Unit and Successive Units are then defined in terms of this Unit.

Eq: Type Current is stange 0 to 10000000 Units na; -- nano amps Ua = 1000 na; -- micro amps ma = 1000 Ha; -- milli amps a = 1000 ma; -- amps end Units; Prepared by

DEF E.Manemma Assf. Poof ECE

July A1/23



4

Semester End Regular/Supplementary Examination, Dec./Jan., 2022-2023

Degree	B. Tech.	Program	CSE, CSE (AI & ML) & CSE	(DS)	Academic Year	2022 - 2023
Course Code	20CS305			Max. Marks		Semester	111
Course	Computer Orga	nization			1		

1

a.

Part A	(Short Answer Questions 5 x 2 = 10 Marks)	<u></u>		
No.	Questions (1 through 5)		Learning Outcome (s)	Del
1	What is the need of a shift Register?	20CS305.1	Dok	
2	What is an Interrupt?	20CS305.1		
3	What is the main purpose of the RISC?		20CS305.2	
4	What is signed magnitude method?			
5	What are Peripherals?		20CS305.4	<u></u>
Part B	(Long Answer Questions 5 x 12 = 60 Marks)		20CS305.5	<u> L1</u>
No.	Questions (6 through 15)	A.4		
6 (a)	Distinguish between fixed point representation and floating	Marks	Learning Outcome (s)	Dok
	point representation.	6M	20CS305.1	L2
6 (b)	Define Decoder. Explain about 3 - to - 8 - line decoder.	6M	20CS305.1	-
7 (-)	ÓR	0		
7 (a)	Explain about the error detection codes.	6M	20CS305.1	L2
7 (b)	With a neat sketch explain about multiplexers.	6M	20CS305.1	L2
0				
8	Explain arithmetic, logic and shift micro-operations.	12M	20CS305.2	L2
	OR			
9 (a)	Explain memory reference instructions with an example.	5M	20CS305.2	L2
9 (b)	Describe the various phases involved in the instruction cycle.	7M	20CS305.2	L2
10	What do you mean by addressing mode? Explain the addressing modes.	12M	20CS305.3	L2
	OR OR			
11 (a)	How address sequencing is done in micro programmed control?	6M	20CS305.3	L2
11 (b)	Explain about stack organization.	6M	20CS305.3	L2
	Explain the bordware implementation of the term			
12 (a)	Explain the hardware implementation of signed magnitude addition and subtraction.	8M	20CS305.4	L2
12 (b)	Draw and explain the flowchart for booth multiplication algorithm.	4M	20CS305.4	L2
:	OR			
13 (a)	What is overflow and underflow in floating point arithmetic?	6M	20CS305.4	110
13 (b)	Explain the difference between signed and unsigned division.	6M	20CS305.4	L2
			2006303,4	L2
14 (a)	Explain the methods employed for establishing priority using Daisy – Chaining priority.	6M	20CS305.5	L2
	Describe the vertice commend in the terms		200000,0	64
14 (b)	Describe the various components like input/output units, memory unit, control unit, arithmetic logic unit connected in the basis emergination of a second	6M	20CS305.5	L2
	the basic organization of a computer.	I		
5 (0)	OR OR			
15 (a)	Explain the method of DMA transfer in a computer system.	5M	20CS305.5	L2
15 (b)	Explain the concept of virtual memory. Why it is significant?	7M	20CS305.5	L2

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N S RAJU INSTITUTE OF TECHNOLOGY

(AUTONOMOUS)

SONTYAM , ANANDAPURAM, VISAKHAPATNAM - 531 173

ANSWER KEY AND SCHEME OF EVALUATION

Semester End Regular/Supplementary Examination, Dec./Jan., 2022-2023

. 61.1		B. Tech.	Program		(AI & ML) & CSE	(DS)	Academic Year	2022 - 202
	rse Code rse	20CS305 Computer Or	Test Duration ganization	3 Hrs.	Max. Marks	70	Semester	HE
Part No. 1	Questio	Answer Questic ns (1 through 5) the need of a si		s)		Le	arning Outcome (s) 20CS305.1	DoK L1
	Definitio	n-1M					2000000.1	L 1
	Need-1	м			\mathcal{L}		<u>6</u> .	
	The shif	it registers are i	used for temporary (data storad	e. The shift regi	stere a	ire also used for de	to transfe
	and data time dela	a manipulation. ay to digital circu	The serial-in serial-o	out and para	illel-in parallel-o	ut shift	registers are used t	lo produce
2		an Interrupt?					20CS305.2	L1
	Definitio							
	Interrupt	s are the signals	s generated by the e	external devi	ces to request t	he mic	roprocessor to perfo	rm a task.
	so an in	iterrupt in com	outer architecture is	s a signal t	hat requests th	e proc	essor to suspend i	its curren
	execution	n and service th	e occurred		-	·		
			eeodiiod		34		7.8	
3		he main purpos				÷	20CS305.3	L1
3		he main purpos			a T	:)=(-	20CS305.3	L1
3	What is t	he main purpos n-1 M			• 		20CS305.3	L1
3	What is t Definition Purpose-	he main purpos n-1 M 1M	e of the RISC?	microproce	ssor that is desire	aned to		
3	What is t Definition Purpose- RISC (red	he main purposi n-1 M 1M duced instruction	e of the RISC?	microproce: an operate	ssor that is designated at a higher spo	gned to) perform a smaller r	number of
3	What is t Definition Purpose- RISC (red types of	he main purposi n-1 M 1M duced instruction computer instru s the analytical p	e of the RISC?	an operate	at a higher spo	eed. S	o perform a smaller r o the processor are	number of
3	What is t Definition Purpose- RISC (red types of that shifts compile ti	he main purposi n-1 M 1M duced instruction computer instru s the analytical p	e of the RISC? In set computer) is a uctions so that it ca process of a comput	an operate	at a higher spo	eed. S	o perform a smaller r o the processor an runtime to the prepa	number of
	What is t Definition Purpose- RISC (red types of that shifts compile ti	he main purposi in-1 M 1M duced instruction computer instru- s the analytical p ime.	e of the RISC? In set computer) is a uctions so that it ca process of a comput	an operate	at a higher spo	eed. S	o perform a smaller r o the processor are	number of
	What is t Definition Purpose- RISC (red types of that shifts compile ti What is si Definition	he main purposite -1 M 1M duced instruction computer instru- the analytical p ime. igned magnitude -2M	e of the RISC? In set computer) is a auctions so that it ca process of a comput e method?	an operate ational task	at a higher spo from the execu	eed. S tion or	o perform a smaller r o the processor an runtime to the prepa 20CS305.4	number of chitecture aration or L1
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	What is t Definition Purpose- RISC (red types of that shifts compile ti What is si Definition- In the sig negative m number.	he main purposition -1 M duced instruction computer instru- the analytical p ime. igned magnitude -2M ned magnitude	e of the RISC? In set computer) is a actions so that it ca process of a comput e method? method number is c	an operate ational task fivided into	at a higher spo from the execu two parts: Sign	eed. S tion or bit and	o perform a smaller r o the processor an runtime to the prepa 20CS305.4 d magnitude. Sign b ed with the binary fo	number of chitecture aration or L1 it is 1 for rm of the
	What is t Definition Purpose- RISC (red types of that shifts compile ti What is si Definition- In the sig negative m number.	he main purposition 1M 1M duced instruction computer instru- is the analytical p ime. igned magnitude -2M ned magnitude humber and 0 for Peripherals?	e of the RISC? In set computer) is a actions so that it ca process of a comput e method? method number is c	an operate ational task fivided into	at a higher spo from the execu two parts: Sign	eed. S tion or bit and	o perform a smaller r o the processor an runtime to the prepa 20CS305.4	number of chitecture aration or L1 it is 1 for

read information from or write in the memory unit on receiving a command from the CPU. They are considered to be a part of the total computer system. As they require a conversion of signal values, these devices can be referred to as electromechanical and electromagnetic devices.

Part B (Long Answer Questions 5 x 12 = 60 Marks)

No.	Questions (6 through 15)	Marks	Learning Outcome (s)	DoK
	Distinguish between fixed point representation and floating	6M	20CS305.1	L2
6 (a)	point representation.			

Difference-6M

Fixed point representation	Floating point representation
Fixed point is a representation of real data type for a number that has a fixed number of digits after the radix point.	Floating point is a formulaic representation of real numbers as an approximation so as to support a tradeoff between range and precision.
While fixed point can be used to represent a limited range of values	floating point can be used to represent a wide range of values.
The performance of the fixed point is higher than floating point.	The performance of the fixed point is higher than floating point.
Floating point representation is more flexible than fixed point representation.	Floating point representation is more flexible than fixed point representation.
Higher Performance	Less Performance
Less flexibility Define Decoder. Explain about 3 – to – 8 – line decoder	More flexible er. 6M 20CS305.1 L1

6 (b)

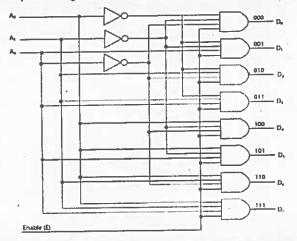
Define Decoder. Explain about 3 – to – 8 – line decoder. Definition of Decoder- 2M

Explanation for 3 - to - 8 - line decoder & Diagram - 4 M

A decoder is a combinational circuit that modifies binary data from n input lines to a maximum of 2n unique output lines. An encoder creates the binary code corresponding to the input activated. A decoder gets a set of binary inputs and activates only the output that complements that input number.

<u>3 - to - 8 - line decoder</u>: The most preferred or commonly used decoders are n-to-m decoders, where m<= 2^n . An n-to-m decoder has n inputs and m outputs and is also referred to as an n * m decoder.

The following image shows a 3-to-8 line decoder with three input variables which are decoded into eight output, each output representing one of the combinations of the three binary input variables.



The three inverter gates provide the complement of the inputs corresponding to which the eight AND gates at the output generates one binary combination for each input. The most common application of this decoder

L2

20CS305.1

Definition of Error detection codes- 2 M

Explanation & Diagram -4M

Error-detecting codes are a sequence of numbers generated by specific procedures for detecting errors in data that has been transmitted over computer networks.

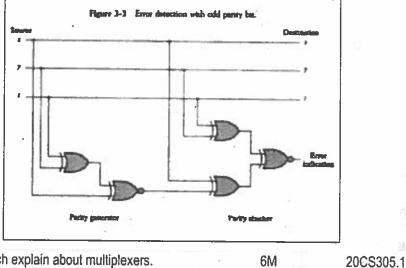
6M

Parity Check

Parity check is done by adding an extra bit, called parity bit to the data to make number of 1s either even in case of even parity, or odd in case of odd parity.

While creating a frame, the sender counts the number of 1s in it and adds the parity bit in following way

- In case of even parity: If number of 1s is even then parity bit value is 0. If number of 1s is odd then
 parity bit value is 1.
- In case of odd parity: If number of 1s is odd then parity bit value is 0. If number of 1s is even then
 parity bit value is 1.



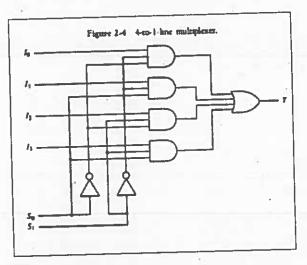
7 (b) With a neat sketch explain about multiplexers. 6M 20CS Definition of Multiplexers- 2 M Explanation & Diagram -4M

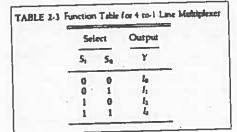
A Multiplexer (MUX) can be described as a combinational circuit that receives binary information from one of the 2ⁿ input data lines and directs it to a single output line. The selection of a particular input data line for the output is decided on the basis of selection lines.

A 4-to-1-line multiplexer is shown in Fig. 2-4. Each of the four data inputs I0 through I, is applied to one input of an AND gate. The two selection inputs 51 and 50 are decoded to select a particular AND gate. The outputs of the AND gates are applied to a single OR gate to provide the single output. To demon strate the circuit operation, consider the case when 5150 = 10. The AND gate associated with input I, has two of its inputs equal to 1. The third input of the gate is connected to I2• The other three AND gates have at least one input equal to 0, which makes their outputs equal to 0. The OR gate output is now equal to the value of I,, thus providing a path from the selected input to the output.

The 4-to-1 line multiplexer of Fig. 2-4 has six inputs and one output. A truth table describing the circuit needs 64 rows since six input variables can have 2 6 binary combinations. This is an excessively long table and will not be shown here. A more convenient way to describe the operation of multiplexers is by means of a function table. The function table for the multiplexer is shown in Table 2-3. The table demonstrates the relationship between the four data inputs and the single output as a function of the selection inputs s1 and

s0.





8 Explain arithmetic, logic and shift micro-operations.

12M

20CS305.2

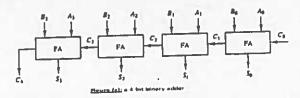
L2

Arithmetic Microoperations -4M Logic Microoperations -Explanation-3M, Applications-1M

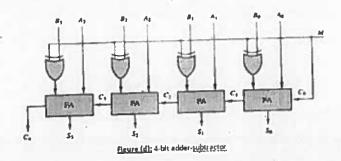
Shift Microoperations -4M

The basic arithmetic micro operations are addition, subtraction, increment, decrement, and shift. Arithmetic shifts are explained later in conjunction with the shift micro operations. The arithmetic micro operation defined by the statement R3 <-R1+R2.

Binary Adder: To implement the add micro operation with hardware, we need the registers that hold the data and the• digital component that performs the arithmetic addition. The digital circuit that forms the arithmetic sum of two bits and a previous carry is called a full-adder. The digital circuit that generates the arithmetic sum of two binary numbers of any lengths is called a binary adder. The below figure (c) shows the interconnections of four full-adders (FA) to provide a 4-bit binary adder.



Binary Adder-Subtractor The subtraction of binary numbers can be done most conveniently by means of complements. A - B can be done by taking the 2's complement of B and added with A



Binary Incrementer: The increment micro operation adds one to a number in a register. For example, if a 4bit

to

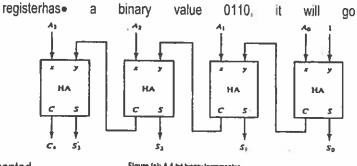
0111

after

32

it.

is



incremented.

Eloure (e): A 4-bit binary Incr

LOGIC MICROOPERATIONS Logic micro operations specify binary operations for strings of bits stored in registers. These operations - consider each bit of the register separately and treat them as binary variables. For example, the exclusive-OR micro operation with the contents of two registers R1 and R2is symbolized. by the statement,

 $P: R1 \leftarrow R1 \oplus R2$

List of Logic Microoperations: There are 16 different logic micro operations that can be performed with two binary variables.

The 16 logic microoperations are derived from these functions by replacing variable x by the binary content of register A and variable y by the binary content of register B.

Boolean function	Microoperation	Name
$F_{\rm e} = 0$.	<i>F</i> ← 0	Clear
$F_1 = xy$ $F_2 = xy'$	FHAAB FHAAB	AND
$F_3 = x$ $F_4 = x^* y$	F-A F-A AB	Transfer A
$F_1 = y$	F +- B	Transfer B
F. = x 🕀 y	F-ABB	Exclusive-OR
$F_{y} = x + y$	F-AVB	OR
$F_{\pm} = (x + y)^{\prime}$	E-AVB	NOR
$F_{n} = (x \oplus y)^{\prime}$	F+-ABB	Exclusive-NOR
$F_{10} = y^*$ $F_{11} = x + y^*$	F-AVB	Complement B
$F_{2} = x'$ $F_{0} = x' + y$	F-AVB	Complement A
$F_{\rm so} = (xy)^{\prime}$	F-773	NAND
$F_{ij} = 1$	F all 1's	Set to all 1's

Logic microoperations are very useful for manipulating individual bits or a portion of a word stored in a register. They can be used to change bit values, delete a group of bits, or insert new bit values into a register.

- 1. selective-set (or) the OR microoperation
- 2. Selective-complement (or) the exclusive-OR microoperation.
- 3. selective-clear (or) the AB' microoperation.
- 4. Mask (or) AND microoperation

SHIFT MICROOPERATIONS

Shift micro operations are used for serial transfer of data. The contents of a register can be shifted to the left or the right. Determines the type of shift.

- There are three types of shifts:
- 1. Logical shift
- 2. Circular shift

3. Arithmetic shift

OR

9 (a) Explain memory reference instructions with an example. 5M 20CS305.2 L2 Explanation- 5M

MEMORY-REFERENCEINSTRUCTIONS:

AND : AND to AC This is an instruction that performs the AND logic operation on pairs of bits in AC and the memory word specified by the effective address. The result of the operation is transferred to AC.

AND : AND to AC This is an instruction that performs the AND logic operation on pairs of bits in AC and the memory word specified by the effective address.

LDA: Load to AC This instruction transfers the memory word specified by the effective address to AC.

STA: Store AC This instruction stores the content of AC into the memory word specified by the effective address.

BUN: Branch Unconditionally This instruction transfers the program to the instruction specified by the effective address.— The BUN instruction allows the programmer to specify an instruction out of sequenceandwe say— that the program branches (or jumps) unconditionally.

BSA: Branch and Save Return Address This instruction is useful for branching to a portion of the program called a subroutine or procedure. When executed, the BSA instruction stores the address of the next instruction in sequence (which is available in PC) into a memory location specified by the effective address. The effective address plus one is then transferred to PC to serve as the address of the first instruction in the subroutine.

ISZ: Increment and Skip if Zero This instruction increments the word specified by the effective address, and if the incremented value is equal to 0, PC is incremented by 1.

9 (b) Describe the various phases involved in cycle.	the instruction	7M ·	20CS305.2	L2
Explanation-2 M Fetch-1 M				
Decode-1M Read-1M	12	¥.	1 A A	
Execute-1M Instruction Cycle flow chart-1M, A program residing in the memory unit of the	computer consists o	f a sequence	of instructions. The	program is

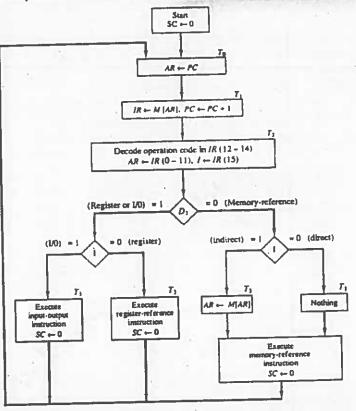
A program residing in the memory unit of the computer consists of a sequence of instruction cycle in turn is executed in the computer by going through a cycle for each instruction. Each instruction cycle in turn is subdivided into a sequence of subcycles or phases. In the basic computer each instruction cycle consists of the following phases:

1. Fetch an instruction from memory.

2. Decode the instruction.

3. Read the effective address from memory if the instruction has an indirect address.

4. Execute the instruction.



What do you mean by addressing mode? Explain the addressing modes. Definition-1 M

L2

L2

20CS305.3

Types of Addressing Modes: 11 M

Addressing mode is the way in which the location of an operand can be specified in an instruction. It generates an effective address (the actual address of the operand). Instruction format with mode field Types of Addressing Modes:

12M

1. Implied Mode 2. Immediate Mode 3. Register Mode 4. Register Indirect Mode: 5. Autoincrement or Autodecrement Mode 6. Direct Address Mode 7. Indirect Address Mode 8. Relative Address Mode 9. Indexed Addressing Mode 10. Base Register Addressing Mode.

OR

 11
 How address sequencing is done in micro programmed
 6M
 20CS305.3
 L2

 (a)
 control?
 Steps-6M

Microinstructions are stored in control memory in groups, with each group specifying a routine.

Step-1: An initial address is loaded into the control address register when power is turned on inthe computer. This address is usually the address of the first microinstruction that activates the instruction fetche routine. The fetch routine may be sequenced by incrementing the control address register through the rest of itse microinstructions. At the end of the fetch routine, the instruction is in the instruction register of the computer. 2 M

Step-2: The control memory next must go through the routine that determines the effective address of theo operand. A machine instruction may have bits that specify various addressing modes, such as indirect address and index registers. The effective address computation routine in control memory can be reached through a branch microinstruction, which is conditioned on the status of the mode bits of the instruction. When the effective address computation routine is completed, the address of the operand is available in the memory address register. 2M

Step-3: The next step is to generate the microoperations that execute the instruction fetched frommemory. The microoperation steps to be generated in processor registers depend on the operationcode part of the instruction. Each instruction has its own micro-program routine stored in a given location of controlmemory. The transformation from the instruction code bits to an address in control memory wherethe routine is located is• referred to as a mapping process. A mapping procedure is a rule that transforms the instruction code into a controlmemory address• 2 M

Step-4: Once the required routine is reached, the microinstructions that execute the instructionmay be sequenced by• incrementing the control address register. Micro-programs that employ subroutines will require an external register for storing thereturn address. Return addresses cannot be stored in ROM because the unit has no writing capability. When the execution of the instruction is completed, control must return to the fetchroutine. This is accomplished by executing an unconditional branch microinstruction to the firstaddress of the fetch• routine. 2M

6M

11 (b)

Explanation: 3 M Daigram-1M Push-1M

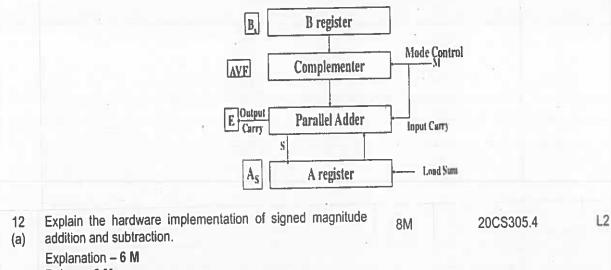
Explain about stack organization.

Pop-1M

STACK ORGANIZATION: A useful feature that is included in the CPU of most computers is a stack or lastin, first-out (LIFO) list. A stack is a storage device that stores information in such a manner that the item stored last is the first item retrieved. The operation of a stack can be compared to a stack of trays. The last tray placed on top of the stack is the first to be taken off. The register that holds the address for the stack is called a stack pointer (SP) because its value always points at the top item in the stack. The two operations of a stack are the insertion and deletion of items. 1. Push or push-down (insertion operation) 2.

Pop or pop-up (deletion operation)

Hardware Algorithm for Signed-magnitude addition and subtraction

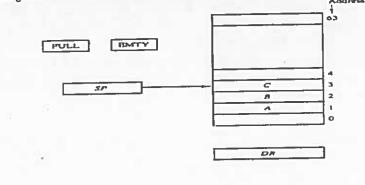


Daigram-2 M Addition and Subtraction with Signed-Magnitude Data

The algorithms for addition and subtraction stated as follows (the words inside + parentheses should be used for the subtraction algorithm): Addition (subtraction) algorithm. • 1. when the signs of A and B are identical (different), add the two magnitudes and attach the sign of A to the result. 2. When the signs of A and B are different (identical), compare the magnitudes and 3. subtract the smaller number from the larger. Choose the sign of the result to be the same as A if A > B or the complement of the sign of A if A < B. If the two magnitudes are equal, subtract B from A and make the sign of the result positive.

Hardware Implementation for Addition and Subtraction with Signed-Magnitude Data

To implement the two arithmetic operations with hardware, it is first necessary that the two- numbers be stored in registers. Let A and B be two registers that hold the magnitudes of the numbers, and As, and Bs between flip-flops that hold the corresponding signs. Here parallel-adder is needed to perform the microoperation A + B. (Consists of Full adder). The complementer for generating the 2's Complement while performing subtraction operation... (Consists of X-OR gate). Where M is Mode of operation. When M = 0, the output of B is transferred to the adder, the input carry is 0, and the adder is equal to the sum A + B. When M =1, the I's complement of B is applied to the adder, the input carry is 1, and • output is = A + + 1. This is equal to A plus the 2's complement of B, which is equivalent to the A - B. Figure 1: Hardware for signed-magnitude addition and subtraction. A+ B. where EA is a register that the magnitudes are added with a microoperation EA combines E and A. The value of E is transferred into the Add-overflow flip-flop AVF, if E is 1.• The magnitudes are subtracted by adding A to the 2's complement of B•



Draw and explain the flowchart for booth multiplication 12 **4**M algorithm.

20CS305.4

L2

(b)

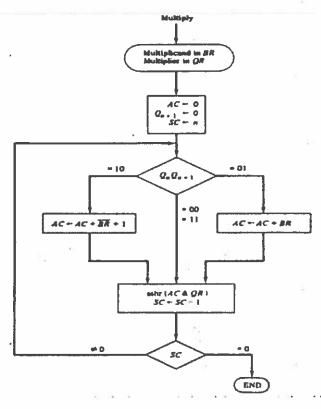
Flow chart-2 M

Explanation-2M

Booth Multiplication Algorithm (for signed-2's complement numbers)

The flowchart for Booth algorithm is shown in below figure. AC and the appended bit Qn+1 are initially

cleared to 0 and the sequence counter SC is set to a • number n equal to the number of bits in the multiplier. The two bits of the multiplier in Qn and Qn+1 are inspected. • If the two bits are equal to 10, it means that the first 1 in a string of 1's has been encountered. • This requires a subtraction of the multiplicand from the partial product in AC. If the two bits are equal to 01, it means that the first 0 in a string of 0's has been encountered. • This requires the addition of the multiplicand to the partial product in AC. When the two bits are equal, the partial product does not change. An overflow cannot occur because the addition and subtraction of the multiplicand follow each other. The next step is to shift right the partial product and the multiplier (including bit Qn+1). This is an • arithmetic shift right (ashr) operation which shifts AC and QR to the right and leaves the sign bitin AC unchanged. The sequence counter is decremented and the computational loop is repeated n times.





 13
 What is overflow and underflow in floating point arithmetic?
 6M
 20CS305.4
 L2

 (a)
 Explanation of overflow-2 M

 Example- 1M

 Explanation of Underflow-2 M

Example-1M

floating-point number in computer registers consists of two parts: a mantissa m and an exponent e. The two parts represent a number obtained from multiplying m times a radix r raised to the value of e; thus m x re. The mantissa may be a fraction or an integer. The location of the radix point and the value of the radix r are assumed and are not included in the registers. For example, assume a fraction representation and a radix 10. The decimal number 537.25 is represented in a register with m = 53725 and e = 3 and is interpreted to represent the floating- point number 0.53725×10^{4} . A floating-point number is normalized if the most

significant digit of the mantissa is nonzero. In this way the mantissa contains the maximumpossible number of significant digits. A zero cannot be normalized because it does not have a nonzero digit. It is represented in floating-point by all 0's in the mantissa and exponent.

A floating-point number that has a 0 in the most significant position of the mantissa is said to have an underflow. To normalize a number that contains an underflow, it is necessary to shift the mantissa to the left and decrement the exponent until a nonzero digit appears in the first position. In the example above, it is necessary to shift left twice to obtain .35000 x 10 3 . In most computers, a normalization procedure is performed after each operation to ensure that all results are in a normalized form.

When two normalized mantissas are added, the sum may contain an overflow digit. An overflow can be corrected easily by shifting the sum once to the right and incrementing the exponent. When two numbers are subtracted, the result may contain most significant zeros.

Explain the difference between signed and unsigned division. 6M 20CS305.4

L2

13 (b)

Signed Division-3 M Unsigned Division-3 M

The Division of two fixed-point binary numbers in the signed-magnitude representation is done by the cycle of successive compare, shift, and subtract operations. Division of two fixed-point binary numbers in signed-magnitude representation is done with paper and pencil by a process of successive compare, shift, and subtract operations. Binary division is simpler than decimal division because the quotient digits are either 0 or 1 and there is no need to estimate how many times the dividend or partial remainder fits into the divisor. The division process is illustrated by a numerical example in Figure. The divisor B consists of five bits and the dividend A, of ten bits. The five most significant bits of the dividend are compared with the divisor. Since the 5- bit number is smaller than B, we try again by taking the six most significant bits of A and compare this number with B. The 6-bit number is greater than B, so we place a 1 for the quotient bit in the sixth position above the dividend. The divisor is then shifted once to the right and subtracted from the dividend. The divisor is then shifted once to the right and subtracted from the dividend. The divisor is then shifted once to the right and subtracted from the dividend. The divisor is then shifted once to the right and subtracted from the dividend. The divisor is then shifted once to the right and subtracted from the dividend. The divisor is then shifted once to the right and subtracted from the dividend. The divisor is then shifted once to the right and subtracted from the dividend. The divisor is then shifted once to the right and subtracted from the dividend. The divisor is called a partial remainder because the division could have stopped here to obtain a quotient of 1 and a remainder equal to the partial remainder. The process is continued by comparing a partial remainder with the divisor. If the partial remainder is greater than or equal to the divisor, the quotient bit is equal to 1.

14Explain the methods employed for establishing priority using
Daisy – Chaining priority.6M20CS305.5L2

Explanation-4 M

Daigram-2 M

The daisy-chaining method of creating priority includes a serial connection of all devices that request an interrupt. The device with the highest priority is located in the first position, followed by lower-priority devices up to the device with the lowest priority, which is situated last in the chain. This technique of connection between three devices and the CPU.

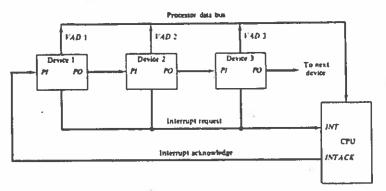


Figure 11-12 Duisy-chain priority interrupt

14 Describe the various components like input/output units,

(b) memory unit, control unit, arithmetic logic unit connected in the basic organization of a computer.

20CS305.5

L2

the basic organization of a Input /Output Unit-2 M Memory Unit- 1 M control unit – 1 M arithmetic logic unit-1 M Diagram -1 M

Input /Output Unit : These components help users enter data and commands into a computer system. Data can be in the form of numbers, words, actions, commands, etc

6M

Memory Unit : Once a user enters data using input devices, the computer system stores this data in its memory unit.

control unit:

This unit is the backbone of computers. It is responsible for coordinating tasks between all components of a computer system.

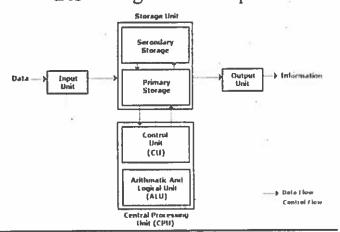
Output Unit

11

The third and final component of a computer system is the output unit. After processing of data, it is converted into a format which humans can understand.

arithmetic logic unit

This part of the CPU performs arithmetic operations. It does basic mathematical calculations like addition, subtraction, division, multiplication, etc.



Block diagram of computer

OR

15 (a)

5M

L2

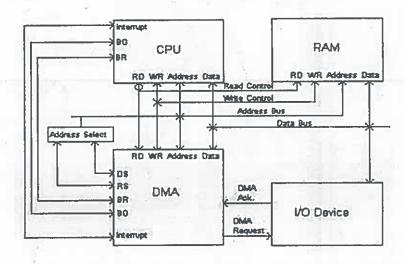
20CS305.5

DMA-1M

Explanation- 3 M

Daigram-1 M

DMA basically stands for Direct Memory Access. It is a process which enables data transfer between the Memory and the IO (Input/ Output) device without the need of or you can say without the involvement of CPU during data transfer.



15 (b)

Explanation-4 M

Diagram- 2 M

Significance-1 M

Virtual memory is a memory management technique where secondary memory can be used as if it were a part of the main memory. Virtual memory is a common technique used in a computer's operating system (OS).

7M

Significance :

It can handle twice as many addresses as main memory.

Explain the concept of virtual memory. Why it is significant?

- It enables more applications to be used at once.
- It frees applications from managing shared memory and saves users from having to add memory

modules when RAM space runs out.

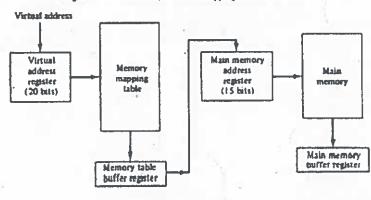


Figure 12-17 Memory table for mapping a virtual address.

